

Name: \_\_\_\_\_

Points: \_\_\_\_\_

Matr.-Nr.: \_\_\_\_\_

## 1 Formalisation (15 pts)

Consider the following statements:

- ① If a car has a diesel engine and is refueled with gas then it gets defect.
- ② A mechanic cannot repair a defect car if she/he does not have the necessary tool for that car.
- ③ There is a blue car which is refueled with gas.
- ④ All defect cars are blue and all blue cars have diesel engine.
- ⑤ For every car with diesel engine there is a mechanic which has the necessary tool for that car.
- ⑥ There is a defect car which cannot be repaired by any mechanic.

a) Give a first-order formula that formalises the statements above. Use therefore the following constants, functions and predicates: [5]

- constants: *diesel*, *gas* and functions: *tool(x)*  
[ first letter suffices, hence you can use *d, g, t(x)* ]
- predicates: *car(x)*, *engine(x, y)*, *refuel(x, y)*, *mechanic(x)*, *repair(x, y)*,  
*defect(x)*, *has(x, y)*, *blue(x)* [first 3 letters suffice to identify the predicate]

Note that the function *tool(x)* returns the tool needed to repair car *x*. Further, the predicate *engine(x, y)* has to be interpreted as “*x* has a *y* engine”, the predicate *refuel(x, y)* as “*x* is refueled with *y*”, the predicate *repair(x, y)* as “*x* can repair *y*”, the predicate *has(x, y)* as “*x* has *y*”.

b) Is your formalization satisfiable? Explain your answer. [5]

c) Change ② so that the answer to b) changes and shortly explain why it changes. [5]

## 2 Resolution (10 pts)

a) Decide for the following set of clauses

$$\begin{array}{cccc} \{\neg p, q, \neg r\} & \{q, \neg r, \neg s\} & \{p, q, \neg s\} & \{\neg p, \neg q, \neg s\} \\ \{p, \neg q, \neg s\} & \{\neg q, r, s\} & \{q, r, s\} & \{\neg r, s\} \end{array}$$

whether the empty clause can be derived or not. If the empty clause can be derived, prove it by resolution. Otherwise give a satisfying truth assignment (valuation). [4]

b) Prove by resolution that:

$$\begin{aligned} & \forall x, y (P(x, f(x), f(y)) \rightarrow Q(y, x)) \quad \wedge \\ & \forall x, y (Q(x, y) \rightarrow Q(g(x), f(y))) \quad \wedge \\ & \forall x, y, z P(x, y, f(g(z))) \quad \vdash \quad \forall x Q(g(g(g(x))), f(f(x))) \end{aligned}$$

Display the resolution tree and give all substitutions employed in its construction. [6]

### 3 Resolution from semantic tree (5pt)

Consider the following set of clauses

$$\{p, \neg q, r\} \quad \{\neg p, \neg q\} \quad \{p, q, r\} \quad \{q\} \quad \{p, \neg q, \neg r\}$$

a) Give the semantic tree. Mark the failure nodes and indicate their associated clauses. [2.5]

b) Construct a resolution derivation of the empty clause in the following way:

- ① In each step choose an inference node from the semantic tree from a)
- ② Resolve associated clauses of the two children of the inference node.
- ③ Mark new failure nodes in the semantic tree and indicate associated clauses.
- ④ Repeat until no more inference node. [2.5]

### 4 Theoretical Questions (10 pts)

Give the (correct) answers to the following questions:

a) Explain the difference between  $\models$  and  $\vdash$  with respect to propositional logic. How are  $\models$  and  $\vdash$  *related* to each other? [3]

b) Is it *decidable* whether a propositional formula is valid or not? Cross the right answer!

[1]  
 Yes No Unknown

c) Give the definition of *Herbrand expansion*, state and prove its associated lemma. Outline the *Ground-Resolution* procedure. Give the proof of *completeness* of this procedure. In particular, name and state the main theorems used in this short proof (no need to prove the main theorems). [5]

d) Is it *decidable* whether a first-order formula is valid or not? Cross the right answer!

[1]  
 Yes No Unknown