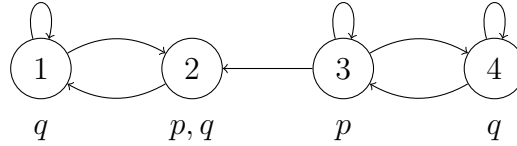


This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [1] Consider the boolean function $f(x, y) = (x + y) \cdot \overline{(x \cdot y)}$.
- [4] (a) Compute the algebraic normal form of f .
- [3] (b) Is $\{f\}$ adequate?
- [6] (c) Show that there is no nullary, unary, or binary boolean function that is self-dual but not affine. (Hint: There are only finitely many boolean functions with arity ≤ 2 .)
- [7] (d) Find a ternary boolean function that is self-dual but not affine.
- [6] [2] (a) Determine whether the terms $f(f(x, y), f(x, z))$ and $f(z, f(y, z))$ are unifiable and compute a most general unifier, if possible.
- [7] (b) Compute a Skolem normal form for the predicate logic formula
- $$P(a) \wedge \neg \forall x (P(x) \rightarrow \exists x (Q(x) \rightarrow \forall y P(y)))$$
- [7] (c) Use resolution to decide if the following set of clauses is satisfiable:
- $$\{\{P(a)\}, \{Q(x), \neg P(f(x))\}, \{P(f(f(x))), \neg P(x)\}, \{\neg Q(f(a))\}\}$$
- [3] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.
- [6] (a) $q \rightarrow p \vdash p \rightarrow \neg q \rightarrow p$
- [7] (b) $\forall x (Q(x) \rightarrow P(x)), \exists x \neg P(x) \vdash \exists x \neg Q(x)$
- [7] (c) $\forall x (Q(x) \rightarrow P(x)), \exists x \neg Q(x) \vdash \exists x \neg P(x)$

4 Consider the following model \mathcal{M} :



- [7] (a) Use the labelling algorithm to determine in which states of the model \mathcal{M} the CTL formula $E[EX p \cup AX q \wedge EG q]$ holds.
- [6] (b) Give a CTL* formula ψ such that $\mathcal{M}, s \models \psi$ if and only if $s = 4$.
- [7] (c) Show that the CTL formulas with fairness constraints $E_{\{p,q\}}G \top$ and $E_{\{p \vee q\}}G \top$ are not equivalent.

[20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

The LTL formulas $\neg(\phi R \psi)$ and $\neg\phi U \neg\psi$ are semantically equivalent.

The compactness theorem holds for universal second-order logic.

A boolean function is affine if it can be written in algebraic normal form.

Every Horn formula can be written in conjunctive normal form.

We have $\mathcal{M} \models \phi$ for sentences ϕ in predicate logic.

The set of connectives $\{\perp, \vee, \leftrightarrow\}$ is adequate for propositional logic.

There is a boolean function f with n arguments such that every OBDD computing f has size exponential in n , for all natural numbers n .

$$s = t, t = u \vdash s = u$$

Validity in predicate logic is decidable using resolution with factoring.

Gödel's incompleteness theorem states that there is a sentence that is true in the standard model of arithmetic but not provable in PA.