Summer Term 2023

Interactive Theorem Proving using Isabelle/HOL

Session 1

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Outline

Organization

Course Info (VU 3)

- LV-Number: 703315
- instructor: René Thiemann
- VU: attendance mandatory, shared lecture and proseminar
- website: `http://cl-informatik.uibk.ac.at/teaching/ss23/itpIsa`
  (slides and Isabelle files are available online)
- consultation hours: Tuesday 10:00 – 11:00 in 3M09 (ICT building)

Grading

- weekly exercises (50 %)
- project (50 %)
  (finished projects must be submitted through OLAT
  deadline: August 1, 2023)

The Exercises

- weakly exercise will be handed out each Thursday
- mark and upload solved exercises in OLAT until Thursday, 6am
- solutions will be discussed at start of each VU

The Project

- list of potential formalization projects will be made available
- projects will be assigned on April 27
- work alone or in small groups (depending on specific project)
- projects have to be finished before August 1
- be able to answer project related questions

Course Information

- two courses on interactive theorem proving provided by CL
- VU3 Interactive Theorem Proving (Cezary Kaliszyk)
  - broader: different proof assistants based on different logics
  - covers foundations of interactive theorem provers
- VU3 Interactive Theorem Proving using Isabelle/HOL (this course)
  - focussed: single proof assistant (Isabelle), one logic (HOL: higher-order logic)
  - practical course to obtain hands-on experience

$\hookrightarrow$ good idea to attend both courses

Literature

- Isabelle documentation
  (`https://isabelle.in.tum.de`)
  - Tobias Nipkow: Programming and Proving in Isabelle/HOL
  - ...
- Tobias Nipkow and Gerwin Klein: Concrete Semantics with Isabelle/HOL
  (`http://www.concrete-semantics.org`)

Acknowledgement

- Several slides have been taken from a previous course on interactive theorem proving given by Christian Sternagel

Motivation and Introduction

Motivation

- bugs in unverified software and hardware may have severe consequences
- these can be costly (crash of Ariane, Pentium bug, ...)
- or fatal (control software of aircrafts, medical devices, ...)

One Solution: Formal Verification

Proving program correctness with respect to given formal specification

State of the Art in Formal Verification

- verified SAT solver wins against unverified SAT solvers in competition
- verified operating system kernel (seL4)
  (no arithmetic exceptions, deadlocks, buffer overflows, ...)
- verification of Kepler conjecture: optimal density of packing spheres is $\pi/\sqrt{18}$
- 99 % of a top 100 mathematical theorems list has been verified

  `https://www.cs.ru.nl/~freek/100/`

Formal Verification via Theorem Proving

- various logics to write formal specifications
  - propositional logic, SMT, first-order logic
  - higher-order logic (HOL), calculus of inductive constructions
- logics differ in expressivity and automation
  - automated theorem proving (ATP)
    - push button verification (SAT solver, SMT solver, first-order resolution prover, ...)
    - limited expressivity
  - interactive theorem proving (ITP)
    - proofs are developed manually (within a proof assistant)
    - less automation
    - high expressivity (mathematical theorems, program verification, ...)
- Isabelle is a popular proof assistant (besides Coq, Lean, PVS, ...) that supports HOL
- HOL is sweet spot between expressivity and automation

What is a Proof Assistant?

- combination of automated theorem prover (ATP) and proof checker
- structure of proofs is designed manually, some subproofs are found automatically
- all proofs are checked rigorously, e.g., in an LCF-style proof assistant such as Isabelle

Examples

- automatic methods
  - logical reasoning (e.g., linear arithmetic, first-order reasoning)
  - equational reasoning
  - ...
- manual steps
  - provide intermediate statements or auxiliary lemmas
  - perform induction or case analysis
  - ...
- proof checking
  - check that all cases have been covered, that inference rules are applied correctly, ...

What is LCF-Style?

- theorems are represented by abstract data type (`thm`)
- set of (basic) logical inferences provided as interface (trusted kernel)
- no other ways to create theorem (value of type `thm`) due to abstraction barrier and strong type system

Example

- kernel provides functions `assume : cterm -> thm`
  and `implies_elim : thm -> thm -> thm`
- implements inference rules

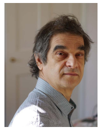
$$\frac{}{A \vdash A} \qquad \frac{\Gamma \vdash A \Longrightarrow B \quad \Delta \vdash A}{\Gamma, \Delta \vdash B}$$

- if desired, inspect implementation of kernel functions to increase trust

History of Isabelle

- 1986: creation of Isabelle, a proof assistant for various logics (University of Cambridge, Technische Universität München)
- 1993: support for higher-order logic: Isabelle/HOL
- 1996: human-readable proof language: Isabelle/Isar
- 2011: prover IDE: Isabelle/jEdit
- since 2004: archive of formal proofs (a library of formalized proofs with currently 447 authors and 231 100 lemmas)

