

This exam consists of five exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [1] In this exercise we consider the boolean function  $\text{ite}$  defined by

$$\text{ite}(x, y, z) = \begin{cases} z & \text{if } x = 0 \\ y & \text{if } x = 1 \end{cases}$$

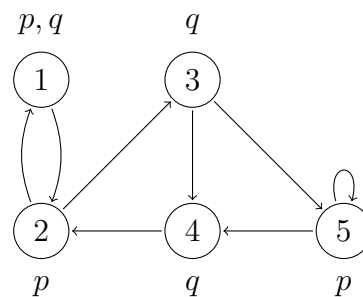
- [7] (a) Compute a reduced OBDD for  $\text{ite}$  with variable ordering  $[y, x, z]$ .  
 [7] (b) Find the missing clauses such that  $a \leftrightarrow \text{ite}(x, y, z)$  and

$$(a \vee \neg x \vee \neg y) \wedge \boxed{\phantom{a \vee \neg x \vee \neg y}} \wedge \boxed{\phantom{a \vee \neg x \vee \neg y}} \wedge (\neg a \vee x \vee z)$$

are equivalent.

- [6] (c) Is  $\text{ite}$  monotone? Is  $\text{ite}$  self-dual? Is  $\text{ite}$  affine?

- [2] Consider the model  $\mathcal{M}$ :



- [7] (a) Use the CTL model checking algorithm to determine in which states of  $\mathcal{M}$  the CTL formula  $\phi = A[p \text{ U } AF q] \wedge AX EF \neg p$  holds.  
 [6] (b) Consider the LTL formulas  $\phi_1 = G(F p \wedge F q)$ ,  $\phi_2 = p \text{ U } q \rightarrow \neg p$ , and  $\phi_3 = G(p \vee X X q)$ . Find a model  $\mathcal{M}$  and a state  $s$  such that  $\mathcal{M}, s \models \phi_i$  for all  $1 \leq i \leq 3$ .  
 [7] (c) Show that the LTL formula  $F p \rightarrow F q$  and the CTL formula  $AF p \rightarrow AF q$  are not equivalent.

- [7] [3] (a) Use resolution to determine whether the formula  $(p \rightarrow q \rightarrow p) \rightarrow p$  is valid or not.  
 [7] (b) Give a Skolem normal form of the formula  $\exists x ((\exists y P(x, y)) \rightarrow (\forall y \exists z Q(y, z)))$ .  
 [6] (c) Determine whether the terms  $h(f(x, x), y, f(a, x))$  and  $h(y, f(x, a), y)$  are unifiable and compute a most general unifier if possible. Here  $a$  is a constant and  $x$  and  $y$  are variables.
- [4] For each of the following sequents, either give a natural deduction proof or find a model which does not satisfy it.
- [6] (a)  $p \rightarrow t, \neg t \vee \neg s \vdash p \rightarrow \neg s$   
 [7] (b)  $Q(c), \exists x (Q(x) \rightarrow \forall x P(x)) \vdash P(c)$   
 [7] (c)  $\forall x Q(x), \exists x (Q(x) \rightarrow \forall x P(x)) \vdash P(c)$
- [20] [5] Determine whether the following statements are true or false. Every correct answer is worth 2 points. For every wrong answer 1 point is subtracted, provided the total number of points is non-negative.

statement

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The sequent  $\vdash \forall x (D(x) \rightarrow \exists y D(y))$  is valid.

Every binary monotone boolean function is affine.

The propositional formulas  $\perp$  and  $\top$  are equisatisfiable.

The LTL formulas  $\phi \mathbf{R} \psi$  and  $\mathbf{G} \psi \vee \psi \mathbf{U} (\phi \wedge \psi)$  are equivalent.

The hereditary base 3 representation of 333 is  $3^{3+2} + 3^{3+1} + 3^2$ .

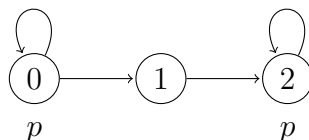
The formula  $\forall x \forall y P(x, y) \vee \neg Q(f(x))$  is a Skolem normal form.

The set  $\{\forall x \exists y P(x, y), \neg \exists x \forall y \neg P(y, x), \forall x \neg P(x, x)\}$  is consistent.

Every sorting network for  $n > 1$  wires has at least  $(n - 1)^2$  comparators.

The substitution  $\{x \mapsto y, y \mapsto x\}$  is a most general unifier of  $f(x, y)$  and  $f(y, x)$ .

In the model  $\mathcal{M}$



$\mathcal{M}, 0 \models \mathbf{AF} \mathbf{AG} p$  holds.