

Prototype Verification System

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Overview

- ▶ Facts
- ▶ Applications
- ▶ Telephone Book
- ▶ Subset Types
- ▶ Type Checking
- ▶ Abstract Data Types
- ▶ Dependent Types
- ▶ Propositional Proof
- ▶ Impressions



Facts

- ▶ **Implemented in Common Lisp**
 - ▶ Functions in C, Tcl/TK and LaTeX
 - ▶ Interface: GNU Emacs

- ▶ **Prototype Verification System**
 - ▶ specification language: strongly typed & based on classical higher-order logic
 - ▶ interactive theorem prover (or proof checker)

- ▶ **current version: PVS 6.0 (start 1993)**
 - ▶ open-source & GPL



Facts

- ▶ **developed by SRI**
 - ▶ Jovial Verification System, Hierarchical Development Model, STP, EHDM
 - ▶ Yices

- ▶ **current work: develop methodologies for**
 - ▶ highly automated hardware verification
 - ▶ integration with model checkers
 - ▶ for concurrent and real-time systems



Applications

“PVS has been installed at hundreds of sites in North America, Europe and Asia” [PVS]

- ▶ DB with reference to history, technology and applications of PVS: ~300 citations
- ▶ **NASA PVS Library**
 - ▶ “The major goals of our research program are to advance the state-of-the-art in formal methods, making it practical for use on life-critical systems developed by the aerospace industry in the United States.” [LFM]



Motivation

- ▶ **specification:** phone book [TUT]
 - ▶ A phone book shall store the phone numbers of a city.
 - ▶ It shall be possible to retrieve a phone number, given a name.
 - ▶ It shall be possible to add and delete entries from a phone book.

`FindPhone` `AddPhone` `DelPhone`



Types

- ▶ represent entities as distinguishable types

N: TYPE % names

P: TYPE % phone number

- ▶ uninterpreted types
 - ▶ distinguishable, nothing about members is known & equality predicate

B: TYPE = [N -> P] % phone book

- ▶ total function?



Types

```
B: TYPE = [N -> P]    % phone book
```

- ▶ For names without phone number introduce “n0”

```
n0: P
```

```
emptybook: B
```

```
emptyax: AXIOM
```

```
  FORALL (nm:N) : emptybook(nm) = n0
```



FindPhone & AddPhone

- ▶ find a phone number

FindPhone: $[B, N \rightarrow P]$

Findax: AXIOM FORALL (bk:B), (nm:N):

FindPhone(bk, nm) = bk(nm)

- ▶ add a phone number

AddPhone: $[B, N, P \rightarrow B]$

Addax: AXIOM FORALL (bk:B), (nm:N), (pn:P):

AddPhone(bk, nm, pn) = bk WITH [(nm) := pn]



Challenge

- ▶ testing vs. challenging

FindAdd: CONJECTURE

FORALL($bk:B$), ($nm:N$), ($pn:P$):

FindPhone(AddPhone(bk, nm, pn), nm) = pn

- ▶ Prove!



Short Reminder

phone-1 : THEORY

N: TYPE % names

P: TYPE % phone number

B: TYPE = [N -> P] % phone book

n0: P

emptybook: B

emptyax: AXIOM FORALL (nm:N):emptybook(nm)= n0

FindPhone: [B,N -> P]

Findax: AXIOM FORALL (bk:B), (nm:N): FindPhone(bk,nm) = bk(nm)

AddPhone: [B,N,P -> B]

Addax: AXIOM FORALL (bk:B), (nm:N), (pn:P):

 AddPhone(bk,nm,pn) = bk WITH [(nm):=pn]

FindAdd: CONJECTURE

FORALL(bk:B), (nm:N), (pn:P): FindPhone(AddPhone(bk,nm,pn),nm) = pn



Prove FindAdd Conjecture

- ▶ start prover

FindAdd :

|-----

{1} FORALL (bk:B) , (nm:N) , (pn:P) :

FindPhone (AddPhone (bk , nm , pn) , nm) = pn

Rule?

- ▶ type

`(grind :theories("phone-1"))`



Grind?

- ▶ ~ 20 basic commands &
- ▶ ~ 20 higher-level commands (strategies)
 - ▶ `assert`, `inst?`, `induct-and-simplify`, `induct-and-rewrite`, `skosimp*`
- ▶ highest level: `grind [args]`
 - ▶ skolemization
 - ▶ heuristic instantiation
 - ▶ propositional simplification (BDD)
 - ▶ if-lifting
 - ▶ rewriting
 - ▶ decision procedure for linear arithmetic & equality



Output

Rule? (grind :theories(phone-1))

Addax rewrites **AddPhone**(bk, nm, pn)
to bk WITH [(nm) := pn]

Findax rewrites

FindPhone(bk WITH [(nm) := pn], nm)
to pn

...