Tobias Nipkow



Lawrence Paulson

Makarius Wenzel

Higher-Order Logic

Higher-Order Logic

- we assume knowledge of first-order logic
- higher-order logic has two main differences to first-order logic
  - terms are typed
  - quantification for each type, including function-types
- higher-order logic will be used both to specify functional programs as well as logical specifications

HOL = Functional Programming + Logic

Types in HOL

- very similar to Haskell types
  - basic types for booleans, natural numbers, integers, ...
  - type variables
  - function types
  - algebraic data types: lists, trees, pairs, tuples, ...
- in Isabelle
  - function types have form $input\_type \Rightarrow output\_type$
  - $ty_1 \Rightarrow ty_2 \Rightarrow ty_3$ is the same as $ty_1 \Rightarrow (ty_2 \Rightarrow ty_3)$
  - type variables are written with a leading prime: `'a`, `'b`, ...
  - most type constructors are written postfix: `'a list`, `nat list list`, ...
  - tuples are encoded as nested pairs: `'a × 'b × 'c` is the same as `'a × ('b × 'c)`
  - new algebraic data types can be created via `datatype` as in
    `datatype ('a,'b)tree = Leaf 'a | Node "('a,'b)tree" 'b "('a,'b)tree"`
  - type synonyms (abbreviations) can be created via `type_synonym` as in
    `type_synonym ('a)special_tree = "(nat × 'a, 'a list)tree"`
    `type_synonym string = "char list"`

Inner and Outer Syntax

- Isabelle contains various languages
  - implementation languages Scala and ML
  - language to write Isabelle theories: outer syntax
    - add a function definition
    - add a type definition
    - state a lemma
    - perform a proof step
    - …
  - language to specify terms and types: inner syntax
    - provide defining equations of a function
    - provide definition of type
    - provide a formula that describes the lemma
    - instantiate some inference rule, e.g., provide a term as existential witness
    - …
- important
  - content of inner syntax needs to be surrounded by double-quotes
  - exception: if content is atomic, then double-quotes can be dropped

Example for Outer and Inner Syntax: Data Type Definitions

- general definition is specified by outer syntax:
  $\texttt{datatype}\ ('a_1,\ \ldots,\ 'a_n)\ ty\ =$
  $\quad C_1\ ty_{11}\ \ldots\ ty_{1k_1}\ |\ \ldots\ |\ C_m\ ty_{m1}\ \ldots\ ty_{mk_m}$
- each ty_{ij} is a type, i.e., something that is specified by inner syntax
- consider concrete data type definition from previous slide
  `datatype ('a,'b)tree = Leaf 'a | Node "('a,'b)tree" 'b "('a,'b)tree"`
  - the first argument of `Node` is `"('a,'b)tree"` – double-quotes required
  - the second argument of `Node` is `'b` – double-quotes not required
  - further examples
    - both `nat` and `"nat"` are okay
    - `"nat ⇒ bool"`
    - `"nat list"`
    - `"(nat × 'a) list"`
  - once we are inside inner syntax, no further double-quotes are allowed:
    `"("nat × 'a") list"` is not permitted

Difference Between Types in Haskell and in HOL

- although HOL types look similar to Haskell types there are two import differences
- data type definitions in Isabelle/HOL do not include infinite applications of constructors
  - consider `datatype 'a list = Nil | Cons 'a "'a list"`
  - in Haskell, lists can be infinite, e.g., `ones = Cons 1 ones`
  - in Isabelle/HOL, only finite lists are covered by type `'a list`
- all types in HOL must be inhabited
  - for each `datatype` invocation, Isabelle internally checks that at least one term of the new type can be created, and if not, the new type is not accepted
  - example: `datatype foo = Bar foo` is refused

Terms in HOL

- terms in Isabelle/HOL are similar to Haskell terms, they include
  - literals: 0, 5, `''hello''`, `CHR ''c''`, …
  - variables: free `x`, `y`, `xs`, …or bound `x`, `y`, `xs`, …
  - constants: `True`, `False`, `Nil`, `Cons`, (∨), (∧), (¬), (⟶), (=), (<), (+), `map`, …
  - application: $t_1\ \ t_2$ – multiple arguments $t_1\ \ t_2\ \ t_3$ are encoded as $(t_1\ \ t_2)\ \ t_3$
  - λ-abstractions: λ `x`. t – multiple arguments λ `x y`. t are encoded as λ `x`. (λ `y`. t)
  - type constraints: t `::` ty
  - Isabelle/HOL provides further syntactic conveniences like if-then-else, let, case, infix-syntax, special syntax for lists and quantifiers, …
- terms are typed, Isabelle performs type inference and type checking
- HOL-formulas are just terms of type `bool`
- example terms
  - `map (λ x :: nat. x + 1) [1, 3]` is a term with type `nat list`
  - `map f (Cons x xs) = Cons (f x) (map f xs)` might be a defining equation of `map`
  - `(x :: nat) + (y + z) = (x + y) + z ∧ x + y = y + x :: bool`
    states that addition of natural numbers is associative and commutative

