

This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [6] [1] (a) Use Tseitin's transformation to compute an equisatisfiable CNF of the following propositional formula:

$$\phi = \neg(r \rightarrow (q \wedge p)) \rightarrow ((\neg p \wedge r) \vee q)$$

- [7] (b) Give a precise definition of the backjump rule in (basic) DPLL.

- [7] (c) Use DPLL to decide satisfiability of the following propositional formula:

$$\psi = (r \vee q) \wedge (t \vee \neg q \vee p) \wedge (\neg t \vee \neg q) \wedge (t \vee \neg p) \wedge (\neg r \vee q)$$

- [2] Consider the boolean function  $f(x, y, z) = xy \oplus z$  and the following two reduced OBDDs:

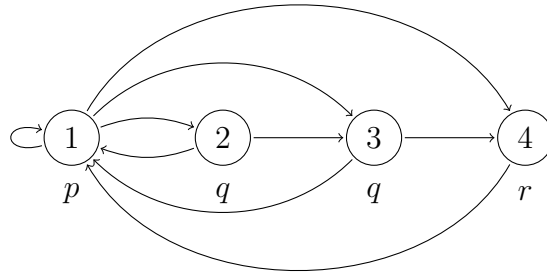


- [6] (a) Construct a reduced OBDD for  $f$  with variable ordering  $[x, y, z]$ .  
[7] (b) Compute  $\text{apply}(\cdot, B_g, B_h)$ .  
[7] (c) Starting from  $B_g$ , compute a reduced OBDD that is equivalent to  $\forall y.g$ .

- [3] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it. (Hint: The sequent of part (b) is valid. Using LEM on  $\forall y P(y)$  as first step does not hurt.)

- [6] (a)  $\vdash (p \rightarrow q) \vee (q \rightarrow p)$   
[7] (b)  $\vdash \exists x \forall y (P(x) \rightarrow P(y))$   
[7] (c)  $\vdash \exists x \forall y (\neg Q(y, y) \rightarrow Q(x, y))$

4 Consider the model  $\mathcal{M}$



- [6] (a) For each of the LTL formulas  $\phi_1 = X(p \text{ U } r)$ ,  $\phi_2 = \text{GF}p \wedge \text{GF}q \wedge \text{GF}r$ , and  $\phi_3 = \text{G}(p \rightarrow Xq)$ , find a path  $\pi_i$  starting at state 1 such that  $\pi_i \models \phi_i$  for each  $i \in \{1, 2, 3\}$ .
- [7] (b) Give an LTL formula that distinguishes states 2 and 3 of  $\mathcal{M}$ .
- [7] (c) Use the labeling algorithm to determine in which states of the model  $\mathcal{M}$  the CTL formula  $\psi = \text{E}[q \text{ U } (\text{EX}p \wedge \text{AF}r)]$  holds.

[20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

---

Peano arithmetic is decidable.

$\forall x P(x) \wedge \forall y Q(y) \vdash \exists x (P(x) \wedge \forall y Q(y))$

Validity in propositional logic is undecidable.

The boolean function  $\text{HWB}_2(x, y)$  is self-dual.

Every Skolem normal form is a prenex normal form.

The boolean function  $f(x, y, z) = \overline{(x + z)} \cdot y$  is adequate.

Every pair of unifiable terms has a unique most general unifier.

The set  $\{\text{AX}, \text{AG}, \text{AU}\}$  is an adequate set of temporal connectives for CTL.

A sentence in predicate logic has a finite model if it has an infinite model.

The set of temporal LTL connectives  $\{\text{U}, \text{G}\}$  is adequate for the LTL fragment consisting of negation normal forms without  $\text{X}$ .