



Interactive Theorem Proving Lecture 1 (VU)

Cezary Kaliszyk

Administration

Teacher

- Cezary Kaliszyk
- Consultation hours: Wednesday midday, 3M12, on demand

Grading

- Exercises: Assignments and participation
- Presentations
- Closed book test (last week)

Practical Assignments

- Software: HOL Light, Coq, Set Theory
- Laptops convenient

Content

- Proof Assistants
- Lambda calculus
 - types, Church vs Curry, derivation formats, well-typedness, term finding
- Second-order typed lambda calculus
 - Π -types, second-order abstraction and application, $\lambda 2$
- Types dependent on types
 - Sorts, weakening, formation, properties
- Dependent types
 - λ P, minimal logic, natural deduction again
- CoC
 - λ -cube, Girard's paradox, classical logic
- Definitions
 - terms, types, δ -conversion, \rightarrow_{Δ} , axioms
- Sets and set theory
- Numbers and arithmetic
 - \mathbb{N} , bits, efficient computation, divisibility, proof irrelevance

Outline

Today

- Proof assistants
- Common uses
- Comparison with other tools

Course Prerequisites

- Propositional and predicate logic
- Functional programming

Automated Reasoning

- Computer used to reason in a logic
- Traditionally part of artificial intelligence
 - (not machine learning)
- Field of research since the fifties
- Applications: program verification, mathematical deduction, ...
- Theorem proving logics, precision, automation, ... very varied.

What is a Proof Assistant? (1/2)

A Proof Assistant is a

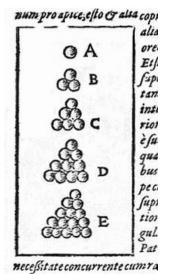
- a computer program
- to assist a mathematician
- in the production of a proof
- that is mechanically checked

What does a Proof Assistant do?

- Keep track of theories, definitions, assumptions
- Interaction proof editing
- Proof checking
- Automation proof search

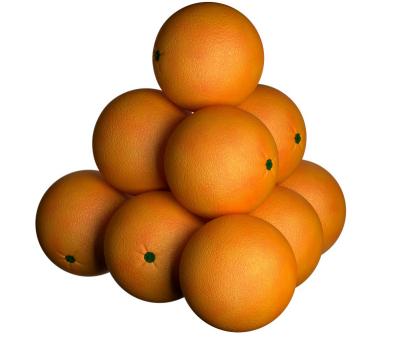
What does it implement? (And how?)

- a formal logical system intended as foundation for mathematics
- decision procedures



The most compact way of stacking balls of the same size in space is a pyramid.

$$V=rac{\pi}{\sqrt{18}}pprox 74\%$$



Proved in 1998

- Tom Hales, 300 page proof using computer programs
- Submitted to the Annals of Mathematics

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- 99% correct... but we cannot verify the programs

1039 equalities and inequalities

For example:

$$\frac{-x_1x_3 - x_2x_4 + x_1x_5 + x_3x_6 - x_5x_6 + \\ +x_2(-x_2 + x_1 + x_3 - x_4 + x_5 + x_6)}{4x_2 \begin{pmatrix} x_2x_4(-x_2 + x_1 + x_3 - x_4 + x_5 + x_6) + \\ +x_1x_5(x_2 - x_1 + x_3 + x_4 - x_5 + x_6) + \\ +x_3x_6(x_2 + x_1 - x_3 + x_4 + x_5 - x_6) - \\ -x_1x_3x_4 - x_2x_3x_5 - x_2x_1x_6 - x_4x_5x_6 \end{pmatrix}} < \tan(\frac{\pi}{2} - 0.74)$$

Solution? Formalized Proof!

