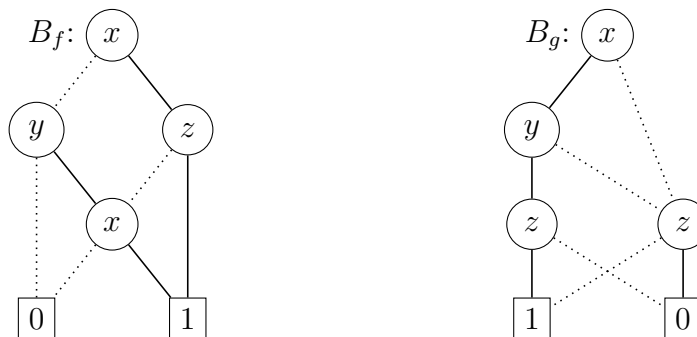


This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [1] Consider the Boolean function $f(x, y, z) = (\bar{y} + z) \oplus y(x + z)$.
- [6] (a) Give precise definitions of the following properties of Boolean functions: monotonicity, self-duality, and affinity.
- [7] (b) Compute an algebraic normal form of f .
- [7] (c) Determine all minimal complete adequate subsets of $\{f, \perp, \rightarrow\}$.

- [2] Consider the BDDs

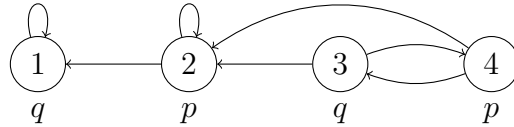


- [7] (a) Use Shannon's expansion to obtain a boolean expression equivalent to B_f . Simplify your answer.
- [6] (b) Transform B_g into an equivalent reduced OBDDs with variable ordering $[y, z, x]$.
- [7] (c) Compute a reduced OBDD for $\forall z.g$, where g is given by B_g .

- [3] Either give a natural deduction proof or find a model which does not satisfy the following sequents. In the natural deduction proofs you may additionally use the quantifier equivalences (e.g. $\neg\forall x \phi(x) \dashv\vdash \exists x \neg\phi(x)$).

- [5] (a) $\forall x \forall y (x = y \rightarrow P(x, y)), \exists x \neg(P(x, a) \vee P(a, x)) \vdash \neg(a = b) \rightarrow P(a, b)$
- [10] (b) $\vdash \exists x (\neg P(x) \rightarrow \forall y \neg P(y))$
- [5] (c) $\vdash \forall x (\neg P(x) \rightarrow \forall y \neg P(y))$

4 Consider the following model \mathcal{M} :



- [7] (a) Determine in which states of \mathcal{M} the CTL formula $\phi = E[(A[EF q U EG p] \vee AX p) U AG q]$ holds.
- [7] (b) Give an LTL formula ψ such that $\mathcal{M}, s \models \psi$ if and only if $s = 4$.
- [6] (c) Show that the two CTL formulas with fairness constraints $E_{\{p,q\}}G \top$ and $E_{\{p \wedge q\}}G \top$ are not equivalent.

[20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

There exists a CTL model \mathcal{M} and a state s such that $\mathcal{M}, s \models E_{\{\neg\phi\}}G \phi$.

The theorem of Löwenheim and Skolem states that a set of sentences of predicate logic is satisfiable if all its finite subsets are satisfiable.

The quantifier-free fragment of the theory of arrays is decidable.

The terms $f(g(x), x, y)$ and $f(z, y, g(z))$ are unifiable. Here x, y , and z are variables.

The set $\{U, R\}$ is an adequate set of connectives for the LTL fragment consisting of negation-normal forms without X .

There exists a sorting network for 16 wires with 60 comparators.

The LTL formulas $\phi R \psi$ and $\psi W (\psi \wedge \phi)$ are semantically equivalent.

The predicate logic formula $\forall x \exists y ((P(x) \wedge Q(x, y)) \vee \neg Q(y))$ is in Skolem normal form.

The boolean function $f(x, y, z) = x + xy$ is self-dual.

The use of the backjump rule in abstract DPLL is known as chronological backtracking.