First Exam Game Theory, LVA 703716

March 15, 2010

Name:

Student number:

The exam consists of 5 exercises with a total of 50 points.



(5 pts)

(7 pts)

- 1. a) Show that transitivity (of \succeq_S) implies (i) if $f \sim_S g, g \sim_S h$, then $g \sim_S h$, and (ii) if $f \succ_S g, g \succeq_S h$, then $g \succ_S h$.
 - b) Consider $u: X \times \Omega \to \mathbb{R}, y \in X$, and the set S of all probabilities $p \in \Delta(\Omega)$ such that y is optimal. Prove or disprove: The set S is convex. (6 pts)
- 2. Consider the following game. Two hunter hunt together for dear or rabbits. If they hunt down a dear they get 4 points, if they hunt down a rabbit they get 2 points. While they can hunt rabbits on their own, they can only hunt a dear if they work together.
 - a) Express the game in extensive form. (5 pts)
 - b) Transform the game to fully reduced strategic form. (5 pts)
- 3. Consider the following two games:

	ma	0.10	~~		x_2	y_2	~ 2	
	x_2	g_2	~ 2	r_1	0 0	5 4	15	
r_1	0 4	56	87	x_1	0, 0	$0, \mathbf{T}$	$\mathbf{I}, 0$	
w_1	$0, \pm$	0, 0	0,1	2/1	4 5	0.0	$5\ 4$	
1/1	2.9	6.5	5.1	91	1,0	0, 0	0, 1	
91	-,0	\circ, \circ	0,1	z_1	5.4	4.5	0.0	
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- a) Compute all Nash equilibria of the game Γ_1 to the left. (6 pts)
- b) Compute all Nash equilibria of the game Γ_2 to the right.

For the computation use the support enumeration algorithm and note that the number of equilibria is odd for both games.

4. Sketch the vertex enumeration algorithm and describe all necessary assumptions for its invocation.

С	prrect answer is worth 1 points (and every wrong -1 points).		
	statement	yes	no
	consistently in pursuit of her own objective.		
	A lottery is a function from states to the probability distribution over a set of prizes. If the lottery is independent on the states it depends only on subjective unknowns.		
	A set of vectors S is convex if for any two vectors p, q also $\lambda p + (1 - \lambda)q \in S$, where $\lambda \in [0, 1]$.		
	Given a finite game Γ in strategic form, there exists at least one pure equilibrium.		
	An auction where the bidders have the same private information is called common value auction.		
	Nash's theorem of the existence of an equilibrium is not extensible to games over infinite strategy sets		
	A game may have multiple equilibria, but at least one of the equilibria is efficient.		
	Let $m, n \in \mathbb{N}$ and $m < n$. A two-person game is called degener- ated if there exists a strategy profile σ with support size m such that σ has n pure best responses.		
	For a Nash equilibrium (σ, ρ) of a degenerated two-person game, σ and ρ have support of equal size.		
	If we can show that $P = NP$, then $P = PPAD$ follows.		

5. Determine whether the statements on the answer sheet are true or false. Every

(10 pts)