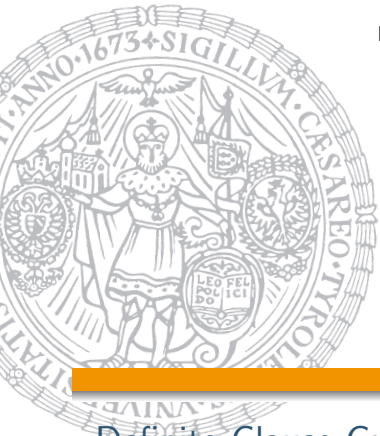


Logic Programming

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Definite Clause Grammars (DCGs for short)

Example

```

sentence(sentence(NP,VP),Num) →
    noun_phrase(N,Num), verb_phrase(V,Num).
:
determiner(det(the),Num) → [the].
determiner(det(a),singular) → [a].
noun(noun(pie-plate),singular) → [pie-plate].
noun(noun(pie-plates),plural) → [pie-plates].
noun(noun(surprise),singular) → [surprise].
noun(noun(surprises),plural) → [surprises].
adjective(adj(decorated)) → [decorated].
verb(verb(contains),singular) → [contains].
verb(verb(contains),plural) → [contain].
sentence(PT) ⇒* ‘‘the decorated pie-plates contain a surprise’’

```

Summary of Last Lecture

Observation

given a list $[1,2,3]$ it can be **represented** as the **difference** of two lists

- 1 $[1,2,3] = [1,2,3] \setminus []$
- 2 $[1,2,3] = [1,2,3,4,5] \setminus [4,5]$
- 3 $[1,2,3,8] = [1,2,3,8] \setminus [8]$
- 4 $[1,2,3] = [1,2,3|Xs] \setminus Xs$

Definition

the difference of two lists is denoted as $As \setminus Bs$ and called **difference list**

Example

```
append_dl(Xs \ Ys, Ys \ Zs, Xs \ Zs).
```

Outline of the Lecture

Logic Programs

introduction, basic constructs, database and recursive programming, theory of logic programs

The Prolog Language

programming in pure prolog, arithmetic, structure inspection, meta-logical predicates, cuts, extra-logical predicates, how to program efficiently

Advanced Prolog Programming Techniques

nondeterministic programming, incomplete data structures, **definite clause grammars**, **meta-programming**, constraint logic programming

Example (Facts)

```

father(andreas,boris).    female(doris).    male(andreas).
father(andreas,christian). female(eva).    male(boris).
father(andreas,doris).    male(christian).
father(boris,eva).    mother(doris,franz). male(franz).
father(franz,georg).    mother(eva,georg). male(georg).

```

Example

```

children(X,Cs) :- children(X,[],Cs).
children(X,A,Cs) :-
    father(X,C), \+ member(C,A), !, children(X,[C|A],Cs).
children(X,Cs,Cs).

```

Example

```

children(X,Kids) :- setof(C,father(X,C),Kids).

```

Second-Order Programming

Definitions

- the predicate *bagof*(*Template*,*Goal*,*Bag*) unifies *Bag* with the alternatives of *Template* that meet *Goal*
- if *Goal* has free variables besides the one sharing with *Template* *bagof* will backtrack
- fails if *Goal* has no solutions
- construct *Var* \uparrow *Goal* tells *bagof* to existentially quantify *Var*
- the predicate *setof*(*Template*,*Goal*,*Bag*) is similar to *bagof* but sorts the obtained multi-set (*bag*) and removed duplicates

Example

```

kids(Kids) :- setof(Y, X  $\uparrow$  (father(X,Y)),Kids).

```

Simple Application of Set Predicates

Example

```

no_doubles(Xs,Ys) :- setof(X,member(X,Xs),Ys).

```

Definition

- the predicate *findall*(*Template*,*Goal*,*Bag*) works as *bagof* if all excessive variables are existentially quantified
- SWI-Prolog notation for \uparrow : \wedge

Meta-Programming and Meta-Interpreters

Definition

- a *meta-program* treats other programs as data; it analyses, transforms, and simulates other programs
- a *meta-interpreter* for a language is an interpreter for the language written in the language itself
- relation *solve*(*Goal*) is true, if *Goal* is true with respect to the program interpreted

Example (simple meta-interpreter)

```

solve(true).
solve((A,B)) :- solve(A), solve(B).
solve(A) :- clause(A,B), solve(B).

```

Example (Simple Meta-Program)

```

eval(X,Y) :-
    number(X), X = Y.
eval(X,Y) :-
    nonvar(X), functor(X,F,N),
    built_in(F,N),
    functor(Z,F,N),
    eval_args(N,X,Z),
    Y is Z.
eval(X,Y) :-
    nonvar(X), functor(X,F,N),
    user_def(F,N),
    N1 is N + 1, functor(Z,F,N1),
    eval_args(N,X,Z),
    Z,
    arg(N1,Z,Y).