- Formalize the proof using Proof Assistants
- Implement the computer code in the system
- Prove the code correct
- Run the programs inside the Proof Assistant

Flyspeck Project

- Completed 2017
- Many Proof Assistants and contributors

Intel Pentium P5 (1994)

FPU unit

- Division lookup table
- For certain inputs division result off

Replacement

- Few customers cared, still 450M\$
- Birth of HOL Light
- Intel and AMD processors formally verified

```
theorem sqrt2_not_rational:
    "sqrt (real 2) ∉ ℚ"
proof
```

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    assume "sqrt (real 2) ∈ ℚ"
```

thus False qed

```
theorem sqrt2_not_rational:
    "sqrt (real 2) ∉ Q"
proof
    assume "sqrt (real 2) ∈ Q"
    then obtain m n :: nat where
    n_nonzero: "n ≠ 0" and sqrt_rat: "¦sqrt (real 2)¦ = real m / real n"
    and lowest terms: "gcd m n = 1" ..
```

thus False qed

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    n_nonzero: "n ≠ 0" and sqrt_rat: "¦sqrt (real 2)¦ = real m / real n"
    and lowest terms: "gcd m n = 1" ..
```

```
have eq: m^2 = 2 * n^2"
hence "2 dvd m<sup>2</sup>" ...
have dvd m: "2 dvd m"
```

```
hence "2 dvd n<sup>2</sup>" ...
have "2 dvd n"
have "2 dvd gcd m n"
```

```
thus False
ged
```

```
theorem sgrt2 not rational:
  "sgrt (real 2) ∉ 0"
proof
  assume "sqrt (real 2) \in \mathbb{O}"
  then obtain m n :: nat where
    n nonzero: "n \neq 0" and sqrt rat: "!sqrt (real 2)! = real m / real n"
    and lowest terms: "qcd m n = 1"...
  from n nonzero and sort rat have "real m = {sort (real 2)} * real n" by simp
  then have "real (m^2) = (sqrt (real 2))^2 * real <math>(n^2)"
    by (auto simp add: power2 eq square)
  also have "(sqrt (real 2))<sup>2</sup> = real 2" by simp
  also have "... * real (m^2) = real (2 * n^2)" by simp
  finally have eq: m^2 = 2 * n^2"...
  hence "2 dvd m<sup>2</sup>"
  with two is prime have dvd m: "2 dvd m" by (rule prime dvd power two)
  then obtain k where "m = 2^* k" ...
  with eq have "2 * n^2 = 2^2 * k^2" by (auto simp add: power2 eq square mult ac)
  hence "n^2 = 2 * k^2" by simp
  hence "2 dvd n<sup>2</sup>" ...
  with two is prime have "2 dvd n" by (rule prime dvd power two)
  with dvd m have "2 dvd gcd m n" by (rule gcd greatest)
  with lowest terms have "2 dvd 1" by simp
  thus False \overline{b}y arith
aed
```

```
:: W The Irrationality of the Square Root of 2
theorem Th1:
 for p being Element of NAT st p is prime holds
 sort p is irrational
proof
 let p be Element of NAT ;
 assume A1: p is prime ;
 then A2: p > 1 by INT 2:def 4:
 assume sqrt p is rational ;
 then consider i being Integer, n being Element of NAT such that
 A3: n \iff 0 and
 A4: sort p = i / n and
 A5: for il being Integer
 for n1 being Element of NAT st n1 <> 0 & sart p = i1 / n1 holds
 n <= n1 by RAT 1:9:
 A6: i = (sqrt p) * n by A3, A4, xCMPLX 1:87;
 sqrt p \ge 0 by SQUARE 1:def 2;
 then reconsider m = i as Element of NAT by A6, INT 1:3:
 A7: m^2 = ((sart p)^2) * (n^2) bv A6
 .= p * (n ^2) by SQUARE_1:def 2 ;
 then p divides m ^2 by NAT D: def 3:
 then p divides m by A1, NEWTON:80;
 then consider m1 being Nat such that
 A8: m = p * m1 bv NAT D: def 3:
 n^{2} = (p^{*}(m1^{2})) / p^{2} hv^{2} A7, A8, x_{CMPLX 1:89}
 .= p * (m1 ^2) by A2, XCMPLX 1:89;
 then p divides n ^2 by NAT D:def 3;
 then p divides n by A1, NEWTON:80;
 then consider n1 being Nat such that
 A9: n = p * n1 bv \text{ NAT D: def 3:}
 A10: m1 / n1 = sqrt p by A2, A4, A8, A9, XCMPLX 1:91;
 A11: n1 <> 0 bv A3, A9:
 then p * n1 > 1 * n1 bv A2. XREAL 1:98:
 hence contradiction by A5, A9, A11, A10:
end;
```

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Proof Assistant (2/2)

- Keep track of theories, definitions, assumptions
 - set up a theory that describes mathematical concepts (or models a computer system)
 - express logical properties of the objects
- Interaction proof editing
 - typically interactive
 - specified theory and proofs can be edited
 - provides information about required proof obligations
 - allows further refinement of the proof
 - often manually providing a direction in which to proceed.
- Automation proof search
 - various strategies
 - decision procedures
- Proof checking
 - checking of complete proofs
 - sometimes providing certificates of correctness
- Why should we trust it?
 - small core

Can a Proof Assistant do all proofs?

