



# Interactive Theorem Proving using Isabelle/HOL

Session 13

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## Topics

calculational reasoning, case analysis, code generation, computation induction, data type invariants, document preparation, finding theorems, first steps, functional programming in HOL, higher-order logic, history and motivation, induction, inductive definitions, Isabelle basics, Isabelle/Isar, Isabelle/ML, **IsaFoR/CeTA**, locales, manual termination proofs, multisets, natural deduction, notation, proof methods, PSL: a high-level proof strategy language, rule induction, rule inversion, session management, sets, simplification, sledgehammer, structural induction, structured proof, **The Archive of Formal Proofs, the certification approach**, total recursive functions, type classes, type definitions, well-foundedness

# Overview

- The Archive of Formal Proofs
- The Certification Approach
- The IsaFoR/CeTA Project
- Exercises

## The Archive of Formal Proofs

- collection of Isabelle/HOL formalizations
- available from <https://www.isa-afp.org/>
- maintained by Isabelle developers (that is, kept up to date with latest Isabelle version)
- submit your own development through web-interface

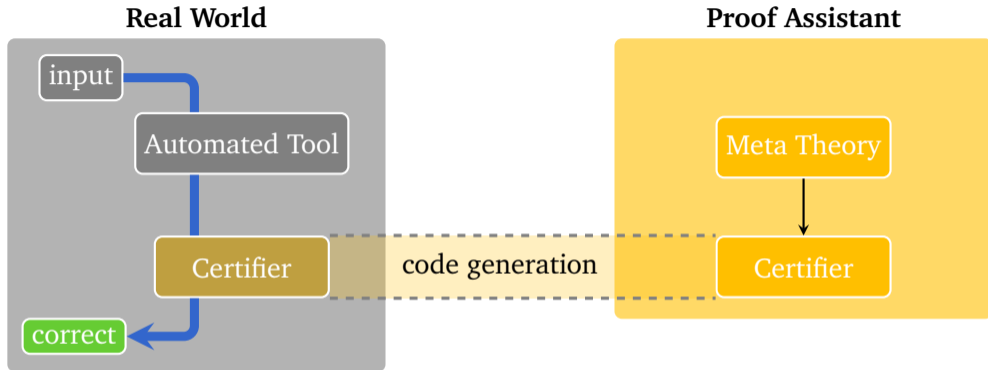
## Using the AFP

- (1) download the AFP tar-archive and extract it into \$AFP
- (2) `$ echo "$AFP/thys" >> ~/.isabelle/Isabelle2018/ROOTS`

## Demo13.thy – First-Order Terms

- apply substitution  $\sigma = \{y \mapsto g(b)\}$  to term  $t = f(x, y)$ :  $t\sigma = f(x, g(b))$
- unify terms  $f(a, x)$  and  $f(x, b)$ : ✗
- unify terms  $f(a, x)$  and  $f(y, b)$ : ✓ with mgu  $\{x \mapsto b, y \mapsto a\}$

# The Certification Approach



- implement untrusted automated tool
- formalize meta-theory of tool in proof assistant
- use formalized meta-theory to verify certifier for tool output
- certifier bridges real world and logical world of proof assistant

## IsaFoR/CeTA

- available from <http://cl-informatik.uibk.ac.at/isafor>
- specific instance of certification approach for term rewriting

### Example – Termination of Odd Even/Odd Function

- `fun evenodd :: "nat  $\Rightarrow$  bool  $\Rightarrow$  bool"`  
`where`  
`"evenodd 0 True  $\longleftrightarrow$  True"`  
`| "evenodd x False  $\longleftrightarrow$   $\neg$  (evenodd x True)"`  
`| "evenodd (Suc x) True  $\longleftrightarrow$  evenodd x False"`
- a corresponding term rewrite system (TRS)
- Tyrolean Termination Tool 2 (ttt2) proves termination automatically
- verified certifier CeTA validates resulting proof

# Exercises

work on your semester projects