

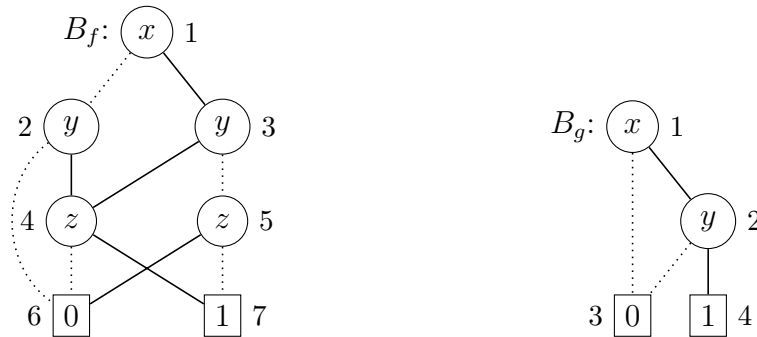
This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [5] 1 (a) Are the terms  $s = f(g(x, a), h(y, b))$  and  $t = f(g(c, z), z)$  unifiable? Here  $a$  and  $b$  are constants and  $x, y$  and  $z$  are variables.
- [7] (b) Compute a Skolem normal form for the formula  $\forall x \exists y P(x, y) \wedge \forall x Q(x)$ .
- [8] (c) Use resolution to decide if the following set of clauses is satisfiable:

$$\{\{P(a, f(b))\}, \{\neg R(c, f(b))\}, \{P(x, y), R(x, y)\}, \{\neg P(x, y), \neg R(x, y)\}, \\ \{P(x, y), \neg P(y, x)\}, \{\neg P(x, y), \neg P(y, z), P(x, z)\}, \{R(a, c)\}\}$$

Here  $a, b$  and  $c$  are constants and  $x, y$  and  $z$  are variables.

- 2 Consider the following two OBDDs.

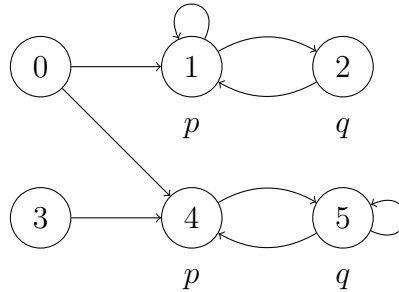


- [6] (a) Compute a reduced OBDD for  $\forall y.g$ , where  $g$  is given by  $B_g$ .
- [7] (b) Use Shannon's expansion to obtain a boolean expression that is equivalent to  $B_f$ . Simplify your answer.
- [7] (c) Compute  $\text{apply}(+, B_f, B_g)$ .

- 3 For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.

- [6] (a)  $p \vdash ((p \rightarrow q) \rightarrow p) \rightarrow p$
- [7] (b)  $P(a), \exists x (x = a \rightarrow \neg P(x)) \vdash \exists x \neg(x = a)$
- [7] (c)  $Q(a), \forall x (Q(x) \rightarrow Q(f(f(x)))) \vdash \neg Q(f(a)) \vee \forall x Q(x)$

4 Consider the model  $\mathcal{M}$ :



- [7] (a) Determine in which states of  $\mathcal{M}$  the CTL formula  $\phi = \text{EXAF A}[p \text{ U } q]$  holds.
- [7] (b) Give an LTL formula  $\psi$  such that  $\mathcal{M}, s \models \psi$  if and only if  $s = 5$ .
- [6] (c) Find a model and identify a state in that model such that the CTL\* formula

$$\text{A}[\text{GF } p \wedge \text{GE}[\text{F } q]] \wedge \neg \text{A}[p \rightarrow \text{F } q]$$

holds in that state.

- [20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

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The depth of any sorting network is larger than its size.

Presburger arithmetic is finitely axiomatizable.

For every boolean function  $f$  and variable  $x$ ,  $f \equiv \bar{x} \cdot f[0/x] \oplus x \cdot f[1/x]$ .

Natural deduction for propositional logic without  $\neg\neg$ e is sound.

Basic DPLL contains the pure literal inference rule.

Every binary self-dual boolean function is monotone.

The CTL formulas  $\text{EF } \phi \wedge \text{EF } \psi$  and  $\text{EF}(\phi \wedge \psi)$  are semantically equivalent.

$\forall x \phi \rightarrow \exists x \psi \dashv\vdash \exists x (\phi \rightarrow \psi)$

The boolean function  $f(x, y, z) = x + xy$  is affine.

Natural deduction for propositional logic without  $\neg\neg$ i is complete.