

First name: _____

Last name: _____

Matriculation number: _____

- Write your name and matriculation number on every page.
- Please answer all exercises in a readable and precise way. Do *not* write with a pencil or a red pen. Please cross out solution attempts which are replaced by another solution.
- Cheating is not allowed. Everyone who is caught will fail the exam.
- Please do not remove the staples of the exam.

Exercise	Maximal points	Points
1	12	
2	24	
3	15	
4	19	
Σ	70	
Grade		

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Exercise 1 (12 points)

Each correct answer is worth four points. A wrong answer results in zero points. Giving no answer is worth one point.

	Yes	No
The CTL formula $(\mathbf{A} \mathbf{G} \mathbf{A} \mathbf{F} \textit{request}) \Rightarrow (\mathbf{A} \mathbf{G} \mathbf{A} \mathbf{F} \textit{response})$ is equivalent to the LTL formula $(\mathbf{G} \mathbf{F} \textit{request}) \Rightarrow (\mathbf{G} \mathbf{F} \textit{response})$.		
Every language $L \subseteq \Sigma^\omega$ can be recognized by some NBA.		
Emptiness of $\mathcal{L}(\mathcal{A})$ for some GNBA \mathcal{A} can directly be decided using an SCC-based analysis, without first translating \mathcal{A} into some NBA.		

Exercise 2 (21 + 3 points)

Consider the following nanoPromela program which has two clients ($i \in \{1, 2\}$) which send their data via a scheduler to a printer. After a clients data d_i is delivered at the printer, client i gets an acknowledgement.

```

----- CLIENT  $i$  -----
do :: true => ic !  $i$ ; dc !  $d_i$ ; ac $_i$  ? ack od

----- SCHEDULER -----
atomic { x := 0; d := "" };
do :: true => ic ? x; dc ? d; pc ! d; if :: x = 1 => ac $_1$  ! ack :: x = 2 => ac $_2$  ! ack fi od

----- PRINTER -----
do :: true => pc ? d; skip od

```

- Construct the channel-system for the nanoPromela program.

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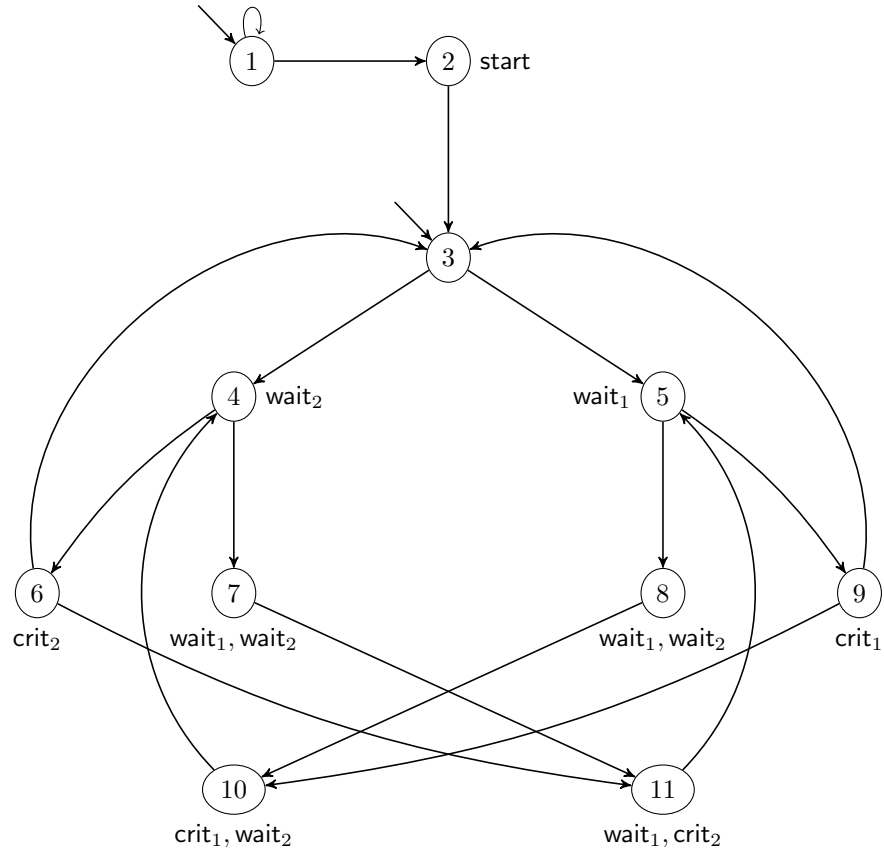
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- Does the program contain a serious bug using asynchronous communication? If so, shortly describe it.

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Exercise 3 (15 points)



Consider the above transition system TS of a mutual exclusion protocol and the following CTL*-formula Φ .

$$\Phi = (A((FG \neg \text{start}) \wedge A(\neg \text{wait}_1 \vee F \text{crit}_1))) \wedge AF(\text{crit}_1 \vee \text{crit}_2)$$

Does $TS \models \Phi$ hold? Justify your answer by performing CTL*-model checking, and write down $Sat(\Psi)$ for every state-subformula Ψ of Φ . Whenever one computes a set $Sat(A\varphi)$, additionally write down the corresponding LTL-formula φ' that is checked. However, it is not necessary to perform LTL-model checking explicitly.

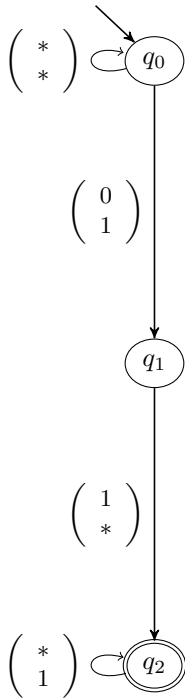
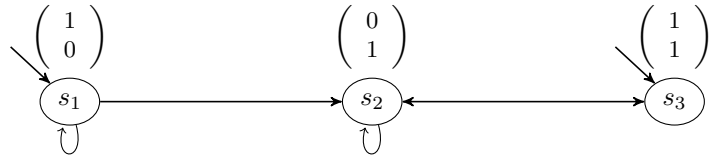
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Exercise 4 (18 + 1 points)

Consider the following NBA \mathcal{A} and the following transition system TS .



- Construct the NBA $\mathcal{B} = TS \otimes \mathcal{A}$ which accepts $\mathcal{L}(TS) \cap \mathcal{L}(\mathcal{A})$.
- Is $\mathcal{L}(\mathcal{B}) = \emptyset$? If not, then provide a word which is contained in $\mathcal{L}(\mathcal{B})$.