Q.E.D. ☺



DelPhone

DelPhone: $[B, N \rightarrow B]$

Delax: AXIOM FORALL (bk:B), (nm:N):

DelPhone(bk, nm) = bk WITH [(nm) := n0]

► challenge

DelAdd: CONJECTURE

FORALL (bk:B), (nm:N), (pn:P)

DelPhone(AddPhone(bk, nm, pn), nm) = bk



Prove DelAdd Conjecture

|-----

FORALL (bk:B), (nm:N), (pn:P):

DelPhone(AddPhone(bk,nm,pn),nm) = bk

Rule? (grind :theories(phone-1))

...

|-----

{1} bk!1 WITH[(nm!1) := pn!1]

WITH[(nm!1) := n0] = bk!1

Rule?

▶ **bk!1** ... representatives for quantified variables



Prove DelAdd Conjecture

|-----

{1} bk!1 WITH[(nm!1) := pn!1]
WITH[(nm!1) := n0] = bk!1

Rule? (apply-extensionality)

▶ to prove that $bk!1 = bk!1$ (functions are the same)

...

|-----

{1} bk!1 WITH[(nm!1) := pn!1]
WITH[(nm!1) := n0](x!1)
= bk!1(x!1)



Prove DelAdd Conjecture

...

Rule? (lift-if)

...

Rule? (ground)

{-1} nm!1 = (x!1)

|-----

{1} n0 = bk!1(x!1)

- ▶ We have to show, that in the original phone book, the phone number was n0.



Observation

- ▶ PVS is designed to be the “rigorous skeptic”
 - ▶ detect errors in reasoning (or specification?), but also
 - ▶ helps you to find the error by providing feedback



(Subset) Types

- ▶ Base types: `bool, int, real`
- ▶ Function types: `[bool, int -> int]`
- ▶ Subset types: `nat: TYPE = {i:int | i >= 0}`
- ▶ `/ : [int, { n: int | n /= 0 } -> int]`
- ▶ `average = sum / numbers : int`
- ▶ only well-typed if `numbers` is not 0
- ▶ type checking conditions (TCCs)



Type Checking

- ▶ type checking conditions (TCCs)
- ▶ type checking in PVS is undecidable
- ▶ proof obligations
- ▶ consistency check

- ▶ most obligations can be discharged automatically



Recursion

- ▶ must be shown to terminate

```
factorial(x:nat): RECURSIVE nat =  
  IF x = 0 THEN 1  
  ELSE x * factorial (x-1) ENDIF  
MEASURE (LAMBDA (x:nat):x)
```

- ▶ generates proof obligation

```
Factorial_TCC2: OBLIGATION  
(FORALL (x:nat): NOT x = 0 IMPLIES x-1 < x)
```



Abstract Data Types

- ▶ automatic generation of complete axiomatization

```
stack [t: TYPE]: DATATYPE
```

```
BEGIN
```

```
    empty: emptystack?
```

```
    push(top: t, pop: stack) : nonemptystack?
```

```
END stack
```

```
empty, push: constructors,
```

```
top, pop: accessors,
```

```
emptystack?, nonemptystack: recognizers
```



Abstract Data Types

- ▶ automatic generation of complete axiomatization:
- ▶ extensionality axioms for constructors
- ▶ eta axiom
- ▶ accessor/constructor axioms
- ▶ induction scheme
- ▶ functions distributing predicates over the stack base type
- ▶ subterm function
- ▶ well-foundedness axiom
- ▶ recursive combinator



Dependent Types

```
date: TYPE = [ yr: year, mon: month,  
  {d: nat | d <= days(mon, yr)} ]
```

- ▶ for function arguments

```
ratio(x, y: real, z: {z: real | z /= x}:real  
  = (x - y) / (x - z)
```

- ▶ generates TCCs



Propositional Proof

|-----
{1} P => Q

Rule: `flatten`

[-1] P
|-----
{1} Q

|-----
{1} P OR Q

|-----
{1} P
{2} Q

[-1] P & Q
|-----

[-1] P
[-2] Q
|-----



Propositional Proof

[-1] P OR Q

|-----

Rule: `split`

[-1] P

|-----

[-1] Q

|-----



Comparison

[17P]	HOL	Mizar	PVS
small proof kernel	+	-	-
extensible/programmable	+	-	+
powerful automation	+	-	+
readable	-	+	-
constructive logic	-	-	-
decidable types	+	+	-
dependent types	-	+	+
based on HOL	+	-	+
large mathematical std. lib.	+	+	+



My Impressions

- ▶ overwhelming
- ▶ good documentation available
 - ▶ especially for using the system
- ▶ on-going development since <1993

Thank you for your attention!



Bibliography

▶ [LFM]: NASA Langley Formal Methods Site

▶ <http://shemesh.larc.nasa.gov/fm/index.html>

[PVS]: PVS Specification and Verification System

▶ <http://pvs.csl.sri.com/index.shtml>

[TUT]: A Tutorial Introduction to PVS, Boca Raton, '95

▶ available at [PVS]

[17P]: The Seventeen Provers of the World,
Freek Wiedijk, '06

