# Logic Programming

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Outline of the Lecture

## Logic Programs

introduction, basic constructs, database and recursive programming, theory of logic programs

## The Prolog Language

programming in pure prolog, arithmetic, structure inspection, meta-logical predicates, cuts, extra-logical predicates, how to program efficiently

# Advanced Prolog Programming Techniques

nondeterministic programming, incomplete data structures, definite clause grammars, meta-programming, constraint logic programming, answer set programming

Summar

# Expert System with Certification Factor

```
solve4/1

solve4(true,1,_Treshold):-
!.

solve4((A,B),C,Treshold):-
!,
    solve4(A,C1,Treshold),
    solve4(B,C2,Treshold),
    minimum(C1,C2,C).

solve4(A,C,Treshold):-
    clause_cf(A,B,C1),
    C1 > Treshold,
    Treshold is Treshold / C1,
    solve4(B,C2,Treshold1),
    C is C1 * C2.
```

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An Abstract Problem

# Cryptarithmetic

#### Definition

- a cryptarithmetic problem is a puzzle in which each letter represents a unique digit  $\,\leqslant\,9$
- the object is to find the value of each letter
- first digit cannot be 0

# Example

#### An Abstract Proble

# First Attempt

# generate and test

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#### Constraint Logic Programming

# Constraint Logic Programming

### Definitions (CLP on finite domains)

- use\_module(library(clpfd)) loads the clpfd library
- Xs ins N .. M specifies that all values in Xs must be in the given range
- all\_different(Xs) specifies that all values in Xs are differnt
- label(Xs) all variables in Xs are evaluated to become values
- #=, #\=, #>, ...like the arithmetic comparison operators, but may contain (constraint) variables

### standard approach

- load the library
- specify all constraints
- call label to start efficient computation of solutions

### Discussion

```
very inefficient
```

```
?— time(solve(X)). % 133,247,057 inferences, % 7.635 CPU in 7.667 seconds (100% CPU, 17452690 Lips) X = [[9, 5, 6, 7], [1, 0, 8, 5], [1, 0, 6, 5, 2]]
```

# explanation

- generate-and-test in it's purest form
- all guesses are performed before the constraints are checked
- arithmetic checks cannot deal with variables

#### improvement

- move testing into generating
- destroys clean structure of program
- any other ideas?

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#### Constraint Logic Programming

# Second Attempt

## constraint logic program

# 8 queens (as before)

```
\begin{array}{lll} \text{queens}(\mathsf{Xs}) := & \operatorname{template}(\mathsf{Xs}), \operatorname{solution}(\mathsf{Xs}). \\ \\ \text{template}([1/\_Y1,2/\_Y2,3/\_Y3,4/\_Y4,\\ & & 5/\_Y5,6/\_Y6,7/\_Y7,8/\_Y8]). \\ \\ \text{solution}([]). \\ \text{solution}([\mathsf{X/Y}|\operatorname{Others}]) := \\ & & \operatorname{solution}(\operatorname{Others}),\\ & & \operatorname{member}(\mathsf{Y}, [1,2,3,4,5,6,7,8]),\\ & & \operatorname{noattack}(\mathsf{X/Y}, \operatorname{Others}). \\ \\ \\ \text{noattack}(\mathsf{X/Y}, [\mathsf{X1/Y1}|\operatorname{Others}]) := \\ & & Y = \setminus = Y1,\\ & & Y1 - Y = \setminus = X1 - X,\\ & & Y1 - Y = \setminus = X - X1,\\ & & \operatorname{noattack}(\mathsf{X/Y}, \operatorname{Others}). \\ \\ \end{array}
```

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Constraint Logic Programming

Sudoku

### Definition

- Sudoku is a well-known logic puzzle; usually played on a  $9 \times 9$  grid
- $\forall$  cells: cells  $\in \{1, \dots, 9\}$
- ∀ *rows*: all entries are different
- ∀ colums: all entries are different
- ∀ blocks: all entries are different

# Main Loop (using clp)

```
sudoku(Puzzle) :-
    show(Puzzle),
    flatten(Puzzle, Cells),
    Cells ins 1 .. 9,
    rows(Puzzle),
    cols(Puzzle),
    blocks(Puzzle),
    label(Cells),
    show(Puzzle).
```

```
n-queens (using clp)
```

```
nqueens(N,Qs) :-
    length(Qs,N),
    Qs ins 1 .. N, all_different(Qs),
    constraint_queens(Qs),
    label(Qs).

constraint_queens([]).
constraint_queens([Q|Qs]) :-
    noattack(Q,Qs,1),
    constraint_queens(Qs).

noattack(_,[],_).
noattack(X,[Q|Qs],N) :-
    X #\= Q+N,
    X #\= Q-N,
    M is N+1,
    noattack(X,Qs,M).
```

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Constraint Logic Programming

Sudoki

## auxiliary predicates

- flatten/2 flattens a list
- show/1 prints the current puzzle

### row/1

```
 \begin{array}{c} \mathsf{rows}\,(\,[\,]\,)\,\,.\\ \mathsf{rows}\,(\,[\,\mathsf{R}\,|\,\mathsf{Rs}\,]\,)\,\,:-\\ & \mathsf{all}_-\mathsf{different}\,(\,\mathsf{R}\,)\,,\;\;\mathsf{rows}\,(\,\mathsf{Rs}\,)\,. \end{array}
```

