

# Game Theory

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# Summary of Last Lecture

## Assumption

decision makers/players act rational and intelligent

Definition lottery

• a lottery f maps  $\Omega$  to the probability distributions  $\Delta(X)$ 

$$\sum_{x \in X} f(x|t) = 1 \qquad t \in \Omega$$

• the set of lotteries is defined as follows:

$$L = \{f \mid f : \Omega \to \Delta(X)\}$$

• let t be a state,  $f(\cdot|t)$  denotes the probability distribution over X in t:

$$f(\cdot|t) = (f(x|t))_{x \in X} \in \Delta(X)$$

• the lottery [x] always get prize  $x \in X$  for sure:

$$[x](y|t) = \begin{cases} 1 & \text{if } y = x \\ 0 & \text{if } y \neq x \end{cases}$$

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## Axiomatic Presentation ①

## Axiom (totality)

- $f \succcurlyeq_S g$  or  $g \succcurlyeq_S f$
- if  $f \succcurlyeq_S g$  and  $g \succcurlyeq_S h$ , then  $f \succcurlyeq_S h$

### Axiom (relevance)

 $\forall t \in S$ :  $f(\cdot|t) = g(\cdot|t)$ , then  $f \sim_S g$ 

### Axiom (monotonicity)

if  $f \succ_S h$  and  $\alpha > \beta$ , then  $\alpha f + (1 - \alpha)h \succ_S \beta f + (1 - \beta)h$ 

### Axiom (continuity)

if  $f \succcurlyeq_S g$  and  $g \succcurlyeq_S h$ , then  $\exists \gamma g \sim_S \gamma f + (1 - \gamma)h$ 

## Axiom (interest)

 $\forall t \in \Omega, \exists x, y \in X [y] \succ_{\{t\}} [x]$ 

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# Axiomatic Presentation 2

## Axiom (objective substitution)

if  $e \succcurlyeq_S f$  and  $g \succcurlyeq_S h$ , then  $\alpha e + (1 - \alpha)g \succcurlyeq_S \alpha f + (1 - \alpha)h$ 

## Axiom (strict objective substitution)

if  $e \succ_S f$  and  $g \succcurlyeq_S h$ , then  $\alpha e + (1 - \alpha)g \succ_S \alpha f + (1 - \alpha)h$ 

## Axiom (subjective substitution)

if  $f \succcurlyeq_S g$  and  $f \succcurlyeq_T g$  and  $S \cap T = \emptyset$ , then  $f \succcurlyeq_{S \cup T} g$ 

## Axiom (strict subjective substitution)

if  $f \succ_S g$  and  $f \succ_T g$  and  $S \cap T = \emptyset$ , then  $f \succ_{S \cup T} g$ 

### Axiom (state neutrality)

optional

 $\forall r, t \in \Omega, \ f(\cdot|r) = f(\cdot|t), \ g(\cdot|r) = g(\cdot|t), \ \text{then} \ (f \succcurlyeq_{\{r\}} g) \to (f \succcurlyeq_{\{t\}} g)$ 

# **Expected Utility**

recall:  $\Xi = \{S \mid S \subseteq \Omega, S \neq \emptyset\}$  denotes the set of all events

#### Definition

• a conditional-probability function  $p: \Xi \to \Delta(\Omega)$ :

$$p(t|S) = 0$$
 if  $t \notin S$  
$$\sum_{r \in S} p(r|S) = 1$$

- $p(R|S) = \sum_{t \in R} p(t|S)$
- a utility function is any function from  $u: X \times \Omega \to \mathbb{R}$
- a utility function u is state independent if  $\exists U \colon X \to \mathbb{R}$  such that u(x,t) = U(x) for all  $x \in X$ ,  $t \in \Omega$

#### **Definition**

let p denote a conditional-probability function and u any utility function, then the expected utility determined by lottery f is defined as:

$$\underline{E_p}(u(f)|S) = \sum_{t \in S} p(t|S) \sum_{x \in X} u(x,t) f(x|t)$$

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### Content

motivation, introduction to decision theory, decision theory

basic model of game theory, dominated strategies, Bayesian games, Nash equilibrium

two-person zero-sum games, Bayesian equilibrium, sequential equilibria of extensive-form games, computing Nash equilibria, sub-game-perfect equilibria

efficient computation of Nash equilibria, complexity class PPAD, complexity of Nash equilibria, refinements of equilibrium in strategic form, persistent equilibria, games with communication, sender-receiver games

