## Exercises.

- 9.0 Study Chapter 5
- 9.1 Exercise 5.3.4

Solution. According to Definition 5.3.6 a formula is true in a model provided its interpretation evaluates to true for all assignments. So it has to be shown that  $(X \equiv Y)^{\mathbf{I},\mathbf{A}} = \mathbf{t}$  iff  $(X^{\mathbf{I},\mathbf{A}} = Y^{\mathbf{I},\mathbf{A}})$ . But this follows from the fact that  $(X^{\mathbf{I},\mathbf{A}} \equiv Y^{\mathbf{I},\mathbf{A}}) = \mathbf{t}$  is equivalent as well to the first equation (by Definition 5.3.5) as to the second (by the definition of  $\equiv$ ).

## 9.2 Exercise 5.3.5

Solution. First we consider the "if" part: Suppose S is satisfiable in a model for L'. Then all formulas in S are true in some model  $\mathbf{M} = (\mathbf{B}, \mathbf{I})$  with assignment **A**. We define a model **N** for L by setting  $\mathbf{N} = (\mathbf{D}, \mathbf{I}|_{\mathsf{L}})$ , where  $\mathbf{I}|_{\mathsf{L}}$  denotes the restriction of the interpretation **I** to the language L. Set  $\mathbf{J} = \mathbf{I}|_{\mathsf{L}}$ . Now it is easy to show by induction that for each  $X \in S$ ,  $[X]^{\mathbf{J},\mathbf{A}} = \mathbf{t}$ . Hence S is satisfiable in a model for L.

Secondly, consider the "only if" part: Suppose S is satisfiable in a model for  $\mathsf{L}$ . Then all formulas in S are true in some model  $\mathbf{M} = (\mathbf{B}, \mathbf{I})$  with assignment  $\mathbf{A}$ . We define  $\mathbf{N}$  for  $\mathsf{L}'$  by setting  $\mathbf{N} = (\mathbf{J}, \mathbf{D})$ , where  $\mathbf{J}$  is chosen such that  $\mathbf{J}_{\mathsf{L}} = \mathbf{I}$ . (The symbols in  $\mathsf{L}' - \mathsf{L}$  can be interpreted arbitrarily.) Then it is again easy to show by induction that for each  $X \in S$ ,  $[X]^{\mathbf{J},\mathbf{A}} = \mathbf{t}$ . Thus S is satisfiable in a model for  $\mathsf{L}'$ .

9.3 Exercise 5.4.1

Solution. Let X be a formula of L, we have to show that  $[X]^{\mathbf{I},\mathbf{A}} = [X\mathbf{A}]^{\mathbf{I}}$  for all assignments **A**. To this avail, we exploit Prop. 5.3.7 in its simple form stating that the equality  $[X]^{\mathbf{I},\mathbf{A}} = [X\{x \mapsto \mathbf{A}(x)\}]^{\mathbf{I},\mathbf{A}}$  holds. As X is finite the number of

variables occurring free in X is finite, too. Suppose  $fvar(X) = \{x_1, \ldots, x_n\}$ . Then

$$[X]^{\mathbf{I},\mathbf{A}} = [X\{x_1 \mapsto \mathbf{A}(x_1)\}]^{\mathbf{I},\mathbf{A}}$$
Prop. 5.3.7  

$$= [X\{x_1 \mapsto \mathbf{A}(x_1)\}\{x_2 \mapsto \mathbf{A}(x_2)\}]^{\mathbf{I},\mathbf{A}}$$

$$\vdots$$

$$= [X\{x_1 \mapsto \mathbf{A}(x_1)\} \cdots \{x_n \mapsto \mathbf{A}(x_n)\}]^{\mathbf{I},\mathbf{A}}$$

$$= [X\mathbf{A}]^{\mathbf{I},\mathbf{A}}$$
as  $X\mathbf{A}$  is closed

Alternatively a direct proof of Proposition 5.4.3 would proceed by induction on X as follows:

- a) For the atomic cases this follows from Proposition 5.4.2.
- b) Case  $X = \neg X_1$ : By induction hypothesis (IH)  $X_1^{\mathbf{I},\mathbf{A}} = [X_1\mathbf{A}]^{\mathbf{I}}$  which implies  $\neg [X_1^{\mathbf{I},\mathbf{A}}] = \neg [[X_1\mathbf{A}]^{\mathbf{I}}]$ . The induction step follows as

$$[\neg X_1]^{\mathbf{I},\mathbf{A}} = \neg [X_1^{\mathbf{I},\mathbf{A}}] = \neg [[X_1\mathbf{A}]^{\mathbf{I}}] = [[\neg X_1]\mathbf{A}]^{\mathbf{I}}.$$

- c) Case  $X = (X_1 \circ X_2)$ : Analogous to b).
- d) Case  $X = (\forall x)X_1$ : By IH we know that  $[X_1]^{\mathbf{I},\mathbf{B}} = [X_1\mathbf{B}]^{\mathbf{I}}$ . Note that we can conceive the assignment **B** as a substitution (of closed terms) and thus obtain the equality:  $\mathbf{B} = (\mathbf{B}_x)\{x \mapsto B(x)\}$ .

Furthermore III is applicable to the formula  $X\mathbf{B}_x$ , hence we obtain  $[X\mathbf{B}_x]^{\mathbf{I},\mathbf{C}} = [(X\mathbf{B}_x)\mathbf{C}]^{\mathbf{I}}$  for any assignment  $\mathbf{C}$ . If we assume further that  $\mathbf{C}$  is defined such that  $\mathbf{C}(x) = \mathbf{B}(x)$ , and  $\mathbf{C}(y)$  is arbitrary for  $x \neq y$ , then we obtain  $[X\mathbf{B}_x]^{\mathbf{I},\mathbf{C}} = [(X\mathbf{B}]^{\mathbf{I}}, \text{ as } (X\mathbf{B}_x)\mathbf{C} = X\mathbf{B}.$ 

Now we show that  $[(\forall x)X]^{\mathbf{I},\mathbf{A}} = \mathbf{t}$  iff  $[(\forall x)X\mathbf{A}]^{\mathbf{I}} = \mathbf{t}$ :

 $[(\forall x)X]^{\mathbf{I},\mathbf{A}} = \mathbf{t}$  iff  $[X]^{\mathbf{I},\mathbf{B}} = \mathbf{t}$  for all x-variants **B** of **A** 

- iff  $[X\mathbf{B}]^{\mathbf{I}} = \mathbf{t}$  for all *x*-variants  $\mathbf{B}$  of  $\mathbf{A}$  (by IH)
- iff  $[X\mathbf{B}_x]^{\mathbf{I},\mathbf{C}} = \mathbf{t}$  for all assignments  $\mathbf{C}$ , s.t.  $\mathbf{C}(x) = \mathbf{B}(x)$  and  $\mathbf{C}$  arbitrary otherwise and all x-variants  $\mathbf{B}$  of  $\mathbf{A}$
- iff  $[(\forall x)X\mathbf{B}_x]^{\mathbf{I},\mathbf{D}} = \mathbf{t}$  for all assignments  $\mathbf{D}$
- iff  $[(\forall x)X\mathbf{B}_x]^{\mathbf{I}} = \mathbf{t} ((\forall x)X\mathbf{B}_x \text{ is closed})$
- iff  $[((\forall x)X)\mathbf{A}]^{\mathbf{I}} = \mathbf{t}$  as **B** is an *x*-variant of **A**.

e) Analogous to d).

9.4 Exercise 5.4.2

Solution. We only consider the first case. Suppose  $\mathbf{M} = (\mathbf{D}, \mathbf{I})$ .

 $(\forall x)X$  is true in **M** iff  $[(\forall x)X]^{\mathbf{I},\mathbf{A}} = \mathbf{t}$  for all assignments **A** 

- iff  $[X]^{\mathbf{I},\mathbf{B}} = \mathbf{t}$  for all *x*-variants **B** of **A**
- iff  $[(X \{x \mapsto \mathbf{B}(x)\})]^{\mathbf{I},\mathbf{A}} = \mathbf{t}$  for all x-variants **B** of **A** as **B** is an x-variant of **A** by the general statement of Prop. 5.3.7
- iff  $[(X\{x \mapsto d\})]^{\mathbf{I},\mathbf{A}} = \mathbf{t}$  for all  $d \in \mathbf{D}$ .