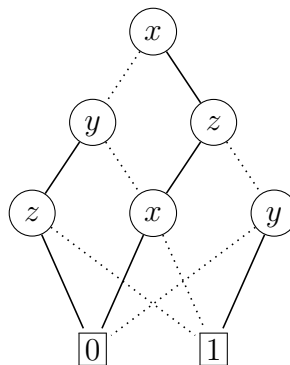


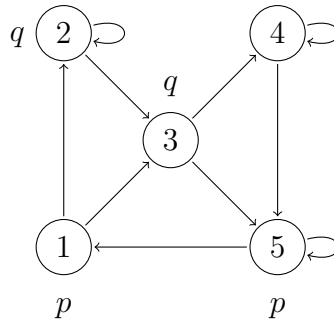
This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass. **Explain your answers to the first four exercises!**

- [1] Consider the boolean function  $f(x, y, z) = (x \oplus \bar{y}) \cdot (\bar{x} + z)$  and the BDD  $B_g$



- [2] (a) Is  $B_g$  reduced? Is  $B_g$  ordered?
- [5] (b) Transform  $B_g$  into an equivalent reduced OBDD with variable ordering  $[x, z, y]$ .
- [5] (c) Compute the algebraic normal forms of  $f$  and  $g$ .
- [5] (d) Which of the properties from Post's adequacy theorem hold for  $f$  and  $g$ ?
- [3] (e) Which subsets of  $\{f, g, \oplus, +\}$  are adequate?
- [6] [2] (a) Determine whether the terms  $f(h(z), g(x, x), z)$  and  $f(h(x), y, h(y))$  are unifiable and compute a most general unifier if possible. Here,  $x, y$  and  $z$  are variables.
- [7] (b) Transform the following formula into an equisatisfiable Skolem normal form:
- $$\varphi = \forall x \exists y (P(x) \rightarrow P(y)) \rightarrow \forall y \exists x Q(x, y)$$
- [7] (c) Use resolution to determine satisfiability of the clausal form
- $$\{ \{Q(x), \neg P(f(x), f(x))\}, \{R(f(x), y), Q(x)\}, \{\neg Q(a)\}, \\ \{R(x, f(y)), R(f(u), v)\}, \{P(x, y), \neg R(x, y)\} \}$$
- Here  $a$  is a constant.
- [3] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.
- [6] (a)  $\neg(\neg q \wedge p) \vdash q \vee \neg p$
- [7] (b)  $\forall x R(x) \vee \forall x \exists y S(x, y) \vdash \forall x \exists y (S(x, y) \vee R(x))$
- [7] (c)  $\forall x R(x) \wedge \forall x \exists y S(x, y) \vdash \exists y \forall x (S(x, y) \wedge R(x))$

4 Consider the following model  $\mathcal{M}$ :



- [6] (a) Use the CTL model checking algorithm to determine in which states of  $\mathcal{M}$  the CTL formula  $\varphi = \mathbf{EG}((p \vee q) \wedge (q \rightarrow \mathbf{EX} \neg q))$  holds.
- [7] (b) Specify a path which satisfies LTL formula  $q \wedge \mathbf{X}(\neg q \wedge (p \mathbf{U} q))$ .
- [7] (c) For each  $1 \leq i \leq 5$  find a CTL formula  $\chi_i$  without  $\mathbf{AX}$  and  $\mathbf{EX}$  which holds only in state  $i$  of  $\mathcal{M}$ .

[20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

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The sequent  $p \rightarrow q \vdash \neg p \rightarrow \neg q$  is not valid.

The function  $f(x, y) = x \oplus y \oplus \bar{y}$  is monotone.

Resolution is complete but not sound for predicate logic.

Satisfaction of CTL formulas in finite models is decidable.

In DPLL any backjump can always be simulated by a backtrack instead.

It is not possible to verify if some comparator network is a sorting network.

The CTL formula  $p \wedge \mathbf{EXEF} p$  is semantically equivalent to the LTL formula  $\mathbf{F} p$ .

For every boolean function there exist at least two different reduced BDD representations.

The Skolem normal form of a predicate logic formula cannot contain any existential quantifiers.

The substitution  $\{x \mapsto h(z), y \mapsto h(z)\}$  is a most general unifier of the terms  $f(x, y, g(x))$  and  $f(h(z), h(z), g(z))$ .