

Logik

Week 9

May 23, 2024

## **Selected Solutions**

2 (a) Because  $f(0, 1, 1) = 0 \oplus (1 + 1) = 1$  and  $f(1, 1, 1) = 1 \oplus (1 + 1) = 0$ , f is not monotone. The dual

 $\hat{f}(x,y,z) = \overline{f(\overline{x},\overline{y},\overline{z})} = \overline{\overline{x} \oplus (\overline{y}+\overline{z})} = \overline{\overline{x} \oplus \overline{yz}} = x \oplus yz \oplus 1$ 

of f is different from f (e.g. f(0,0,0) = 0 and  $\hat{f}(0,0,0) = 1$ ), so f is not self-dual. Finally, f is not affine since its algebraic normal form is  $x \oplus y \oplus z \oplus yz$ .

|(b)| No. Since f(0,0,0) = 0, it follows from Post's adequacy theorem that  $\{f\}$  is not adequate. However,  $\{f, -\}$  is adequate and hence - cannot be expressed in terms of f.

Alternatively, one can prove by induction that any expression constructed from f and variables evaluates to 0 when all variables are set to 0. Since  $\overline{0} = 1$ ,  $\overline{\phantom{0}}$  cannot be expressed.

- 3 Yes. By Post's adequacy theorem,  $f(0, \ldots, 0) = 1$ ,  $f(1, \ldots, 1) = 0$ , and f is neither monotone nor self-dual nor affine. Let n be the arity of f. We prove that  $\{\hat{f}\}$  is adequate by showing that  $\hat{f}$  satisfies the five conditions of Post's adequacy theorem.
  - (1) We have  $\hat{f}(0, ..., 0) = \overline{f(1, ..., 1)} = \overline{0} = 1.$
  - (2) We have  $\hat{f}(1, ..., 1) = \overline{f(0, ..., 0)} = \overline{1} = 0.$
  - (3) Since f is not monotone,

$$f(b_1,\ldots,b_{i-1},x,b_{i+1},\ldots,b_n) = \overline{x}$$

for some i and  $b_1, \ldots, b_{i-1}, b_{i+1}, \ldots, b_n \in \{0, 1\}$ . Hence also

$$f(b_1,\ldots,b_{i-1},\overline{x},b_{i+1},\ldots,b_n) = x$$

and thus

$$\hat{f}(\overline{b}_1,\ldots,\overline{b}_{i-1},x,\overline{b}_{i+1},\ldots,\overline{b}_n) = \overline{f(b_1,\ldots,b_{i-1},\overline{x},b_{i+1},\ldots,b_n)} = \overline{x}$$

It follows that  $\hat{f}$  is not monotone.

Alternatively, the non-monotonicity of f is a direct consequence of (1) and (2).

- (4) Since the dual  $\hat{f}$  of  $\hat{f}$  is f, and  $\hat{f} \neq f$  because f is not self-dual,  $\hat{f}$  cannot be self-dual.
- (5) If  $\hat{f}$  is affine then there exist bits  $c, c_1, \ldots, c_n$  such that

$$\hat{f}(x_1,\ldots,x_n) = c \oplus c_1 x_1 \oplus \cdots \oplus c_n x_n$$

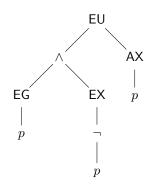
Hence

$$f(\overline{x}_1,\ldots,\overline{x}_n) = \overline{c \oplus c_1 x_1 \oplus \cdots \oplus c_n x_n} = 1 \oplus c \oplus c_1 x_1 \oplus \cdots \oplus c_n x_n$$

and thus

$$f(x_1, \ldots, x_n) = 1 \oplus c \oplus c_1 \overline{x_1} \oplus \cdots \oplus c_n \overline{x_n} = (1 \oplus c \oplus c_1 \oplus \cdots \oplus c_n) \oplus c_1 x_1 \oplus \cdots \oplus c_n x_n$$
  
contradicting the fact that  $f$  is not affine. We conclude that  $\hat{f}$  is not affine.

[4] (a) From the parse tree of  $\varphi$ 



we obtain 7 subformulas:

$$p \qquad \qquad \mathsf{EG}\,p \qquad \neg p \qquad \qquad \mathsf{EX}\,\neg p \qquad \qquad \mathsf{EG}\,p \wedge \mathsf{EX}\,\neg p \qquad \qquad \mathsf{AX}\,p \qquad \varphi$$

(b) Applying the CTL model checking algorithm results in the table

	p	$\neg p$	EGp	$EX\neg p$	$EGp\wedgeEX\neg p$	AXp	$\varphi$
1	$\checkmark$		$\checkmark$	$\checkmark$	$\checkmark$		$\checkmark$
2	$\checkmark$		$\checkmark$	$\checkmark$	$\checkmark$		$\checkmark$
3		$\checkmark$		$\checkmark$			
4	$\checkmark$		$\checkmark$			$\checkmark$	$\checkmark$

Hence  $\varphi$  holds in states 1, 2 and 4.