

Automatic Deduction — Introduction to Isabelle

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1 Two Grammars

The most natural definition of valid sequences of parentheses is this:

$$S \rightarrow \epsilon \mid '(S)'\mid SS$$

where ϵ is the empty word.

A second, somewhat unusual grammar is the following one:

$$T \rightarrow \epsilon \mid T'(T)'$$

▷ Make the definitions of sets S and T in Isabelle and give structured Isar proofs leading to the theorem $S = T$.

Hint: use lists over a specific type to represent words.

The alphabet:

```
datatype alpha = A | B
```

2 Polynomial sums

▷ Produce structured proofs of the following theorems, using induction and calculational reasoning in Isar.

Note that the given tactic scripts are of limited use in reconstructing structured proofs; nevertheless the hints of automated steps below can be re-used to finish sub-problems. The \sum symbol can be entered as “\`\<Sum>`”; note that numerals in Isabelle/HOL are polymorphic.

```
theorem — problem  
  fixes n :: nat  
  shows "2 * ( $\sum$  i=0..n. i) = n * (n + 1)"  
  by (induct n) simp_all
```

```
theorem — problem  
  fixes n :: nat  
  shows " $(\sum$  i=0..<n. 2 * i + 1) = n2"
```

```
by (induct n) (simp_all add: power_eq_if nat_distrib)
```

theorem — problem

```
fixes n :: nat
```

```
shows " $(\sum_{i=0..<n} 2^i) = 2^n - (1::nat)$ "
```

```
by (induct n) (simp_all split: nat_diff_split)
```

theorem — problem

```
fixes n :: nat
```

```
shows " $2 * (\sum_{i=0..<n} 3^i) = 3^n - (1::nat)$ "
```

```
by (induct n) (simp_all add: nat_distrib)
```

theorem — problem

```
fixes n :: nat
```

```
assumes k: "0 < k"
```

```
shows " $(k - 1) * (\sum_{i=0..<n} k^i) = k^n - (1::nat)$ "
```

```
by (induct n) (insert k, simp_all add: nat_distrib)
```