```

Example (Yet Another Simple Meta-Program)

```

edit :- edit(file([],[])).
edit(File) :- read(Command), edit(File,Command).

edit(_File,exit) :- !.
edit(File,Command) :-
    apply(Command,File,File1),
    !,
    edit(File1).
edit(File,Command) :-
    write(Command),
    write(' is not applicable '),
    !,
    edit(File).

apply(up,file([X|Xs],Ys),file(Xs,[X|Ys])).
apply(up(N),file(Xs,Ys),file(Xs1,Ys1)) :-
    N > 0,
    up(N,Xs,Ys,Xs1,Ys1).

```

Example (And the Last Example)

```

accept(S) :-
    initial(Q),
    accept(Q,S).
accept(Q,[X|Xs]) :-
    delta(Q,X,Q1),
    accept(Q1,Xs).
accept(Q,[]) :-
    final(Q).

initial(q0).
final(q2).
delta(q0,0,q0).
delta(q0,0,q1).
delta(q0,1,q0).
delta(q1,1,q2).

```

Example (meta-interpreter with proofs)

```

solve(true,true).
solve((A,B),(ProofA,ProofB)) :-
    solve(A,ProofA),
    solve(B,ProofB).
solve(A,(A :- Proof)) :-
    clause(A,B),
    solve(B,Proof).

```

Example

```

father(andreas,boris).      female(doris).           male(andreas).
father(andreas,christian).  female(eva).             male(boris).
father(andreas,doris).     male(christian).
father(boris,eva).         mother(doris,franz).     male(franz).
father(franz,georg).       mother(eva,georg).      male(georg).
son(X,Y) :- father(Y,X), male(X).

```

Example (A Meta-Interpreter with Proofs (cont'd))

```
:- solve(son(christian, andreas), Proof).
Proof  $\mapsto$  (son(christian, andreas) <--
  (father(andreas, christian) <-- true,
  male(christian) <-- true))
```

Example (Tracing Pure Prolog)

```
trace(Goal) :- trace(Goal, 0).
trace(true, Depth).
trace((A, B), Depth) :-
  trace(A, Depth), trace(B, Depth).
trace(A, Depth) :-
  clause(A, B),
  display(A, Depth),
  Depth1 is Depth + 1,
  trace(B, Depth1).
display(A, Depth) :- tab(Depth), write(A), nl.
```

Example

```
system(A is B).      system(read(X)).      system(integer(X)).
system(clause(A,B)). system(A < B).      system(A >= B).
system(write(X)).    system(funcator(T,F,N)). system(system(X)).
```

Example

```
trace(Goal) :- trace(Goal, 0).
trace(true, Depth) :- !.
trace((A, B), Depth) :- !, trace(A, Depth), trace(B, Depth).
trace(A, Depth) :- system(A), A, !, display2(A, Depth), nl.
trace(A, Depth) :-
  clause(A, B), display2(A, Depth), nl,
  Depth1 is Depth + 1, trace(B, Depth1).
trace(A, Depth) :-
  \+ clause(A, B), display2(A, Depth),
  tab(8), write(f), nl, fail.
display2(A, Depth) :- Spacing is 3*Depth, tab(Spacing), write(A).
```

Meta-Interpreters for Debugging

Example (Control Execution)

```
solve(true, _D, no_overflow) :-
  !.
solve(_A, 0, overflow([])).
solve((A, B), D, Overflow) :-
  D > 0,
  solve(A, D, OverflowA),
  solve_conjunction(OverflowA, B, D, Overflow).
solve(A, D, no_overflow) :-
  D > 0,
  system(A), !, A.
solve(A, D, Overflow) :-
  D > 0,
  clause(A, B),
  D1 is D - 1,
  solve(B, D1, OverflowB),
  return_overflow(OverflowB, A, Overflow).
```

Example (Control Execution (cont'd))

```
solve_conjunction(overflow(S), _B, _D, overflow(S)).
solve_conjunction(no_overflow, B, D, Overflow) :-
  solve(B, D, Overflow).

return_overflow(no_overflow, _A, no_overflow).
return_overflow(overflow(S), A, overflow([A|S])).

isort([X|Xs], Ys) :- isort(Xs, Zs), my_insert(X, Zs, Ys).
isort([], []).

my_insert(X, [Y|Ys], [X, Y|Ys]) :-
  X < Y.
my_insert(X, [Y|Ys], [Y|Zs]) :-
  X >= Y,
  my_insert(Y, [X|Ys], Zs).
my_insert(X, [], [X]).
```