Type Inference Functional Programming SS2007

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The correct way of labeling

letrec $all = \operatorname{fun} l \to \operatorname{match} l$ with $[] \to \operatorname{true} |$ true $:: xs \to all xs |$ false $:: xs \to \operatorname{false} ;;$

is

$$\operatorname{letrec} \underbrace{all}_{\alpha_1} = \overbrace{\operatorname{fun} \underset{\alpha_2}{l} \rightarrow \operatorname{match} \underset{\alpha_2}{l} \xrightarrow{\alpha_1} \underset{\alpha_3}{l} \xrightarrow{\alpha_1} \underset{\alpha_4}{l} \underbrace{\operatorname{true}}_{\alpha_5} \underset{\alpha_6}{l} \underbrace{\operatorname{true}}_{\alpha_6} \xrightarrow{\alpha_8} \underset{\alpha_9}{l} \underbrace{\operatorname{all}}_{\alpha_9} \underset{\alpha_{10}}{x_{10}} \underset{\alpha_{11}}{\underline{\alpha_{11}}} \xrightarrow{\alpha_{13}} ;;$$

The first two of occurrences of xs get the label α_6 because the second occurrence is bound to the first. The third occurrence is independent of the other two, so it gets a fresh label.

Similarly, a correct labeling of

letrec
$$s = (\lambda x.x z) (\lambda x.x z) ;;$$

is

letrec
$$s_{\alpha_1} = \overbrace{(\lambda_{\alpha_2}}^{\alpha_2} \cdot \underbrace{x_{\alpha_2}}_{\alpha_2} \cdot \underbrace{z_{\alpha_4}}_{\alpha_5} \cdot \underbrace{(\lambda_{\alpha_3}}_{\alpha_3} \cdot \underbrace{x_{\alpha_2}}_{\alpha_7} \cdot \underbrace{z_{\alpha_6}}_{\alpha_7});;$$

The first and second occurrences of x are independent of the third and fourth. The occurrences of z are dependent, but z is global and therefore get new labels for every occurrence.