Can a Proof Assistant do all proofs?

Decidability!

- Validity of formulas is undecidable
- (for non-trivial logical systems)

Automated Theorem Provers

- Specific domains
- Adjust your problem
- Answers: Valid (Theorem with proof)
- Or: Countersatisfiable (Possibly with counter-model)

Proof Assistants

- Generally applicable
- Direct modelling of problems

Other Tools

Computer Algebra

- Solving equations, simplifications, numerical approximations
- Maple, Mathematica, ...

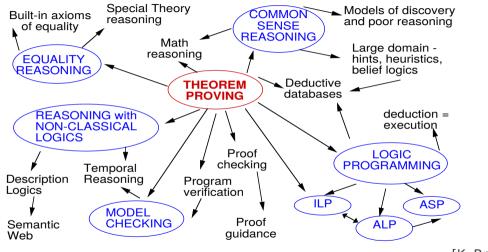
Model Checkers

- Space state abstraction
- Spin, Uppaal, ...

ATPs

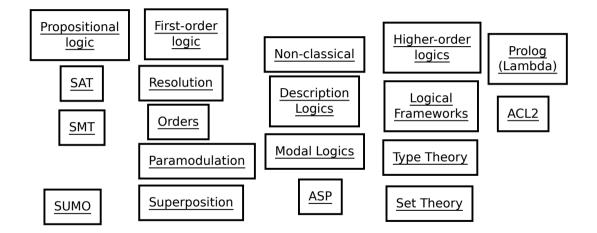
- Built in automation (model elimination, resolution)
- ACL2, Vampire, Eprover, SPASS, ...

Spread of theorem proving (1/2)



[K. Broda]

Spread of theorem proving (2/2)



Users of Proof Assistants

Computer Science

- Modelling and specifying systems
- Proving properties of systems
- Proving software correct

Mathematics

- Defining concepts and theories
- Proving (mostly verifying) proofs
- (currently less common)

Theorems and programs that use ITP

Theorems

- Kepler Conjecture (2014)
- 4 color theorem
- Feit-Thomson theorem (2012)

Software

- Processors and Chips
- Security Protocols
- Project Cristal (Comp-Cert)
- L4-Verified
- Java Bytecode

History of Proof Assistants

λ -calculus (Church, 1940)

- Simple Type Theory
- Higher-Order Logic

Formulas as Types (Curry-Howard, de Bruijn)

- Proofs as Terms
- Reduce Proof Checking to Type Checking

Automath

First implementation

LCF (Milner)

ML programming language

Multitude of Proof Assistants

Characterized by various

- Foundations
- Interaction models
- Automation strategies
- Libraries
- Size of trusted core

Examples

 HOL (HOL4, HOL-Light, ProofPower, HOL0), Mizar (and variants), PVS, Coq, Otter/Ivy, Isabelle/Isar (HOL, ZF, CTT, ...), Alfa/Agda, ACL2, IMPS, Metamath, Theorema, Lego, Nuprl, Ωmega, B method, Minlog

Coverage of Basic Mathematics

Freek Wiedijk's list of 100 th	neorems
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Isabelle	89
HOL Light	87
Coq	79
Lean	76
MetaMath	74
Mizar	69
any	99

http://www.cs.ru.nl/~freek/100/

Summary

This Lecture

- What is a Proof Assistant
- Common Uses
- Comparison with other tools
- Formal proof examples
- De Bruijn factor
- History
- Characteristics
- Coverage of Basic Mathematics

Next

- LCF and HOL Light
- Introduction to the λ-calculus

Homework / Work here

Have a look at an OCaml introduction and familiarize yourself with:

- Toplevel interaction (loading of files)
- Algebraic Datatypes
- Pattern Matching

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Tasks:

- How would you define the type of propositional logic terms?
- What about first-order logic?
- Can you define some basic operation like checking if a term is in CNF?