Quantors and Equality in HOL

- unlike in Haskell, equality is available for all types
- two consequences
  - equality is not necessarily executable
  - quantors are not primitive in HOL, but can be encoded
- example
  - define universal quantification as a function `All :: ('a ⇒ bool) ⇒ bool` via `definition "All P = (P = (λ x. True))"`
  - ∀-quantifier is nothing else than syntactic sugar, e.g. `∀ x. P x y` is syntax for `All (λ x. P x y)`
  - properties of universal quantifiers (introduction and elimination rules) can be derived ↪ we will work with these derived properties and ignore the internal definition
- facts
  - Isabelle/HOL contains only very few axiomatized types and constants (`bool` and some infinite type, `(⟶)`, `(=)` and `The, Eps :: ('a ⇒ bool) ⇒ 'a`)
  - all other types and constants are defined on top of these
  - we won't cover the details of these foundations in this course

Examples Beyond First Order Logic

```
(* well-foundedness of a binary relation can be expressed *)
type_synonym 'a rel = "'a ⇒ 'a ⇒ bool"
definition "well_founded (R :: 'a rel)
  = (¬ (∃ f :: nat ⇒ 'a . ∀ n :: nat. R (f n) (f (n + 1))))"

(* the transitive closure of a relation can be expressed *)
definition "trans_cl (R :: 'a rel) a b
  = (∃ (f :: nat ⇒ 'a) (n :: nat).
     f 0 = a ∧ f n = b ∧ n ≠ 0 ∧
     (∀ i. i < n ⟶ R (f i) (f (i + 1))))"

lemma "well_founded (trans_cl R) = well_founded R" oops

(* induction on natural numbers is sound *)
lemma "∀ P :: nat ⇒ bool.
  P 0 ⟶ (∀ n. P n ⟶ P (n + 1)) ⟶ (∀ n. P n)" oops
```


Color-Codes of Isabelle

- `keyword` a keyword of outer syntax
- `command` a command of outer syntax
- `const` a constant that has been defined before
- `free` a free variable
- `bound` a bound variable (of λ or quantifier)
- `fixed` a fixed variable (e.g., after ∀-introduction in a proof)
- colors help to identify mistakes, e.g. in
  `definition "select_first fst _ = fst"`
  the black color of `fst` indicates that `fst` is an already defined constant
  (and not a bound variable `fst`), so that a name clash needs to be resolved
- at the time of a definition, the used name is free (`name`),
  only afterwards it turns to black (`name`)
- `free` variables in lemmas are implicitly universally quantified
  (and can be instantiated after the lemma has been proven)

Functional Programming in HOL

- functional programs can be written similarly to Haskell
- already seen: type definitions
- new: function definitions
- non-recursive function (or constant) definitions
  - outer syntax: `definition` *name* `::` *ty* `where` *eqn* or just `definition` *eqn*
  - *eqn* is a boolean term of the shape
    name x_1 … x_n = *t*
  - important: often *t* needs to be put in parenthesis
    `definition "sorted_triple x y z = (x ≤ y ∧ y ≤ z)"`
- recursive function definitions
  - outer syntax: `fun` *name* `::` *ty* `where` eqn_1 | … | eqn_m
  - each eqn_i is a boolean term of the shape
    name pat_1 … pat_n = *t*
  - example

```
fun append :: "'a list ⇒ 'a list ⇒ 'a list" where
  "append Nil ys = ys"
| "append (Cons x xs) ys = Cons x (append xs ys)"
```

Function Definitions in Isabelle/HOL

- syntactic differences to Haskell
  - let-expressions are of form `let` x_1 = t_1; … ; x_n = t_n `in` t
  - case-expressions are of form `case` t `of` $pat_1 \Rightarrow t_1$ | … | $pat_n \Rightarrow t_n$
  - let, case, and if-then-else often have to be surrounded by parenthesis
  - let-expressions are sequential and non-recursive: t_i may not refer to $x_i, \ldots, x_n$
  - no local recursive function definitions
  - no restriction to executable functions:
    `fun f where "f P = (if (∀ x :: nat. P x) then 0 else 1)"`
- semantic difference to Haskell
  - functions defined by `fun` have to be terminating
  - if Isabelle is not able to prove termination, a function definition is not accepted

