

# Logic: Errata

## Chapter “Why Logic is Good For You”

- page 3, “the the Nelson-Oppen” → “the Nelson-Oppen”
- page 3, “If two theories [...]. Then also [...]” → “If two theories  $T_1, T_2$  fulfil certain conditions and it is known that satisfiability of quantifier-free formulas with respect to the theories  $T_1$  and  $T_2$  is decidable, then satisfiability of quantifier-free formulas with respect to the union  $T_1 \cup T_2$  is decidable.”
- page 4, “as a unique model” → “has a unique model”
- page 5, “bad should happy” → “bad should happen”

## Chapter “Propositional Logic”

- page 8, “We write  $\vdash A$  instead of  $\emptyset \vdash A$  and call the formula  $A$  *valid* in this case.” → “We write  $\vdash A$  instead of  $\emptyset \vdash A$  and call the formula  $A$  *provable* in this case.”
- page 10, “A formula  $F$  is in said to be in *conjunctive normal form*” → “A formula  $F$  is said to be in *conjunctive normal form*”
- page 11, “otherwise it is  $C$  is said to be *inconsistent*” → “otherwise  $C$  is said to be *inconsistent*”
- page 11, Thm. 2.3, “propositional axioms” → “propositional atoms”

## Chapter “Syntax and Semantics of First-Order Logic”

- page 21, “The input  $x$  of  $M$ ” → “The input  $x$  of  $M$  such that the length of  $x$  is  $n$ ,”
- page 24, “then  $m$  is called an *isomorphism*” → “Then  $m$  is called an *isomorphism*”
- page 26, “Define a two formulas  $F$  and  $G$ , such that  $F \not\models G$  holds and  $F \not\models \neg G$  holds” → “Define two formulas  $F$  and  $G$ , such that  $F \not\models G$  holds and  $F \not\models \neg G$  holds”

## Chapter “Soundness and Completeness of First-Order Logic”

- page 31, “In we obtain  $\mathcal{J} \models \mathcal{G}$ ”  $\rightarrow$  “In sum we obtain  $\mathcal{J} \models \mathcal{G}$ ”
- page 32, “We have to prove that either every finite subset of  $\mathcal{G} \cup \{E\}$  is in  $S$  or every finite subset of  $\mathcal{G} \cup \{F\}$  is in  $S$ . As this would imply that either  $\mathcal{G} \cup \{E\} \in S^*$  or  $\mathcal{G} \cup \{F\} \in S^*$ ”  $\rightarrow$  “We have to prove that either every finite subset of  $\mathcal{G} \cup \{E\}$  is in  $S$  or every finite subset of  $\mathcal{G} \cup \{F\}$  is in  $S$ , as this would imply that either  $\mathcal{G} \cup \{E\} \in S^*$  or  $\mathcal{G} \cup \{F\} \in S^*$ .”
- page 37, “call the formula  $A$  *valid* in this case”  $\rightarrow$  “call the formula  $A$  *provable* in this case”

## Chapter “Extensions of First-Order Logic”

- page 46, “there exists exists a path”  $\rightarrow$  “exists a path”

## Chapter “Normal Forms and Herbrand’s Theorem”

- page 50, “Each step performed will preserves logical equivalence of formulas.”  $\rightarrow$  “Each step performed preserves logical equivalence of formulas.”
- page 53, add “A term  $t$  is called *closed* or *ground*, if  $t$  does not contain (free) variables.”
- page 53, “the  $\mathcal{I}$ ”  $\rightarrow$  “the interpretation  $\mathcal{I}$ ”
- page 54, “the negation of this conjunction  $D$ ”  $\rightarrow$  “the negation of this conjunction  $C$ ”
- page 56, “ $\mathcal{I} \not\models G$ ”  $\rightarrow$  “ $\mathcal{I} \not\models F$ ”
- page 58, “Let  $\mathcal{G}_2 = \{P(c), \neg P(x)\}$ ”  $\rightarrow$  “Let  $\mathcal{G}_2 = \{P(c), \neg P(a)\}$ ”

## Chapter “Automated Reasoning with Equality”

- page 61, the following definition is missing: “An equality problem  $E = x_1 \stackrel{?}{=} v_1, \dots, x_n \stackrel{?}{=} v_n$  in solved form *induces* the substitution  $\sigma_E := \{x_1 \mapsto v_1, \dots, x_n \mapsto v_n\}$ .”
- page 61, “ $\sigma'_E$ ”  $\rightarrow$  “ $\sigma_E$ ”
- page 62, “terminated”  $\rightarrow$  “terminates”

- page 62, Definition 8.4: The definition of  $\text{Res}(\mathcal{C})$  should read:

$$\text{Res}(\mathcal{C}) = \{D \mid D \text{ is conclusion of an inference in Figure 8.2 with premises in } \mathcal{C}\} \cup \mathcal{C}$$

The same mistake occurs in Def. 8.6, 8.7, Def. 8.8, and Def. 8.9.

- page 63, “More”  $\rightarrow$  “Moreover”
- page 63, “We only sketch the proof of the next theorem”  $\rightarrow$  “We only sketch the proof of the theorem”
- page 63, “the derivation is  $D$ ”  $\rightarrow$  “the derivation  $D$ ”
- page 67, the paramodulation rules should have been:

$$\frac{C \vee s \neq s'}{C\sigma'} \quad \frac{C \vee s = t \quad D \vee L[s']}{(C \vee D \vee L[t])\sigma'}$$

similar for ordered paramodulation.

- page 69, “In order to show show completeness”  $\rightarrow$  “In order to show completeness”
- page 69, “For that we show that the set of consistent set of ground clauses fulfils the satisfaction properties we need to take into account”  $\rightarrow$  “For that we need to take into account”
- page 70, change the paramodulation rules as above and “ $D[t]\sigma$  is maximal with respect to  $D[t]\sigma'$ ”  $\rightarrow$  “ $L[t]\sigma$  is maximal with respect to  $D\sigma'$ ”

## Chapter “Issues of Security”

- page 73, “Neumann-Stubblebinde”  $\rightarrow$  “Neuman-Stubblebinde”
- page 77, “the firs of the message”  $\rightarrow$  “the first of the message”
- page 77, Predicates  $\text{Store}_a$ ,  $\text{Store}_b$  are missing in the definition of the language of  $\mathcal{L}$ . “The last predicate (Nonce) denotes that its argument is a nonce.”  $\rightarrow$  “The predicate Nonce denotes that its argument is a nonce and the predicates  $\text{Store}_a$ ,  $\text{Store}_b$  denote information that is in the store of Alice or Bob.”
- page 77, “We collect these 12 sentences”  $\rightarrow$  “We collect these 15 sentences”
- page 78, “ $\forall u \forall v ((\text{Im}(u) \wedge \text{P}(v)) \rightarrow \text{Ik}(\text{key}(v, w)))$ ”  $\rightarrow$  “ $\forall u \forall v ((\text{Im}(u) \wedge \text{P}(v)) \rightarrow \text{Ik}(\text{key}(u, v)))$ ”