# row/1 (alternative)

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```
rows(Rs) := maplist(all_distinct, Rs).
```

cols/1 (alternative)

use *maplist*/2

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Sudol

# Demo

# blocks/1

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## Example

```
:- sudoku([[1,_,_,,,,,,,,,],

[-,_2,7,4,_,,,],

[-,_,5,_,,,,,,],

[-,3,_,,,,,,,],

[7,5,_,,,,,,,],

[-,-,,,,,],

[-,4,_,,,6,,,,],

[-,4,,,,,6,,,,],

[-,-,,,,,7,1],

[-,-,,,,,,1,,3,,]]).
```

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Answer Set Programming

# The New Kid on the Block

### Answer Set Programming

- novel approach to modeling and solving search and optimisation problems
- ¬ programming, but a specification language
- ¬ Turing complete
- purely declarative
- restricted to finite models

### Success Stories

- team building for cargo at Gioia Tauro Seaport
- expert system in space shuttle
- natural language processing
- . . .

# **Propositional Setting**

#### **Definitions**

- atoms, facts, rules are defined as before
- only constants (= propositions) are allowed as atoms
- negation is negation as failure
- disjunctions may appear in the head
- an answer set is a set of atoms corresponding to the minimal model of the program

```
Example (Negation as Failure)

light_on :- power_on, not broken.
power_on.
```

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answer set: { power\_on, light\_on}

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#### Answer Set Programming

#### Definition

constraints are negative assertions, representing fact that must not occur in any model of the program

# Example

```
a :- not a, b.
```

any answer set must not contain b and simplifies to

:- b.

#### Additional Features

- finite choice functions: {fact<sub>1</sub>, fact<sub>2</sub>, fact<sub>3</sub>}.
- choice and counting: 1{fact<sub>1</sub>, fact<sub>2</sub>, fact<sub>3</sub>}2. "1" or "2" may be missing

```
Example (Disjunctive Heads)

open | closed :- door.

answer sets: {open}, {closed}
```

# Example

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#### Answer Set Programming

# First-Order Setting

### Definition

- extension of first-order language
- no function symbols

# Example (3-colouring)

```
 \begin{array}{llll} \text{red} \, (X) & | & \text{green} \, (X) & | & \text{blue} \, (X) \, . \\ :- & \text{red} \, (X) \, , & \text{red} \, (Y) \, , & \text{edge} \, (X,Y) \, . \\ :- & \text{green} \, (X) \, , & \text{green} \, (Y) \, , & \text{edge} \, (X,Y) \, . \\ :- & \text{blue} \, (X) \, , & \text{blue} \, (Y) \, , & \text{edge} \, (X,Y) \, . \end{array}
```

```
Example ((part of) 8-queens problem)
:- not (1 = count(Y : queen(X,Y))), row(X)
expresses that exactly one queen appears in every row and column
```

## Grounders and Solvers



#### Grounders

- DLV (DLV Systems, Calabria)
- Gringo (University of Potsdam)
- Iparse (University of Helsinki)

#### Solvers

- clasp (University of Potsdam)
- cmodels (University of Austin)
- smodels (University of Helsinki)

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#### Answer Set Programming

```
Example
```

```
disk (1..n).
                                   peg(a;b;c).
transition (0...pathlength -1).
                                   situation (0.. pathlength).
location(Peg) :- peg(Peg).
                                   location(Disk) :- disk(Disk).
\#domain disk(X;Y).
                        #domain peg(P;P1;P2).
#domain transition(T). #domain situation(I).
#domain location(L;L1).
on(X,L,T+1) := on(X,L,T), not otherloc(X,L,T+1).
otherloc(X,L,I): - on(X,L1,I), L1!=L.
:- on(X,L,I), on(X,L1,I), L!=L1.
inpeg(X,P,I) := on(X,L,I), inpeg(L,P,I). inpeg(P,P,I).
top(P,L,I) := inpeg(L,P,I), not covered(L,I).
covered(L, I) := on(X, L, I).
:- on(X,Y,I), X>Y.
on(X,L,T+1) := move(P1,P2,T), top(P1,X,T), top(P2,L,T).
:- \text{ move}(P1, P2, T), \text{ top}(P1, P1, T). \text{ movement}(P1, P2) :- P1 != P2.
1 \{ move(A,B,T) : movement(A,B) \} 1.
on(n,a,0).
                                   on (X, X+1, 0) :- X < n.
onewrong: - not inpeg(X,c,pathlength).
:- onewrong.
```

Answer Set Programm

# Prolog and Answer Set Programming

- proof search
- Turing complete
- control
- efficiency

- model search
- finite domain
- specification language
- generality

### Example

```
\begin{split} & \text{hanoi} \left(0\,,_{-\,,\,-\,,\,-\,,\,-}\,,\,[\,]\,\right)\,, \\ & \text{hanoi} \left(N,X,Y,Z\,,Ls\,\right)\,:-\\ & N > 0\,,\,\,M\,\,\text{is}\,\,N\,-\,1\,,\\ & \text{hanoi} \left(M,X,Z,Y\,,Ls0\,\right)\,,\\ & \text{append} \left(Ls0\,,[\,\text{move}(N,X,Z)\,]\,,Ls1\,\right)\,,\\ & \text{hanoi} \left(M,Y,X,Z\,,Ls2\,\right)\,,\\ & \text{append} \left(Ls1\,,Ls2\,,Ls\,\right)\,. \end{split}
```

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Thank You for Your Attention!