

This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass. ***Explain your answers to the first four exercises!***

- [1] Consider the three boolean functions f_1 , f_2 and f_3 defined as follows:

$$f_i(x_1, x_2, x_3) = \begin{cases} x_1 & \text{if } s = 0 \\ x_i & \text{if } s = 1 \\ x_s & \text{if } s > 1 \end{cases}$$

for $i \in \{1, 2, 3\}$. Here $s = x_1 + x_2 + x_3$ is the sum of the inputs, which evaluates to a natural number between 0 and 3.

- [7] (a) Compute the algebraic normal forms of f_1 , f_2 and f_3 .
 [7] (b) Is f_1 monotone? Is f_2 self-dual? Is f_3 affine? Is $\{f_1, f_2, f_3\}$ adequate?
 [6] (c) How many different affine boolean functions of arity 3 exist?

- [7] [2] (a) Use resolution to determine satisfiability of the CNF

$$\varphi = (p \vee q \vee r) \wedge (p \vee \neg q \vee r) \wedge (q \vee \neg r) \wedge (\neg p \vee r)$$

- [7] (b) Transform the following formula ψ into an equisatisfiable Skolem normal form:

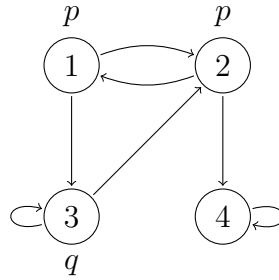
$$\psi = \neg(\forall x \exists y P(x, f(y, x)) \rightarrow \forall z Q(z)) \vee \forall x R(x, x)$$

- [6] (c) Use the unification algorithm to determine whether the terms $f(f(a, g(x)), f(g(y), z))$ and $f(f(z, w), f(g(x), w))$ are unifiable and compute a most general unifier if possible. Here x, y, z, w are variables and a is a constant.

- [3] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.

- [7] (a) $\forall x Q(x, g(x)) \vdash \exists y \forall x Q(x, y)$
 [6] (b) $p \rightarrow \neg r, q \rightarrow \neg r \vdash r \rightarrow \neg(q \vee p)$
 [7] (c) $\forall x (Q(a, x) \rightarrow a = x) \vdash \forall x (Q(a, b) \wedge P(a, x) \rightarrow P(b, x))$

4 Consider the model \mathcal{M} :



- [6] (a) Use the CTL model checking algorithm to determine in which states of \mathcal{M} the CTL formula $\varphi = \text{AX E}[p \text{ U } \text{AG } \neg q]$ holds.
- [7] (b) Determine in which states of \mathcal{M} the LTL formulas $\psi_1 = \text{FX}(p \vee q)$, $\psi_2 = p \rightarrow \text{G} q$, and $\psi_3 = q \text{ U } p$ hold.
- [7] (c) Give a CTL formula without the connectives **AX** and **AG** that is equivalent to **AX AG** p .

[20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

The sequent $\perp \vdash p \rightarrow q$ is valid.

Every Skolem normal form is a prenex normal form.

The term $f(x, y)$ is free for x in $\forall x P(f(x, y)) \vee \forall y P(y)$.

The LTL formulas $\text{G}(\varphi \vee \psi)$ and $\text{G} \varphi \vee \text{G} \psi$ are semantically equivalent.

The set $\{\text{AX}, \text{AF}, \text{AG}\}$ is an adequate set of temporal connectives for CTL.

An adequate set of boolean functions cannot contain any affine boolean functions.

The reduced OBDD representation of a boolean function for a given variable order is unique.

Every propositional formula can be transformed into an equisatisfiable DNF using Tseitin's transformation.

The substitution $\{z \mapsto g(x, y), y \mapsto g(x, y), x \mapsto h(g(x, y))\}$ is a unifier of the terms $f(g(x, y), y, x)$ and $f(z, z, h(z))$.

If there exists a model which satisfies a given predicate logic formula φ then there must also exist a natural deduction proof for φ .