Experiments in Verification

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Simplification On One Slide

- basic methods: `simp`, `simp_all`
- simp-modifiers: `add:` ⟨*thms*⟩, `del:` ⟨*thms*⟩, `only:` ⟨*thms*⟩, `(no_asm)`, `(no_asm_simp)`, `(no_asm_use)`
- modifying the simpset: **`declare`** ⟨*thm*⟩`[simp]`, **`declare`** ⟨*thm*⟩`[simp del]`
- unfolding specific simp-rules: **`unfolding`** ⟨*thms*⟩

This Time

Session 1
formal verification, Isabelle/HOL basics, functional programming in HOL

Session 2
simplification, function definitions, induction, calculational reasoning

Session 3
natural deduction, propositional logic, predicate logic

Session 4
sets, relations, inductively defined sets, advanced topics

Isabelle's Meta-Logic

Description
minimal intuitionistic higher-order logic

Connectives
- $\bigwedge$: universal quantifier
- $\Longrightarrow$: implication
- $\equiv$: equality

Example

$$\bigwedge x\ y.\ x \equiv y \Longrightarrow y \equiv x$$

Some Remarks

Schematic Variables

free variables and (meta) universally quantified variables (at the outermost level) are both turned into schematic variables after a proof

Meta-Equality

in almost any case, equality (=) may be used instead of meta-equality (≡)

Meta-Implication

- nested implications associate to the right and
- may be abbreviated by $[\![A_1; \ldots; A_n]\!] \Longrightarrow B$ instead of $A_1 \Longrightarrow \ldots \Longrightarrow A_n \Longrightarrow B$
- **assumes** A **shows** B is turned into $A \Longrightarrow B$ after a proof

Natural Deduction

Inference Rules

- $\dfrac{A_1 \quad \ldots \quad A_n}{B}\,\langle name\rangle$
- premises $A_1, \ldots, A_n$
- conclusion B

In Isabelle

$$\textbf{theorem}\ \langle name\rangle\text{:}\ [\![A_1; \ldots; A_n]\!] \Longrightarrow B$$

resulting in

$$[\![?A_1; \ldots; ?A_n]\!] \Longrightarrow ?B$$

Example

Conjunction Rules and An Easy Proof

$$\frac{\phi \quad \psi}{\phi \wedge \psi}\wedge\mathrm{i} \qquad \frac{\phi \wedge \psi}{\phi}\wedge\mathrm{e}_1 \qquad \frac{\phi \wedge \psi}{\psi}\wedge\mathrm{e}_2$$

1	$p \wedge q$	premise
2	r	premise
3	q	$\wedge e_2$ 1
4	p	$\wedge e_1$ 1
5	$q \wedge r$	$\wedge i$ 3, 2
6	$p \wedge (q \wedge r)$	$\wedge i$ 4, 5

The Same Rules in Isabelle

conjI: $[\![?P; ?Q]\!] \Longrightarrow ?P \wedge ?Q$

conjunct1: $?P \wedge ?Q \Longrightarrow ?P$

conjunct2: $?P \wedge ?Q \Longrightarrow ?Q$

The Method `rule`

- synopsis: `rule` ⟨*name*⟩
- applies to a goal provided it is the instance of the conclusion of ⟨*name*⟩
- solves the goal if there are current facts that are instances of the premises of ⟨*name*⟩
- the number and order of those facts has to be exactly the same as for the premises of ⟨*name*⟩

The Above Proof in Isabelle

State What You Want To Prove

```
lemma
  assumes pq: "p ∧ q" and "r"
  shows "p ∧ (q ∧ r)" (is ?goal)
```

Prove It

```
proof -
  from pq have "q" by (rule conjunct2)
  from pq have "p" by (rule conjunct1)
  moreover
  from ‘q‘ and ‘r‘ have "q ∧ r" by (rule conjI)
  ultimately
  show ?goal by (rule conjI)
qed
```


Some Notes

- referring to facts is possible via name (if one was defined), e.g., **from** pq ...
- or by explicitly writing the fact between backticks (this is then called a literal fact), e.g., **from** ‘q‘ ...
- for every term (between double quotes) an abbreviation can be introduced using an is-pattern, e.g., "$p \wedge (q \wedge r)$" (is ?goal)
- **moreover** is used to collect a list of facts
- afterwards the list is used by **ultimately**

Introduction/Elimination Rules

Idea

For every logical connective there are several rules for introducing it and for eliminating it.

