

Selected Solutions

1 (b) We use 9 propositional atoms x_1, \dots, x_9 :

x_i is true/false \iff number i is colored blue/red

For each solution of the equation $a + b = c$, two clauses are added:

$$(x_a \vee x_b \vee x_c) \quad (\neg x_a \vee \neg x_b \vee \neg x_c)$$

It follows that the answer for $n = 9$ is yes if and only if the formula

$$\begin{aligned} \varphi = & (x_1 \vee x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2 \vee \neg x_3) \wedge (x_2 \vee x_4 \vee x_6) \wedge (\neg x_2 \vee \neg x_4 \vee \neg x_6) \wedge \\ & (x_1 \vee x_3 \vee x_4) \wedge (\neg x_1 \vee \neg x_3 \vee \neg x_4) \wedge (x_2 \vee x_5 \vee x_7) \wedge (\neg x_2 \vee \neg x_5 \vee \neg x_7) \wedge \\ & (x_1 \vee x_4 \vee x_5) \wedge (\neg x_1 \vee \neg x_4 \vee \neg x_5) \wedge (x_2 \vee x_6 \vee x_8) \wedge (\neg x_2 \vee \neg x_6 \vee \neg x_8) \wedge \\ & (x_1 \vee x_5 \vee x_6) \wedge (\neg x_1 \vee \neg x_5 \vee \neg x_6) \wedge (x_2 \vee x_7 \vee x_9) \wedge (\neg x_2 \vee \neg x_7 \vee \neg x_9) \wedge \\ & (x_1 \vee x_6 \vee x_7) \wedge (\neg x_1 \vee \neg x_6 \vee \neg x_7) \wedge (x_3 \vee x_4 \vee x_7) \wedge (\neg x_3 \vee \neg x_4 \vee \neg x_7) \wedge \\ & (x_1 \vee x_7 \vee x_8) \wedge (\neg x_1 \vee \neg x_7 \vee \neg x_8) \wedge (x_3 \vee x_5 \vee x_8) \wedge (\neg x_3 \vee \neg x_5 \vee \neg x_8) \wedge \\ & (x_1 \vee x_8 \vee x_9) \wedge (\neg x_1 \vee \neg x_8 \vee \neg x_9) \wedge (x_3 \vee x_6 \vee x_9) \wedge (\neg x_3 \vee \neg x_6 \vee \neg x_9) \wedge \\ & (x_2 \vee x_3 \vee x_5) \wedge (\neg x_2 \vee \neg x_3 \vee \neg x_5) \wedge (x_4 \vee x_5 \vee x_9) \wedge (\neg x_4 \vee \neg x_5 \vee \neg x_9) \end{aligned}$$

is satisfiable. The corresponding DIMACS encoding is

```
p cnf 9 32
1 2 3 0 -1 -2 -3 0
2 4 6 0 -2 -4 -6 0
1 3 4 0 -1 -3 -4 0
2 5 7 0 -2 -5 -7 0
1 4 5 0 -1 -4 -5 0
2 6 8 0 -2 -6 -8 0
1 5 6 0 -1 -5 -6 0
2 7 9 0 -2 -7 -9 0
1 6 7 0 -1 -6 -7 0
3 4 7 0 -3 -4 -7 0
1 7 8 0 -1 -7 -8 0
3 5 8 0 -3 -5 -8 0
1 8 9 0 -1 -8 -9 0
3 6 9 0 -3 -6 -9 0
2 3 5 0 -2 -3 -5 0
4 5 9 0 -4 -5 -9 0
```

Executing `z3 -dimacs color.txt`, with the file `color.txt` containing the above DIMACS encoding, yields the output

```
s UNSATISFIABLE
```

So the answer for $n = 9$ is no.

3	(b)	1	$(p \rightarrow q) \rightarrow q$	assumption
		2	$q \rightarrow p$	assumption
		3	$\neg p$	assumption
		4	$\neg q$	MT 2,3
		5	$\neg(p \rightarrow q)$	MT 1,4
		6	p	assumption
		7	\perp	\neg e 6,3
		8	q	\perp e 7
		9	$p \rightarrow q$	\rightarrow i 6-8
		10	\perp	\neg e 9,5
		11	p	PBC 3-10
		12	$(q \rightarrow p) \rightarrow p$	\rightarrow i 2-11
		13	$((p \rightarrow q) \rightarrow q) \rightarrow ((q \rightarrow p) \rightarrow p)$	\rightarrow i 1-12

(c)	1	$p \wedge q$	assumption
	2	$\neg p \vee \neg q$	assumption
	3	$\neg p$	assumption
	4	p	\wedge e ₁ 1
	5	\perp	\neg e 4,3
	6	$\neg q$	assumption
	7	q	\wedge e ₂ 1
	8	\perp	\neg e 7,6
	9	\perp	\vee e 2,3-5,6-8
	10	$\neg(\neg p \vee \neg q)$	\neg i 2-9
	11	$p \wedge q \rightarrow \neg(\neg p \vee \neg q)$	\rightarrow i 1-10

4 We have one introduction rule and two elimination rules for \leftrightarrow :

$$\frac{\begin{array}{|l} \varphi \\ \vdots \\ \psi \end{array}}{\varphi \leftrightarrow \psi} \leftrightarrow i \qquad \frac{\varphi \leftrightarrow \psi \quad \varphi}{\psi} \leftrightarrow e_1 \qquad \frac{\varphi \leftrightarrow \psi \quad \psi}{\varphi} \leftrightarrow e_2$$

By interpreting $\varphi \leftrightarrow \psi$ as $(\varphi \rightarrow \psi) \wedge (\psi \rightarrow \varphi)$, the above rules are derivable from the basic rules of natural deduction:

$$\leftrightarrow i: \quad \begin{array}{|l} \varphi \\ \vdots \\ \psi \end{array} \quad \text{given} \\ \hline \varphi \rightarrow \psi \quad \rightarrow i \\ \begin{array}{|l} \psi \\ \vdots \\ \varphi \end{array} \quad \text{given} \\ \hline \psi \rightarrow \varphi \quad \rightarrow i \\ (\varphi \rightarrow \psi) \wedge (\psi \rightarrow \varphi) \quad \wedge i$$

$$\leftrightarrow e_1: \quad \begin{array}{|l} (\varphi \rightarrow \psi) \wedge (\psi \rightarrow \varphi) \\ \varphi \\ \varphi \rightarrow \psi \\ \psi \end{array} \quad \begin{array}{l} \text{given} \\ \text{given} \\ \wedge e_1 \\ \rightarrow e \end{array} \qquad \leftrightarrow e_2: \quad \begin{array}{|l} (\varphi \rightarrow \psi) \wedge (\psi \rightarrow \varphi) \\ \psi \\ \psi \rightarrow \varphi \\ \varphi \end{array} \quad \begin{array}{l} \text{given} \\ \text{given} \\ \wedge e_2 \\ \rightarrow e \end{array}$$