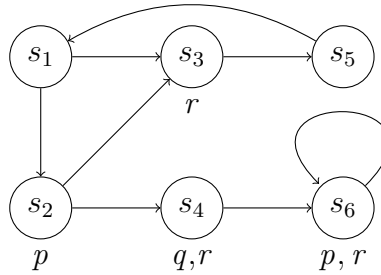




This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [1] Consider the boolean function  $f(x, y, z) = x \oplus xy \oplus yz$ .
- [9] (a) Give a binary decision tree for  $f$  with the variable ordering  $[x, y, z]$  and use the reduce algorithm to construct an equivalent reduced OBDD.
- [9] (b) Determine all minimal adequate subsets of  $\{\rightarrow, \vee, \perp, f\}$ .
- [8] [2] (a) Transform the formulas
- i.  $\forall x \neg \forall y (\forall z P(x, z) \rightarrow Q(f(y), x)) \vee \exists x \forall y P(x, y)$
  - ii.  $\exists x ((\forall y P(x, y) \rightarrow \forall y (P(x, y) \wedge Q(y, x))) \rightarrow \forall y Q(x, y))$
- into equisatisfiable clausal forms.
- [12] (b) Use resolution to determine whether the clausal forms
- i.  $\{ \{P(x, y), Q(x, f(y)), R(x, y)\}, \{ \neg P(x, b), Q(x, f(y)), \neg R(x, f(y))\}, \{ \neg R(x, b)\}, \{R(f(a), f(x))\}, \{ \neg Q(f(a), f(b))\} \}$
  - ii.  $\{ \{ \neg P(x, y), Q(b, x), R(x, x)\}, \{P(x, y), \neg Q(z, f(u)), R(f(a), z)\}, \{P(x, y), Q(f(a), y), R(z, b)\}, \{ \neg R(x, x)\}, \{ \neg Q(b, f(x))\} \}$
- are satisfiable. (Here  $a, b$  are constants and  $u, x, y, z$  variables.)
- [3] For each of the following formulas of predicate logic, either give a natural deduction proof or find a model which does not satisfy it:
- [7] (a)  $\phi_1 = \forall x \exists y (P(y) \rightarrow Q(x)) \rightarrow \forall x (\exists y P(y) \rightarrow Q(x))$
- [7] (b)  $\phi_2 = \forall x \exists y (P(x) \rightarrow Q(y)) \rightarrow \forall x (P(x) \rightarrow \exists y Q(y))$
- [7] (c)  $\phi_3 = \forall x (\exists y P(y) \rightarrow Q(x)) \rightarrow \exists y \forall x (P(y) \rightarrow Q(x))$

- 4 Consider the CTL formulas  $\phi = E[p \vee \neg r \ U \ EX \ AG \ \neg q]$  and  $\psi = AG \ AF(p \vee q)$ , the LTL formula  $\chi = GFp \vee GFq$ , and the model  $\mathcal{M}$ :



- [7] (a) Apply the CTL model checking algorithm to determine the states of  $\mathcal{M}$  which satisfy  $\phi$ .
- [7] (b) Provide an LTL formula  $\xi$  which does not contain the proposition  $r$  such that  $\mathcal{M}, s \models \xi$  if and only if  $s = s_3$ .
- [7] (c) Determine whether  $\psi$  is equivalent to  $\chi$ . In the negative case also provide a counterexample.

- [20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

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Every clause with the property that two of its literals unify has a factor.

To represent the transition relation of a CTL model with 16 states we need at most 8 boolean variables.

$$\forall x (\exists x \phi \rightarrow \psi) \dashv\vdash \exists x \phi \rightarrow \forall x \psi$$

Executing the Prolog query `?- select(b, [a,b,c,d], X).` produces the answer `X = [a,c,d]`.

A set of sentences of predicate logic is satisfiable if and only if all finite subsets are satisfiable.

A boolean function  $f: \{0, 1\}^n \rightarrow \{0, 1\}$  is self-dual if  $f(x_1, \dots, x_n) = f(\overline{x_1}, \dots, \overline{x_n})$ .

Every CTL\* path formula is a CTL\* state formula.

If the set  $S$  has  $n$  elements then  $F^n(S)$  is the least fixed point of a monotone function  $F: \mathcal{P}(S) \rightarrow \mathcal{P}(S)$ .

Every Horn formula is satisfiable.

The LTL formulas  $\neg(\phi R \psi)$  and  $\neg\psi \ U \ \neg\phi$  are semantically equivalent.