

This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [1] For each of the following sequents, either give a natural deduction proof or explain why such a proof does not exist.

- [5] (a) $p \vee q \vdash \neg(\neg p \vee q)$
 [6] (b) $\neg(\neg p \vee \neg q) \vdash (p \wedge q)$
 [6] (c) $\vdash ((p \rightarrow q) \rightarrow q) \rightarrow ((q \rightarrow p) \rightarrow p)$

- [2] Consider the boolean function

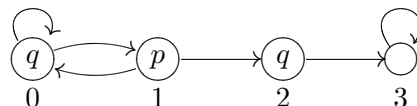
$$f(x, y, z) = \begin{cases} 0 & \text{if } x = z \text{ or } y = \min\{x, z\} \\ 1 & \text{otherwise} \end{cases}$$

- [7] (a) Give a binary decision tree for f with the variable ordering $[z, y, x]$ and use the reduce algorithm to construct an equivalent reduced OBDD.
 [7] (b) Use Shannon's expansion to obtain a boolean expression that is equivalent to f .
 [7] (c) Construct a reduced OBDD for $\forall z.f$ by using **apply** and **restrict**. Give all intermediate OBDDs.

- [3] For each of the following formulas of predicate logic, either give a natural deduction proof or find a model which does not satisfy it:

- [7] (a) $\phi_1 = \exists x (P(x) \vee Q(x)), \forall x \neg Q(x), \forall x (R(x) \rightarrow \neg P(x)) \vdash \exists x \neg R(x)$
 [7] (b) $\phi_2 = \exists x (P(x) \vee Q(x)), \exists x \neg Q(x), \forall x (R(x) \rightarrow \neg P(x)) \vdash \exists x \neg R(x)$
 [7] (c) $\phi_3 = \forall x (P(x) \vee Q(x)), \exists x \neg Q(x), \forall x (R(x) \rightarrow \neg P(x)) \vdash \exists x \neg R(x)$

- [4] Consider the model \mathcal{M} :



- [7] (a) Determine in which states of \mathcal{M} the CTL formula $\phi = A[p \text{ U } q] \rightarrow E[(p \rightarrow AX q) \text{ U } EG q]$ holds.
 [7] (b) Transform ϕ into an equivalent CTL formula ψ without the connective EG.
 [7] (c) Give an LTL formula χ that holds only in states 1 and 2 of \mathcal{M} .

- [20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

LTL is more expressive than CTL.

Every boolean function admits a minimum size reduced BDD which is ordered.

Every predicate logic formula has an equisatisfiable Skolem normal form.

Every function $F : \mathcal{P}(A) \rightarrow \mathcal{P}(A)$ with finite A has a greatest fixed-point.

A logic program is a finite non-empty set of clauses.

$\{\oplus, \wedge, \vee\}$ is an adequate set of propositional connectives.

$\forall x \phi \wedge \forall x \psi \dashv\vdash \neg \exists x (\neg \phi \vee \neg \psi)$

A set Γ of sentences of predicate logic is unsatisfiable if and only if all finite subsets of Γ are unsatisfiable.

The natural deduction proof rules LEM, PBC and $\neg\neg$ i are inter-derivable with respect to the basic proof rules.

$\{\neg, \vee, \text{AX}, \text{EF}, \text{EU}\}$ is an adequate set of connectives for CTL.