



- 1 Consider the formula

$$(\bar{1} \vee 3) \wedge (1 \vee 2) \wedge (2 \vee \bar{4}) \wedge (4 \vee 1) \wedge (\bar{3} \vee 2) \wedge (4 \vee \bar{2} \vee \bar{3})$$

Give a DPLL inference sequence to determine its satisfiability. Indicate which rule was applied in each step.

- 2 Consider the formula

$$(\bar{7} \vee 6) \wedge (\bar{6} \vee 5) \wedge (\bar{6} \vee \bar{5} \vee \bar{4}) \wedge (\bar{2} \vee 1) \wedge (\bar{8} \vee \bar{7}) \wedge (\bar{5} \vee 3) \wedge (\bar{3} \vee \bar{2} \vee \bar{1} \vee 4)$$

and suppose a DPLL inference sequence reached the state $7^d 6 5 \bar{4} 3 \bar{8} 2^d 1$. Construct a conflict graph and indicate three different cuts together with the induced backjump clauses. Indicate at least one UIP.

- 3 Find an MUC of the following formula:

$$(x \vee y \vee z) \wedge \neg x \wedge (x \vee y) \wedge (x \vee \neg y \vee z) \wedge (x \vee z) \wedge (\neg x \vee \neg y \vee z) \wedge \neg z$$

- 4 Consider the following formula ϕ_3 :

$$(f(a) \approx b \vee a \approx b) \wedge (f(b) \approx b \vee a \approx c) \wedge f(a) \not\approx f(b) \wedge (f(a) \not\approx f(c) \vee a \not\approx b)$$

- (a) Construct a propositional skeleton of ϕ_3 .
 (b) Check satisfiability of ϕ_3 using DPLL(T).

- 5 Consider the following set of equations E :

$$\begin{array}{lll} a \approx b & f(a) \approx a & f(b) \approx c \\ g(b) \approx g(g(a)) & f(a) \approx g(a) & f(c) \approx f(g(c)) \end{array}$$

and use congruence closure to determine whether the following hold:

- (a) $E \models f(b) \approx c$ (b) $E \models f(c) \approx c$

- 6 Suppose a run of the Simplex algorithm reached the following intermediate state:

$$\begin{array}{l} x \quad s_3 \\ s_1 \left(\begin{array}{cc} -1 & -1 \\ -4 & -1 \\ -1 & -1 \\ 3 & 1 \end{array} \right) \begin{array}{l} s_1 \leq 4 \\ s_2 \leq -1 \\ s_3 \leq -5 \\ s_4 \leq 3 \end{array} \end{array}$$

with assignment $s_3 = -5$, $x = 0$, $s_1 = 15$, $s_2 = 5$, $y = 5$, $s_4 = -5$.

- (a) Which variable pairs are suitable for a pivot step, and why?
- (b) Perform one possible pivot step.

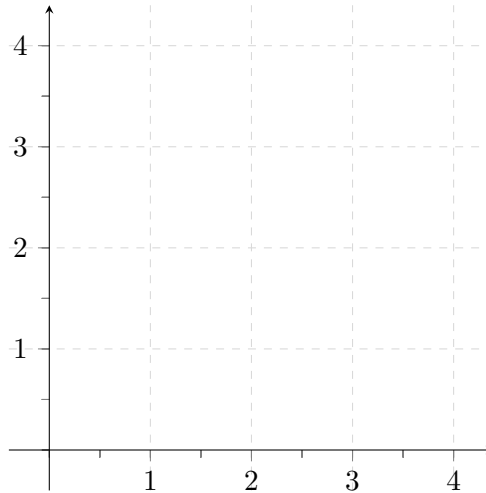
7 Consider the following system of linear inequalities:

$$2y - x \leq 2$$

$$2y + x \geq 2$$

$$y + 2x \leq 4$$

Draw the constraints in the following diagram. How many solutions are there in \mathbb{Z}^2 ?



8 Use equality graphs to determine whether the following formula is satisfiable.

$$x_1 \approx x_2 \wedge x_2 \approx x_3 \wedge x_3 \approx x_4 \wedge x_4 \not\approx x_6 \wedge x_1 \not\approx x_5 \wedge x_6 \not\approx x_1 \wedge x_4 \approx x_9 \wedge$$

$$x_4 \approx x_1 \wedge x_1 \not\approx x_9 \wedge x_6 \approx x_7 \wedge x_7 \approx x_8 \wedge x_5 \not\approx x_8 \wedge x_2 \approx x_9$$

9 Consider three bitvector variables \mathbf{a}_3 , \mathbf{b}_3 , and \mathbf{c}_3 of three bits each.

- (a) Give a propositional formula which encodes $\mathbf{a}_3 + \mathbf{b}_3 = \mathbf{c}_3$.
- (b) Give a propositional formula which encodes that the addition $\mathbf{a}_3 + \mathbf{b}_3$ overflows.

10 Use the Nelson-Oppen procedure to determine satisfiability of the following formula combining uninterpreted functions and linear real arithmetic. You can use either the deterministic or the nondeterministic version.

$$z = 0 \wedge x = y + z \wedge f(y) = 1 \wedge f(x) = 2$$

11 Consider the following C program P :

```

1 int main() {
2   int i = 2;
3   int m = 1;
4   while (i <= 10) {
5     m = m * 2;
6   }
7 }
```

- (a) Draw the program graph.
- (b) Describe an encoding of a state of P .
- (c) Define a predicate that captures the transition relation of P (for your encoding of a state).