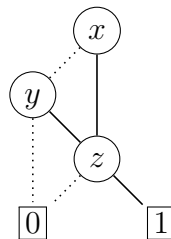


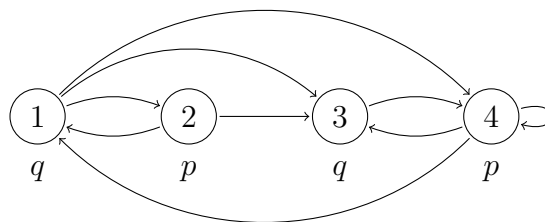
This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- 1 Consider the boolean function f defined by $f(x, y, z) = xz + yz$ with OBDD B_f :



- [7] (a) Give reduced OBDDs for f with variable orderings $[y, x, z]$ and $[z, x, y]$.
- [7] (b) Compute a reduced OBDD for $\exists y.f$.
- [6] (c) Give an expression only using f and the variables x and y that is equivalent to the boolean function xy .

- 2 Consider the model \mathcal{M}



- [7] (a) Use the CTL model checking algorithm to determine in which states of \mathcal{M} the CTL formula $\phi = E[(AX \neg p) \cup A[q \cup (EG \neg q)]]$ holds.
- [7] (b) Give an LTL formula ψ such that $\mathcal{M}, s \models \psi$ if and only if $s \in \{2, 3\}$.
- [6] (c) Are the CTL* formulas $E[XFG \neg p]$ and $\neg A[XGE[F p]]$ equivalent?

[7] [3] (a) Consider the following CNF φ :

$$(\neg 1 \vee \overset{\alpha}{\neg 3} \vee 4) \wedge (\neg 2 \vee \overset{\beta}{\neg 3} \vee \neg 5) \wedge (\neg 3 \vee \overset{\gamma}{6}) \wedge (\neg 4 \vee \overset{\delta}{\neg 6} \vee \neg 7) \wedge (\neg 6 \vee \overset{\epsilon}{8}) \wedge (\neg 4 \vee \overset{\zeta}{7} \vee \neg 8)$$

Starting from the state $\overset{d}{1} \overset{d}{2} \overset{d}{3} \parallel \varphi$, use DPLL to determine satisfiability of φ .

[7] (b) Transform the formula $\forall x \exists y \forall z \exists w (P(x, w) \rightarrow H(y) \wedge \neg Q(z, w))$ into clausal form.

[6] (c) Use resolution to show unsatisfiability of the following clausal form.

$$\{ \{ \neg P(x, f(x)), Q(f(x)) \}, \{ \neg Q(x), R(f(x)) \}, \{ \neg R(f(f(f(x)))) \}, \\ \{ P(a, f(a)) \}, \{ P(f(f(f(a))), y) \} \}$$

For each resolvent give the mgu for the clashing literals. Here a is a constant.

[4] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it. In parts (b) and (c), k and m are constants.

[7] (a) $(p \rightarrow q) \rightarrow r, q \vdash p \vee r$

[7] (b) $\forall x (P(x, m) \rightarrow \neg Q(x)), Q(m) \vdash Q(k) \rightarrow \neg P(k, m)$

[6] (c) $\forall x (P(x, m) \rightarrow \neg Q(x)), Q(m) \vdash Q(k) \rightarrow \neg(k = m)$

[20] [5] Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

If f is a boolean function then $f[0/x] \cdot f[1/x] = 0$.

The PCP instance $(1, 11), (01, 10), (00, 0)$ has a solution.

The boolean function $f(x, y, z) = \overline{(\bar{x} \oplus y)} \cdot (x + z)$ is monotone.

The propositional formula $p \vee q$ has an equisatisfiable Horn formula.

Every LTL formula is semantically equivalent to a CTL* path formula.

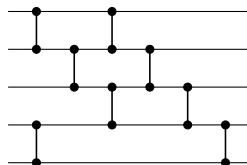
$P(f(y)) \rightarrow \forall x \exists y (\neg Q(x, y) \vee P(y)) \not\equiv \forall x (P(f(y)) \rightarrow \exists y (\neg Q(x, y) \vee P(y)))$

The CTL formulas $\mathbf{A}[\phi \mathbf{U} \psi]$ and $\neg(\mathbf{E}[\neg\psi \mathbf{U} (\neg\psi \wedge \neg\phi)] \vee \neg \mathbf{A}\mathbf{F} \psi)$ are equivalent.

Any valid sequent has a natural deduction proof in which the proof rule $\neg\neg e$ is not used.

The natural numbers can be colored with two colors such that not all of x, y, z have the same color whenever $x^2 + y^2 = z^2$ with $1 \leq x, y, z \leq 100$.

The comparator network



is a sorting network.