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# **Expected Utility Maximisation Theorem**

#### **Theorem**

the axioms (with state neutrality) are satisfied if and only if there exists a (state-independent) utility function u and a conditional-probability function p such that

- 1  $\max_{x \in X} u(x, t) = 1$  and  $\min_{x \in X} u(x, t) = 0$
- 2  $p(R|T) = p(R|S)p(S|T) \ \forall \ R, S, T \text{ so that } R \subseteq S \subseteq T \text{ and } S \neq \emptyset$
- **3**  $f \succcurlyeq_S g$  if and only if  $E_p(u(f)|S) \geqslant E_p(u(g)|S)$

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### Proof of the Theorem ①

### **Special Lotteries**

• define  $a_1$  for all  $t \in \Omega$ :

$$a_1(y|t) = 1$$
 such that  $\forall x \in X$ ,  $[y] \succcurlyeq_{\{t\}} [x]$ 

• define  $a_0$  for all  $t \in \Omega$ :

$$a_0(y|t) = 1$$
 such that  $\forall x \in X$ ,  $[x] \succcurlyeq_{\{t\}} [y]$ 

#### More Special Lotteries

for every event  $S \in \Xi$ 

$$egin{aligned} m{b}_{\mathcal{S}}(\cdot|t) &= egin{cases} a_1(\cdot|t) & ext{if } t \in \mathcal{S} \ a_0(\cdot|t) & ext{if } t 
ot\in \mathcal{S} \end{cases} \end{aligned}$$

## Proof of the Theorem 2

## Definition u(x, t)

• ask: what is the correct number  $\beta$  such that

$$[x] \sim_{\{t\}} \beta a_1 + (1 - \beta) a_0$$

• set  $u(x, t) = \beta$ 

## Definition p(t|S)

• ask: what is the correct number  $\gamma$  such that

$$b_{\{t\}} \sim_{\mathcal{S}} \gamma a_1 + (1 - \gamma) a_0$$

• set  $p(t|S) = \gamma$ 

#### **Proof Plan**

show that u(x, t) and p(t|S) fulfil the 3rd condition:

$$f \succcurlyeq_S g \Leftrightarrow E_p(u(f)|S) \geqslant E_p(u(g)|S)$$

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## Proof of the Theorem 3

Claim 
$$\forall r \in \Omega: \ c_{x,t} \sim_{\{r\}} u(x,t)b_{\{t\}} + (1 - u(x,t))a_0$$

#### **Proof**

on blackboard

Claim 
$$\forall S \in \Xi$$
:  $c_{x,t} \sim_S u(x,t)b_{\{t\}} + (1-u(x,t))a_0$ 

### Proof

from Claim 1 with Subjective Substitution

Claim

$$f \succcurlyeq_S g \Leftrightarrow \frac{1}{n}f + (1 - \frac{1}{n})a_0 \succcurlyeq_S \frac{1}{n}g + (1 - \frac{1}{n})a_0$$

where  $n = |\Omega|$ 

#### **Proof**

with Objective Substitution

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## Proof of the Theorem 4

Claim

$$\frac{1}{n}f + (1 - \frac{1}{n})a_0 = \frac{1}{n}\sum_{t \in \Omega}\sum_{x \in X}f(x|t)c_{x,t}$$

#### **Proof**

on blackboard

using the definitions of u(x, t), p(t|S) we have

$$\frac{1}{n} \sum_{t \in \Omega} \sum_{x \in X} f(x|t) c_{x,t} 
\sim_{S} \frac{1}{n} \sum_{t \in \Omega} \sum_{x \in X} f(x|t) \left( u(x,t) b_{\{t\}} + (1 - u(x,t)) a_{0} \right) 
\sim_{S} \frac{1}{n} \sum_{t \in \Omega} \sum_{x \in X} f(x|t) \left( u(x,t) \left[ p(t|S) a_{1} + (1 - p(t|S)) a_{0} \right] + (1 - u(x,t)) a_{0} \right)$$