Natural Deduction - Propositional Logic

$$\frac{\phi \quad \psi}{\phi \wedge \psi}\,(\wedge\mathrm{i}) \qquad \frac{\phi_i}{\phi_1 \vee \phi_2}\,(\vee\mathrm{i}_i) \qquad \frac{\boxed{\begin{array}{c}\phi\\ \vdots\\ \psi\end{array}}}{\phi \to \psi}\,(\to\mathrm{i}) \qquad \frac{\boxed{\begin{array}{c}\phi\\ \vdots\\ \bot\end{array}}}{\neg\phi}\,(\neg\mathrm{i})$$

$$\frac{\phi_1 \wedge \phi_2}{\phi_i}\,(\wedge\mathrm{e}_i) \qquad \frac{\phi \vee \psi \quad \boxed{\begin{array}{c}\phi\\ \vdots\\ \chi\end{array}} \quad \boxed{\begin{array}{c}\psi\\ \vdots\\ \chi\end{array}}}{\chi}\,(\vee\mathrm{e}) \qquad \frac{\phi \to \psi \quad \phi}{\psi}\,(\to\mathrm{e}) \qquad \frac{\neg\phi \quad \phi}{\psi}\,(\neg\mathrm{e})$$

Some Derived Rules

Double Negation Introduction

$$\frac{\phi}{\neg\neg\phi}\,(\neg\neg\mathrm{i})$$

Proof.

1	ϕ	premise
2	$\neg\phi$	assumption
3	$\bot$	$\neg$e 2, 1
4	$\neg\neg\phi$	$\neg$i 2–3

□

Some Derived Rules (cont'd)

Law Of The Excluded Middle

$$\frac{}{\phi \vee \neg\phi}\ (\mathsf{lem})$$

Proof.
Exercise □

Some Derived Rules (cont'd)

Double Negation Elimination

$$\frac{\neg\neg\phi}{\phi}\ (\neg\neg\mathsf{e})$$

Proof.

1	$\neg\neg\phi$	premise
2	$\phi \vee \neg\phi$	lem
3	ϕ	assumption
4	$\neg\phi$	assumption
5	ϕ	$\neg$e 1, 4
6	ϕ	$\vee$e 2, 3, 4–5

□

Some Derived Rules (cont'd)

Proof By Contradiction

$$\frac{\boxed{\begin{array}{c}\neg\phi\\ \vdots\\ \bot\end{array}}}{\phi}\ (\mathsf{pbc})$$

Proof.

1	$\neg\phi$	assumption
$\vdots$	$\vdots$	
n	$\bot$	
$n+1$	$\neg\neg\phi$	$\neg$i 1–n
$n+2$	ϕ	$\neg\neg$e $n+1$

□

A Word On Destruction Rules

Loosing Information

- usually rules like $\wedge\mathsf{e}_1$ are known as elimination rules
- in Isabelle they are called destruction rules
- using such rules destroys information
- thus it can turn a goal unprovable
- use destruction rules with care

Example (Conjunction Elimination)

$$\frac{\phi \wedge \psi \qquad \boxed{\begin{array}{c}\phi\\ \psi\\ \vdots\\ \chi\end{array}}}{\chi}\ (\wedge\mathsf{e})$$

Raw Proof Blocks

In-Place Proofs

- enclose between { and }
- does not work on current goal but introduces new facts
- any '**assume**'s are premises of the resulting fact
- the last '**have**' is the conclusion of the resulting fact
- like boxes in the 'pen 'n' paper' natural deduction rules

Universal Quantification

Introduction and Elimination Rules

$$\frac{\boxed{\begin{array}{l} x_0 \\ \quad\vdots \\ \phi(x_0) \end{array}}}{\forall x.\ \phi(x)}\ (\forall i) \qquad \frac{\forall x.\ \phi(x)}{\phi(t)}\ (\forall e)$$

Isabelle Idiom for Meta Universal Quantification

$$\textbf{fix}\ x_0\ \ldots\ \textbf{show}\ \text{"}?P(x_0)\text{"}\ \langle proof \rangle$$

results in

$$\bigwedge x.\ ?P(x)$$

Existential Quantification

Introduction and Elimination Rules

$$\frac{\phi(t)}{\exists x.\ \phi(x)}\ (\exists i) \qquad \frac{\exists x.\ \phi(x) \qquad \boxed{\begin{array}{c} x_0\ \phi(x_0) \\ \vdots \\ \psi \end{array}}}{\psi}\ (\exists e)$$

Isabelle Idiom For ∃-Elimination

$$\text{"}\exists x.\ ?P(x)\text{"}\ \textbf{then obtain}\ y\ \textbf{where}\ \text{"}?P(y)\text{"}\ \langle proof \rangle$$

results in

$$?P(y)$$

An Example Proof

```
lemma
  assumes ex: "∃x. ∀y. P x y"
    shows "∀y. ∃x. P x y"
proof
  fix y
  from ex obtain x where "∀y. P x y" by (rule exE)
  hence "P x y" by (rule spec)
  thus "∃x. P x y" by (rule exI)
qed
```

Exercises

http://isabelle.in.tum.de/exercises/logic/elimination/ex.pdf
http://isabelle.in.tum.de/exercises/logic/propositional/ex.pdf
http://isabelle.in.tum.de/exercises/logic/predicate/ex.pdf