

## Homework

1. Consider the formula

$$\varphi \equiv (x \neq v \vee v \neq y) \wedge (x = v \vee x \neq y) \wedge (x = y \vee x \neq z) \wedge z = w \wedge (x = z \vee w \neq y)$$

Using DPLL(T) check the satisfiability of  $\varphi$  for the following theories.

- (a)  $T_1$ : The equality logic over the natural numbers where the single predicate  $=$  is interpreted as the identity over  $\mathbb{N}$ . (2 P)
- (b)  $T_2$ : The equality logic over the Boolean values, where the single predicate  $=$  is the identity over  $\mathbb{B} = \{0, 1\}$ . (2 P)
2. Consider the following instance  $\gamma$  of the Chinese Remainder Theorem: for all  $a$  and  $b$  there is an  $x$  (2 P) such that  $x \equiv a \pmod{7}$  and  $x \equiv b \pmod{5}$ .
- (a) Can  $\gamma$  be expressed in Peano arithmetic? If yes, how?
- (b) Can  $\gamma$  be expressed in Presburger arithmetic? If yes, how?
3. Consider the following *Greater Than Killer Sudoku*. It follows the same rules as Killer Sudoku, but additionally adds equality ( $=$ ) and greater than ( $>$ ) constraints on the sum of some cages.

		3	12		19		14
#			23		8		
18	10			20	14		
					5		
						20	
14							
9		#					
4	15						
	11						7

- (a) Encode this puzzle in Presburger arithmetic. (You are allowed to use  $>$  as a predicate). (2 P)
- (b) Solve the puzzle using an SMT-solver (for example Z3). Is the solution unique? (2 P)

Hint: To encode that some variables  $x_1, \dots, x_n$  are distinct you may just write  $distinct(x_1, \dots, x_n)$ . SMT-LIB 2 also supports the constraint (`assert (distinct x1 x2 x3 ...)`).