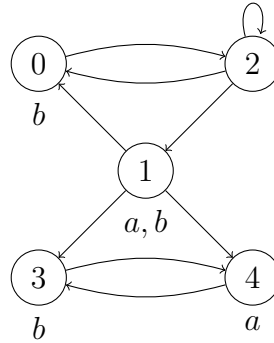


This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [1] Consider the propositional formula $\phi = ((p \vee q) \wedge (\neg p \vee \neg r)) \wedge (\neg p \vee r)$.
- [6] (a) Test the satisfiability of ϕ using DPLL.
- [7] (b) Test the satisfiability of ϕ using the linear SAT solver.
- [7] (c) Test the satisfiability of ϕ using the cubic SAT solver.
- [6] [2] (a) Are the terms $s = f(g(x, y), h(a), h(x))$ and $t = f(g(h(a), h(h(a))), x, h(y))$ unifiable?
- [7] (b) Compute a Skolem normal form for
- $$P(a) \wedge \exists x (P(x) \wedge \neg \forall x (Q(x) \rightarrow \forall y P(y)))$$
- [7] (c) Use resolution to decide if the following set of clauses is satisfiable:
- $$\{\{Q(a)\}, \{P(x), \neg Q(x)\}, \{Q(f(f(a))), \neg Q(a)\}, \{\neg Q(f(a))\}\}$$
- [3] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.
- [7] (a) $\forall x (P(x) \rightarrow P(f(x))), \forall x (P(x) \rightarrow x = a), \exists x P(x) \vdash f(a) = a$
- [7] (b) $\exists x (P(x) \rightarrow P(f(x))), \forall x (P(x) \rightarrow x = a), \exists x P(x) \vdash f(a) = a$
- [6] (c) $\vdash ((p \rightarrow q) \rightarrow p) \rightarrow p$

4 Consider the model \mathcal{M} :



- [8] (a) Determine in which states of \mathcal{M} the CTL formula $\phi = AX E[a U b \rightarrow a]$ holds.
- [6] (b) Give a CTL formula which holds only in states 2 and 3.
- [6] (c) Are the CTL formulas $AF((AG p) \vee q)$ and $AF AG(p \vee q)$ equivalent? Prove your answer.

[20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

The propositional formula $(p \vee q) \wedge (\neg p \vee \neg q)$ is satisfiable.

$HWB_5(0, 1, 0, 1, 0) = 1$

The set $\{X, F, G\}$ is an adequate set of temporal LTL connectives.

The propositional formula $\top \vee P \vee \perp \rightarrow Q$ is a Horn clause.

The CTL* formulas $EG F p$ and $EG EF p$ are semantically equivalent.

The term $f(x, y)$ is free for y in $\exists x (P(x, y) \vee \forall y Q(x, f(y, z)))$.

Every sorting network with n wires has at least n comparators.

Peano arithmetic is consistent.

The algebraic normal form of the boolean function $f(x, y) = x\bar{y} \oplus (x + y)$ is y .

For every CTL formula ϕ the set $\llbracket EF \phi \rrbracket$ is a greatest fixed point of $F(X) = \llbracket \phi \rrbracket \cup \text{pre}_{\exists}(X)$.