First Steps with Isabelle/HOL

Isabelle 2022 in RR 20

- from now on prefix "`$ `" indicates bash prompt
- start Isabelle via
  `$ isabelle jedit (optional: File.thy)`
- in case `isabelle` is not found, add `$ISABELLE_HOME/bin/` to your `PATH` where `ISABELLE_HOME` is `/usr/site/isabelle/2022/`

Isabelle 2022 on Your Machine

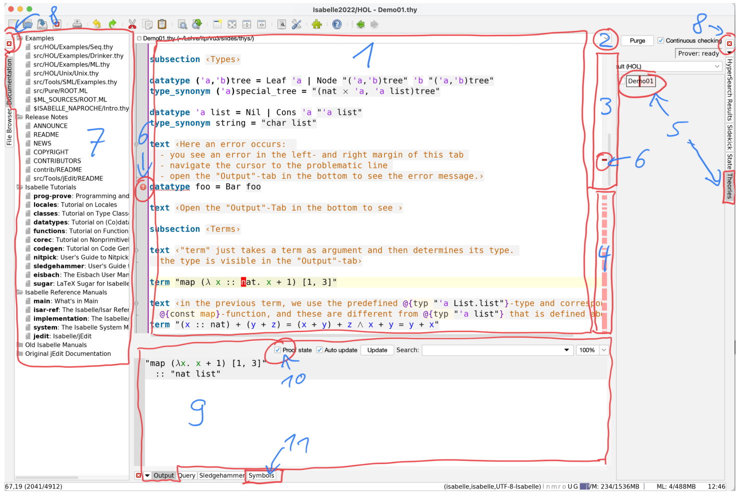
- download and installation instructions available at

  https://isabelle.in.tum.de

Demo

```
$ isabelle jedit Demo01.thy
```


Isabelle/jEdit – Overview of User Interface



Explanation of Previous Slide

1. main text area
2. switch between different theories
3. processed part of theory
4. unprocessed part of theory
5. progress indicator of several theories
6. indication of problem
7. documentation
8. click to close left or right panel
9. main output window
10. enable to view proof state in output (and not just errors)
11. symbol panel for information on special symbols

Theory Files – General Structure

```
theory T
  imports T_1 ... T_n
begin
(* definitions, theorems and proofs *)
...
end
```

Notes

- store theory T in file T.thy
- definitions and theorems from theories $T_1, \ldots, T_n$ available after begin
- new definitions, theorems and proofs go between begin and end
- qualify identifiers by theory name (like T.f) to disambiguate names
- theory Main is collection of basic definitions (like Haskell's Prelude) and should always be imported

Entering Special Symbols

- aim: enter symbols like ∀, ×, λ, ...
- four methods
  - switch to Symbols-panel in Isabelle/jEdit, find and click on symbol; important: hovering over symbol will reveal internal name and abbreviations
  - enter internal name prefixed by backslash and use auto-completion via $\boxed{\text{TAB}}$; example: $\boxed{\backslash}\,\boxed{f}\,\boxed{o}\,\boxed{r}$ will result in \<forall>, i.e., ∀
  - enter abbreviation followed by $\boxed{\text{TAB}}$, e.g., ∀ is also obtained via $\boxed{!}\,\boxed{\text{TAB}}$
  - some abbreviations have an auto completion where no $\boxed{\text{TAB}}$ is required, e.g., $\boxed{/}\,\boxed{\backslash}$ will immediately result in ∧

Frequently Used Symbols

symbol	internal	auto completion	abbreviations
λ	\<lambda>		$\boxed{\%}$
⇒	\<Rightarrow>	$\boxed{=}\,\boxed{>}$	$\boxed{.}\,\boxed{>}$
¬	\<not>		$\boxed{\sim}$
∧	\<and>	$\boxed{/}\,\boxed{\backslash}$	$\boxed{\&}$
∨	\<or>	$\boxed{\backslash}\,\boxed{/}$	$\boxed{\mid}$
⟶	\<longrightarrow>	$\boxed{-}\,\boxed{-}\,\boxed{>}$	$\boxed{.}\,\boxed{>}$
⟷	\<longleftrightarrow>	$\boxed{<}\,\boxed{-}\,\boxed{-}\,\boxed{>}$	$\boxed{<}\,\boxed{>}$
∀	\<forall>		$\boxed{!}$ and $\boxed{A}\,\boxed{L}\,\boxed{L}$
∃	\<exists>		$\boxed{?}$ and $\boxed{E}\,\boxed{X}$
×	\<times>	$\boxed{<}\,\boxed{*}\,\boxed{>}$	
≤	\<le>	$\boxed{<}\,\boxed{=}$	