

Functional Programming

1

(a)

explanation

$$\begin{aligned}
 (\lambda xyz.x z (y z)) (\lambda xy.x) (\lambda x.x) (\lambda x.x) &\rightarrow_{\beta} (\lambda yz.(\lambda xy.x) z (y z)) (\lambda x.x) (\lambda x.x) \\
 &\rightarrow_{\beta} (\lambda yz.(\lambda y.z) (y z)) (\lambda x.x) (\lambda x.x) \\
 &\rightarrow_{\beta} (\lambda yz.z) (\lambda x.x) (\lambda x.x) \\
 &\rightarrow_{\beta} (\lambda z.z) (\lambda x.x) \\
 &\rightarrow_{\beta} \lambda x.x
 \end{aligned}$$

(b)

explanation

$$\begin{aligned}
 (\lambda xyz.x z (y z)) (\lambda xy.x) (\lambda x.x) (\lambda x.x) &\rightarrow_{\beta} (\lambda yz.(\lambda xy.x) z (y z)) (\lambda x.x) (\lambda x.x) \\
 &\rightarrow_{\beta} (\lambda z.(\lambda xy.x) z ((\lambda x.x) z)) (\lambda x.x) \\
 &\rightarrow_{\beta} (\lambda xy.x) (\lambda x.x) ((\lambda x.x) (\lambda x.x)) \\
 &\rightarrow_{\beta} (\lambda y.(\lambda x.x)) ((\lambda x.x) (\lambda x.x)) \\
 &\rightarrow_{\beta} \lambda x.x
 \end{aligned}$$

2

(a)

base case

Let $t = E$. The base case concludes by the derivation

$$\begin{aligned}
 \text{sum (preorder } t) &= \text{sum (preorder } E) && \text{(since } t = E) \\
 &= \text{sum (} [] \text{)} && \text{(by definition of preorder)} \\
 &= 0 && \text{(by definition of sum)} \\
 &= \text{sum_tree } E && \text{(by definition of sum_tree)} \\
 &= \text{sum_tree } t
 \end{aligned}$$

(b)

step case

Let $t = N (t_1, v, t_2)$. By IH it holds that

$$\begin{aligned}
 \text{sum (preorder } t_1) &= \text{sum_tree } t_1 \\
 \text{sum (preorder } t_2) &= \text{sum_tree } t_2.
 \end{aligned}$$

The step case concludes by the derivation

$$\begin{aligned}
 \text{sum (preorder } t) &= \text{sum (preorder (N (} t_1, v, t_2 \text{)))} && (t = N (t_1, v, t_2)) \\
 &= \text{sum (} v :: (\text{preorder } t_1 @ \text{preorder } t_2) \text{)} && \text{(def. of preorder)} \\
 &= v + (\text{sum (preorder } t_1 @ \text{preorder } t_2) \text{)} && \text{(def. of sum)} \\
 &= v + (\text{sum (preorder } t_1) + \text{sum (preorder } t_2) \text{)} && \text{(by } \star \text{)} \\
 &= v + (\text{sum_tree } t_1 + \text{sum_tree } t_2) && \text{(by IH)} \\
 &= \text{sum_tree (N (} t_1, v, t_2 \text{))} && \text{(def. of sum_tree)} \\
 &= \text{sum_tree } t
 \end{aligned}$$

3(a) *implementation*

```

let unique xs =
  let rec rev acc = function
    | [] -> acc
    | x :: xs -> rev (x :: acc) xs
  in
  let rec unique acc = function
    | [] -> rev [] acc
    | x :: xs ->
      if mem x xs then unique acc xs else unique (x :: acc) xs
  in
  unique [] xs
;;

