

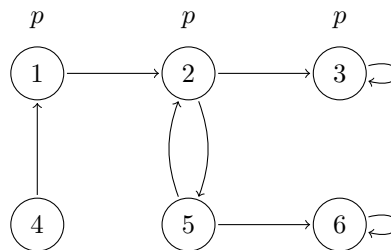
This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [1] Consider the boolean function  $f(x, y, z) = x \oplus (y + z)$ .
- [7] (a) Is  $f$  monotone? Is  $f$  self-dual? Is  $f$  affine?
- [6] (b) Can  $\bar{x}$  be expressed (only) using  $f$ ?
- [7] (c) Let  $g$  be a boolean function such that  $\{g\}$  is adequate. Is  $\{\hat{g}\}$  adequate? Here  $\hat{g}$  is the dual of  $g$ .
- [7] [2] (a) Use the unification algorithm from the lecture to determine whether the two terms  $h(f(x, z), b, f(z, y))$  and  $h(y, z, f(b, f(a, w)))$  are unifiable and compute a most general unifier if possible. Note that  $a$  and  $b$  are constants while  $w, x, y$  and  $z$  are variables.
- [5] (b) Explain how resolution can be used to determine the validity of a formula  $\phi$  in predicate logic. You do not have to explain resolution.
- [8] (c) Use resolution to determine whether the formula

$$\phi = \neg \forall x \exists y \neg P(x, y) \wedge (\forall z \exists w P(w, z) \rightarrow \forall z' Q(z'))$$

is satisfiable.

- [3] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.
- [6] (a)  $p \vee q, \neg r \rightarrow \neg p \vdash q \vee r$
- [7] (b)  $\exists x P(x), x = y \vdash P(y)$
- [7] (c)  $\vdash \forall x ((\neg P(x) \rightarrow \perp) \rightarrow P(x))$
- [4] Consider the CTL formula  $\phi = AX(AF p \vee AG \neg p)$
- [7] (a) Use the labeling algorithm to determine in which states  $\phi$  holds, for the following model:



- [7] (b) Give a CTL formula  $\psi$  that is equivalent to  $\phi$  but does not use the temporal connectives in  $\phi$  (so  $\psi$  may not contain  $AX$ ,  $AF$ , and  $AG$ ).
- [6] (c) Given a path  $\pi$  in a model  $\mathcal{M}$ , define  $\pi \models G\varphi$ , i.e., satisfaction of the LTL formula  $G\varphi$  with respect to  $\pi$ . Using the definition, show that  $G\varphi \equiv GG\varphi$  for every LTL formula  $\varphi$ .

- [20] 5 Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

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The dual of any affine boolean function is monotone.

The symbols  $\zeta$  and  $\Gamma$  are letters from the Greek alphabet.

If  $F = (\neg p \vee q \vee r) \wedge \neg q$  then  $\overset{d}{p} \parallel F \implies \overset{d}{p} r \parallel F$  in DPLL.

The LTL formulas  $\phi \mathbf{R} \psi$  and  $\psi \mathbf{U}(\phi \wedge \psi) \vee \mathbf{G} \psi$  are equivalent.

The sequent  $\neg \forall x (P(x) \vee Q(y)) \vdash \neg Q(y) \wedge \exists x \neg P(x)$  is valid.

If  $f$  is a binary boolean function then  $f(x, y) = \bar{y} \cdot f(x, 0) \oplus y \cdot f(x, 1)$ .

Sorting  $n$  inputs with insertion sort requires at most  $2n + 3$  comparators.

If  $\phi \vdash \psi$  is a valid sequent in propositional logic then  $\neg \phi \vee \psi$  is a tautology.

All occurrences of the variable  $x$  in the formula  $P(y) \wedge \forall y (\exists x Q(x, y) \vee P(x))$  are bound.

If  $X$  is the set of all self-dual boolean functions, then any function constructed from functions in  $X$  is also self-dual.