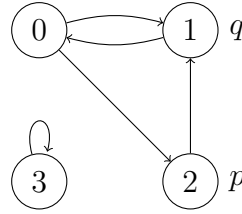


This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [1] Consider the boolean function  $f(x, y, z) = 1 \oplus xy \oplus yz \oplus xyz$ .
- [8] (a) Give a binary decision tree for  $f$  with the variable ordering  $[x, y, z]$  and use the reduce algorithm to construct an equivalent reduced OBDD.
- [4] (b) Show that the constant function  $z(x) = 0$  can be expressed in terms of  $f$ .
- [8] (c) Determine all minimal adequate subsets of  $\{\wedge, \vee, \rightarrow, f\}$ .
- [12] [2] (a) Use first-order resolution to determine the validity of the predicate logic formula
- $$\phi = \forall x (\neg P(x) \rightarrow P(f(x))) \rightarrow (P(a) \rightarrow \exists x (P(x) \wedge P(f(f(x))))))$$
- [4] (b) Solve the following instance of Post's correspondence problem given as a sequence of pairs  $(s_i, t_i)$ :
- $$(11, 1), (1, 111), (0111, 10), (10, 0).$$
- [4] (c) Are the terms  $s = h(g(f(a)), f(g(y)), f(x))$  and  $t = h(g(x), f(g(y)), y)$  unifiable?
- [3] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.
- [6] (a)  $\vdash \forall x \exists y P(x, y) \rightarrow \exists y \forall x P(x, y)$
- [7] (b)  $\vdash \exists y \forall x P(x, y) \rightarrow \forall x \exists y P(x, y)$
- [7] (c)  $\neg(\neg p \wedge \neg q) \vdash p \vee q$

4 Consider the model  $\mathcal{M}$ :



[8] (a) Determine in which states of  $\mathcal{M}$  the CTL formula  $\phi = A[EX p \cup EX q]$  holds.

[8] (b) Determine in which states of  $\mathcal{M}$  the LTL formula  $\psi = Xp \cup Xq$  holds.

[4] (c) Give a CTL formula  $\phi$  such that  $\mathcal{M}, s \models Xp \cup Xq$  implies  $\mathcal{M}, s \models AF\phi$  for all models  $\mathcal{M}$  and states  $s$ .

[20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

---

Peano arithmetic is decidable.

Double negation elimination is a valid proof rule of natural deduction.

Restarts are part of basic DPLL.

The LTL formulas  $\phi W \neg\phi$  and  $\phi \vee X\top$  are semantically equivalent.

A Horn clause  $P_1 \wedge P_2 \rightarrow Q$  is satisfiable if and only if  $Q \in \{P_1, P_2\}$ .

The unification problem  $\{x \stackrel{?}{=} y, y \stackrel{?}{=} z\}$  is in solved form.

The boolean functions  $f(x, y, z) = x \oplus yz \oplus xy$  and  $g(x, y, z) = y \oplus xz \oplus xy$  have the same algebraic normal form.

The CTL formulas  $A[\phi \cup \neg\phi]$  and  $\neg EG \phi$  are semantically equivalent.

Every prenex normal form is in Skolem normal form.

There exists a sorting network with 16 wires and 60 comparators.