```

(b) *implementation*

```

let percentage x ys =
  let rec p x = function
    | [] -> (0, 0)
    | y :: ys ->
      let (i, j) = p x ys in
      if x = y then (i + 1, j + 1) else (i, j + 1)
  in
  if ys = [] then 0.0 else let (i, j) = p x ys in
  float_of_int i /. float_of_int j
;;

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4(a) *explanation*

$$t \rightarrow_{\beta} y (\lambda yz.z y) w$$
(b) *explanation*

$$\mathcal{FVar}(t) = \{w, y\}$$
(c) *explanation*

$$\mathcal{BVar}(t) = \{x, y, z\}$$
(d) *explanation*

$$\text{Sub}(t) = \{y, x, y x, \lambda x.y x, z, z y, \lambda z.z y, \lambda yz.z y, (\lambda x.y x) (\lambda yz.z y), w, t\}$$

5

(a) *explanation*

$$\frac{\frac{E, x : \text{int} \vdash \mathbf{p} : \text{int} \rightarrow \text{int} \rightarrow \text{pair}(\text{int}, \text{int}) \quad E, x : \text{int} \vdash x : \text{int}}{E, x : \text{int} \vdash \mathbf{p} \ x : \text{int} \rightarrow \text{pair}(\text{int}, \text{int})} \text{(app)} \quad \star}{\frac{E \vdash 1 : \text{int} \quad E, x : \text{int} \vdash \mathbf{p} \ x (x + x) : \text{pair}(\text{int}, \text{int})}{E \vdash \mathbf{let} \ x = 1 \ \mathbf{in} \ \mathbf{p} \ x (x + x) : \text{pair}(\text{int}, \text{int})} \text{(let)}} \text{(app)}$$

and \star is

$$\frac{\frac{E, x : \text{int} \vdash (+) : \text{int} \rightarrow \text{int} \rightarrow \text{int} \quad E, x : \text{int} \vdash x : \text{int}}{E, x : \text{int} \vdash (+) \ x : \text{int} \rightarrow \text{int}} \text{(app)} \quad E, x : \text{int} \vdash x : \text{int}}{E, x : \text{int} \vdash (x + x) : \text{int}} \text{(app)}$$

(b) *explanation*

$$\begin{aligned} & E \triangleright \mathbf{let} \ x = 1 \ \mathbf{in} \ \mathbf{p} \ x (x + x) : \alpha_0 \\ & \quad \xRightarrow{\text{let}} \\ & E \triangleright 1 : \alpha_1; E, x : \alpha_1 \triangleright \mathbf{p} \ x (x + x) : \alpha_0 \\ & \quad \xRightarrow{\text{con}} \\ & \text{int} \approx \alpha_1; E, x : \alpha_1 \triangleright \mathbf{p} \ x (x + x) : \alpha_0 \\ & \quad \xRightarrow{\text{app}} \\ & \text{int} \approx \alpha_1; E, x : \alpha_1 \triangleright \mathbf{p} \ x : \alpha_2 \rightarrow \alpha_0; E, x : \alpha_1 \triangleright (x + x) : \alpha_2 \\ & \quad \xRightarrow{\text{app}} \\ & \text{int} \approx \alpha_1; E, x : \alpha_1 \triangleright \mathbf{p} : \alpha_3 \rightarrow \alpha_2 \rightarrow \alpha_0; E, x : \alpha_1 \triangleright x : \alpha_3; E, x : \alpha_1 \triangleright (x + x) : \alpha_2 \\ & \quad \xRightarrow{\text{con}} \\ & \text{int} \approx \alpha_1; \text{int} \rightarrow \text{int} \rightarrow \text{pair}(\text{int}, \text{int}) \approx \alpha_3 \rightarrow \alpha_2 \rightarrow \alpha_0; \\ & \quad E, x : \alpha_1 \triangleright x : \alpha_3; E, x : \alpha_1 \triangleright (x + x) : \alpha_2 \\ & \quad \xRightarrow{\text{con}} \\ & \text{int} \approx \alpha_1; \text{int} \rightarrow \text{int} \rightarrow \text{pair}(\text{int}, \text{int}) \approx \alpha_3 \rightarrow \alpha_2 \rightarrow \alpha_0; \\ & \quad \alpha_1 \approx \alpha_3; E, x : \alpha_1 \triangleright (x + x) : \alpha_2 \\ & \quad \xRightarrow{\text{app}} \\ & \text{int} \approx \alpha_1; \text{int} \rightarrow \text{int} \rightarrow \text{pair}(\text{int}, \text{int}) \approx \alpha_3 \rightarrow \alpha_2 \rightarrow \alpha_0; \\ & \quad \alpha_1 \approx \alpha_3; E, x : \alpha_1 \triangleright (+) \ x : \alpha_4 \rightarrow \alpha_2; E, x : \alpha_1 \triangleright x : \alpha_4 \\ & \quad \xRightarrow{\text{app}} \\ & \text{int} \approx \alpha_1; \text{int} \rightarrow \text{int} \rightarrow \text{pair}(\text{int}, \text{int}) \approx \alpha_3 \rightarrow \alpha_2 \rightarrow \alpha_0; \\ & \quad \alpha_1 \approx \alpha_3; E, x : \alpha_1 \triangleright (+) : \alpha_5 \rightarrow \alpha_4 \rightarrow \alpha_2; E, x : \alpha_1 \triangleright x : \alpha_5; E, x : \alpha_1 \triangleright x : \alpha_4 \\ & \quad \xRightarrow{\text{con}} \\ & \text{int} \approx \alpha_1; \text{int} \rightarrow \text{int} \rightarrow \text{pair}(\text{int}, \text{int}) \approx \alpha_3 \rightarrow \alpha_2 \rightarrow \alpha_0; \\ & \quad \alpha_1 \approx \alpha_3; \text{int} \rightarrow \text{int} \rightarrow \text{int} \approx \alpha_5 \rightarrow \alpha_4 \rightarrow \alpha_2; E, x : \alpha_1 \triangleright x : \alpha_5; E, x : \alpha_1 \triangleright x : \alpha_4 \\ & \quad \xRightarrow{\text{con}} \\ & \text{int} \approx \alpha_1; \text{int} \rightarrow \text{int} \rightarrow \text{pair}(\text{int}, \text{int}) \approx \alpha_3 \rightarrow \alpha_2 \rightarrow \alpha_0; \\ & \quad \alpha_1 \approx \alpha_3; \text{int} \rightarrow \text{int} \rightarrow \text{int} \approx \alpha_5 \rightarrow \alpha_4 \rightarrow \alpha_2; \alpha_1 \approx \alpha_5; E, x : \alpha_1 \triangleright x : \alpha_4 \\ & \quad \xRightarrow{\text{con}} \\ & \text{int} \approx \alpha_1; \text{int} \rightarrow \text{int} \rightarrow \text{pair}(\text{int}, \text{int}) \approx \alpha_3 \rightarrow \alpha_2 \rightarrow \alpha_0; \\ & \quad \alpha_1 \approx \alpha_3; \text{int} \rightarrow \text{int} \rightarrow \text{int} \approx \alpha_5 \rightarrow \alpha_4 \rightarrow \alpha_2; \alpha_1 \approx \alpha_5; \alpha_1 \approx \alpha_4 \end{aligned}$$