

Selected Solutions

- 1 (b) We start by renaming the variables with multiple bindings. In this case we replace the second x by v and the second y by w :

$$\varphi \equiv (\forall x \exists y P(f(x, y), x) \wedge \neg \forall z Q(z)) \rightarrow \forall v \neg \forall w R(w, v)$$

Now we apply quantifier equivalences to pull the quantifiers outside:

$$\begin{aligned} &\equiv (\forall x \exists y P(f(x, y), x) \wedge \exists z \neg Q(z)) \rightarrow \forall v \exists w \neg R(w, v) \\ &\equiv \forall x \exists y (P(f(x, y), x) \wedge \exists z \neg Q(z)) \rightarrow \forall v \exists w \neg R(w, v) \\ &\equiv \forall x \exists y \exists z (P(f(x, y), x) \wedge \neg Q(z)) \rightarrow \forall v \exists w \neg R(w, v) \\ &\equiv \forall v \exists w (\forall x \exists y \exists z (P(f(x, y), x) \wedge \neg Q(z)) \rightarrow \neg R(w, v)) \\ &\equiv \forall v \exists w \exists x \forall y \forall z ((P(f(x, y), x) \wedge \neg Q(z)) \rightarrow \neg R(w, v)) \end{aligned}$$

At this point we have a prenex normal form. In the next step we transform the quantifier free part to CNF:

$$\begin{aligned} &\equiv \forall v \exists w \exists x \forall y \forall z (\neg(P(f(x, y), x) \wedge \neg Q(z)) \vee \neg R(w, v)) \\ &\equiv \forall v \exists w \exists x \forall y \forall z (\neg P(f(x, y), x) \vee Q(z) \vee \neg R(w, v)) \end{aligned}$$

In a final step we replace w and x by the unary Skolem function $g(v)$ and $h(v)$, respectively:

$$\approx \forall v \forall y \forall z (\neg P(f(h(v), y), h(v)) \vee Q(z) \vee \neg R(g(v), v))$$

- 2 (b) The clausal form is not satisfiable as seen by the following refutation:

- | | | | |
|----|----------------------------------|--------------|-----------------------------------|
| 1. | $\{P(x), P(f(a))\}$ | | |
| 2. | $\{\neg P(y), \neg Q(f(x), y)\}$ | | |
| 3. | $\{Q(f(a), a)\}$ | | |
| 4. | $\{\neg Q(x, y), Q(f(y), x)\}$ | | |
| 5. | $\{P(f(a))\}$ | factor 1 | $\{x \mapsto f(a)\}$ |
| 6. | $\{\neg Q(f(x), f(a))\}$ | resolve 2, 5 | $\{y \mapsto f(a)\}$ |
| 7. | $\{Q(f(a), f(a))\}$ | resolve 3, 4 | $\{x \mapsto f(a), y \mapsto a\}$ |
| 8. | \square | resolve 6, 7 | $\{x \mapsto a\}$ |

3 We first calculate the truth tables of f_1 , f_2 and f_3 :

x_1	x_2	x_3	s	f_1	f_2	f_3
0	0	0	0	0	0	0
0	0	1	1	0	0	1
0	1	0	1	0	1	0
0	1	1	2	1	1	1
1	0	0	1	1	0	0
1	0	1	2	0	0	0
1	1	0	2	1	1	1
1	1	1	3	1	1	1

From the table we obtain

$$\begin{aligned}
 f_1(x_1, x_2, x_3) &= \overline{x_1}x_2x_3 + x_1\overline{x_2}\overline{x_3} + x_1x_2 = \overline{x_1}x_2x_3 \oplus x_1\overline{x_2}\overline{x_3} \oplus x_1x_2 \\
 &= (x_1 \oplus 1)x_2x_3 \oplus x_1(x_2 \oplus 1)(x_3 \oplus 1) \oplus x_1x_2 \\
 &= x_1x_2x_3 \oplus x_2x_3 \oplus (x_1x_2 \oplus x_1)(x_3 \oplus 1) \oplus x_1x_2 \\
 &= x_1x_2x_3 \oplus x_2x_3 \oplus x_1x_2x_3 \oplus x_1x_2 \oplus x_1x_3 \oplus x_1 \oplus x_1x_2 \\
 &= x_1 \oplus x_1x_3 \oplus x_2x_3
 \end{aligned}$$

$$f_2(x_1, x_2, x_3) = x_2$$

$$f_3(x_1, x_2, x_3) = \overline{x_1}x_3 + x_1x_2 = \overline{x_1}x_3 \oplus x_1x_2 = (x_1 \oplus 1)x_3 \oplus x_1x_2 = x_1x_2 \oplus x_1x_3 \oplus x_3$$

4 Using Shannon's expansion we obtain

$$\begin{aligned}
 f(x, y, z) &= \overline{x}(\overline{y} \oplus y\overline{z}) \oplus x\overline{y}\overline{z} \\
 &= (x \oplus 1)((y \oplus 1) \oplus y(z \oplus 1)) \oplus x(y \oplus 1)(z \oplus 1) \\
 &= (x \oplus 1)(yz \oplus 1) \oplus x(yz \oplus y \oplus z \oplus 1) \\
 &= xyz \oplus x \oplus yz \oplus 1 \oplus xyz \oplus xy \oplus xz \oplus x \\
 &= xy \oplus xz \oplus yz \oplus 1
 \end{aligned}$$