Computational Logic
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21 November 2008
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Week 8 - Efficien

## Induction on Lists

Induction Principle (without Types)

$$
(\underbrace{P([])}_{\text {base case }} \wedge \underbrace{\forall x . \forall x s .(P(x s) \rightarrow P(x:: x s))}_{\text {step case }}) \rightarrow \forall / s . P(I s)
$$

Lemma
@ is associative, i.e.,

$$
x s @(y s @ z s)=(x s @ y s) @ z s
$$

Proof.
Blackboard

## Mathematical Induction

$$
\begin{aligned}
& \text { Induction Principle } \\
& (\underbrace{P(m)}_{\text {base case }} \wedge \underbrace{\forall k \geq m \cdot(P(k) \rightarrow P(k+1))}_{\text {step case }}) \rightarrow \forall n \geq m \cdot P(n) \\
& \text { Example }
\end{aligned}
$$

- first domino will fall
- if a domino falls also its right neighbor falls

Summary of Week 7


## Structural Induction

## Usage

- can be used on every variant type
- base cases correspond to non-recursive constructors
- step cases correspond to recursive constructors

Example

- lists
- trees
- $\lambda$-terms
- 


## This Week

Practice I
OCaml introduction, lists, strings, trees
Theory I
lambda-calculus, evaluation strategies, induction,
reasoning about functional programs
Practice II
efficiency, tail-recursion, combinator-parsing
Theory II
type checking, type inference
Advanced Topics
lazy evaluation, infinite data structures, monads, ...

## (ICS@UIBK)

## Mathematical (cont'd)

Mathematical
Definition ( $n$-th Fibonacci number)

$$
\begin{aligned}
& \text { fibn } \stackrel{\text { def }}{=} \begin{cases}1 & \text { if } n \leq 1 \\
\text { fib }(n-1)+\operatorname{fib}(n-2) & \text { otherwise }\end{cases} \\
& \text { Graph } \\
& \text { fib(n) }
\end{aligned}
$$

(ICS@UIBK) FP

## OCaml

Definition
let rec fib $n=$ if $n<2$ then 1 else $f i b(n-1)+f i b(n-2)$
Example

$1,1,2,3,5,8,13,21,34,55,89,144,233,377,610,987,1597$, 2584, 4181 ,6765, 10946, 17711, 28657, 46368, 75025, 121393, 196418, 317811, 514229, 832040, 1346269, 2178309, 3524578, 5702887, 9227465, 14930352, 24157817, 39088169, 63245986, 102334155, 165580141, 267914296, 433494437, 701408733, 1134903170, 1836311903, 2971215073, ...

## Combining Several Results

## Idea

- use tuples to return more than one result
- make results available as return values instead of recomputing them

Goal
compute average value of an integer list

## Approach 1

- let average xs = IntLst.sum xs / Lst.length xs
- 2 traversals of xs are done


## Combined Function

- 

```
let rec sumlen = function
    | [] -> (0,0)
    | x::xs -> let (s,l) = sumlen xs in ( }\textrm{x}+\textrm{s},\textrm{l}+1
```

- one traversal of xs suffices

Fibonacci Numbers

Example
let rec fibpair $n=$ if $n<1$ then $(0,1)$ else if $\mathrm{n}=1$ then (1,1)
else let (f1,f2) = fibpair ( $\mathrm{n}-1$ ) in (f2,f1+f2)
)

- this function is linear

Lemma

$$
\operatorname{fibpair}(n+1)=(\text { fib } n, \operatorname{fib}(n+1))
$$

Proof.
Blackboard

## Recursion vs. Tail Recursion

Idea

- a function calling itself is recursive
- functions that mutually call each other are mutually recursive
- special kind of recursion is tail recursion

Definition (Tail recursion)
a function is called tail recursive if the last action in the function body is the recursive call

## Even/Odd

Length

```
let rec length = function [] -> 0
                | x::xs -> 1 + length xs
```

- not tail recursive


## Parameter Accumulation

## Idea

- make function tail recursive
- provide data as input instead of computing it before recursive call
- Why? (tail recursive functions can automatically be transformed into space-efficient loops)


## Average

- 

let sumlen xs =
let rec sumlen sum len = function
| [] -> (sum,len)
| x::xs $->$ sumlen ( $x+$ sum) (len+1) xs
in
sumlen 0 0 xs

- tail recursive