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$$\sim_{S} \frac{1}{n} \sum_{t \in \Omega} \sum_{x \in X} f(x|t) (u(x,t) [p(t|S)a_{1} + (1-p(t|S))a_{0}] + (1-u(x,t))a_{0}) =$$

$$= \frac{1}{n} \sum_{t \in \Omega} \sum_{x \in X} f(x|t) u(x,t) p(t|S)a_{1}$$

$$+ \left(1 - \frac{1}{n} (\sum_{t \in \Omega} \sum_{x \in X} f(x|t) u(x,t) p(t|S))\right) a_{0} =$$

$$= \left(\frac{E_{p}(u(f)|S)}{n}\right) a_{1} + \left(1 - \frac{E_{p}(u(f)|S)}{n}\right) a_{0}$$

in a similar spirit, we have

$$\frac{1}{n}\mathbf{g} + (1 - \frac{1}{n})a_0 \sim_S \left(\frac{E_p(u(\mathbf{g})|S)}{n}\right)a_1 + \left(1 - \frac{E_p(u(\mathbf{g})|S)}{n}\right)a_0$$

## Proof of the Theorem ⑤

Claim

$$f \succcurlyeq_{S} g \Leftrightarrow$$

$$\left(\frac{E_{p}(u(f)|S)}{n}\right) a_{1} + \left(1 - \frac{E_{p}(u(f)|S)}{n}\right) a_{0} \succcurlyeq_{S}$$

$$\succcurlyeq_{S} \left(\frac{E_{p}(u(g)|S)}{n}\right) a_{1} + \left(1 - \frac{E_{p}(u(g)|S)}{n}\right) a_{0}$$

#### **Proof**

by Transitivity

we conclude

$$f \succcurlyeq_S g \Leftrightarrow E_p(u(f)|S) \geqslant E_p(u(g)|S)$$

Proof (that 3rd property in theorem follows from axioms) we use

- Interest and Strict Substitution to conclude  $a_1 \succ_S a_0$  and
- Monotonicity

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Limitation

# Violations of Strict Objective Substitution

#### Example

consider four lotteries

$$f_1 = 0.1 \cdot [ \in 12m] + 0.9 \cdot [ \in 0]$$
  $f_2 = 0.11 \cdot [ \in 1m] + 0.89 \cdot [ \in 0]$   $f_3 = [ \in 1m]$   $f_4 = 0.10 \cdot [ \in 12m] + 0.89 \cdot [ \in 1m] + +.01 \cdot [ \in 0]$ 

**Preferences** 

$$f_1 \succ f_2$$

$$f_3 \succ f_4$$

#### Observation

this violates the axiom, as

$$0.5 \cdot f_1 + 0.5 \cdot f_3 = 0.5 \cdot f_2 + 0.5 \cdot f_4$$

### Example

let 
$$X = \{-€100, €100\}$$
,  $\Omega = \{L, W\}$ 

$$b_{\mathsf{L}}(\leqslant 100 | \mathsf{L}) = 1 = b_{\mathsf{L}}(-\leqslant 100 | \mathsf{W})$$
  
 $b_{\mathsf{W}}(-\leqslant 100 | \mathsf{L}) = 1 = b_{\mathsf{W}}(\leqslant 100 | \mathsf{W})$ 

- L is the event where SC Austria Lustenau wins the ADEG-cup
- W is the event where Wacker Innsbruck wins
- suppose only L or W can occur

#### Preferences

(if someone doesn't know anything about Wacker or Lustenau)

$$0.5 \cdot [ \in 100 ] + 0.5 \cdot [ - \in 100 ] \succ b_{\mathsf{L}} \quad 0.5 \cdot [ \in 100 ] + 0.5 \cdot [ - \in 100 ] \succ b_{\mathsf{W}}$$

#### Observation

this violates the axiom at least one state in  $\Omega$  must have probability greater or equal than  $\frac{1}{2}$ 

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### Cannot be Modelled

## Example

situation A

- you buy a ticket to the movies in advance (for €10)
- on the counter you realise you've lost your ticket
- you have €10, do you buy a new ticket or go home?

#### situation B

- you plan to see a movie and put €10 in your pocket
- on the counter you realise you've lost your money
- do you buy a ticket with your credit card or go home?

#### Question

what is your preference?

$$A + go home$$
 ?  $B + buy ticket$ 

#### Answer

no (strict) preference between A and B is modelled by the axiom: there is no difference

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