

**First name:** \_\_\_\_\_

**Last name:** \_\_\_\_\_

**Matriculation number:** \_\_\_\_\_

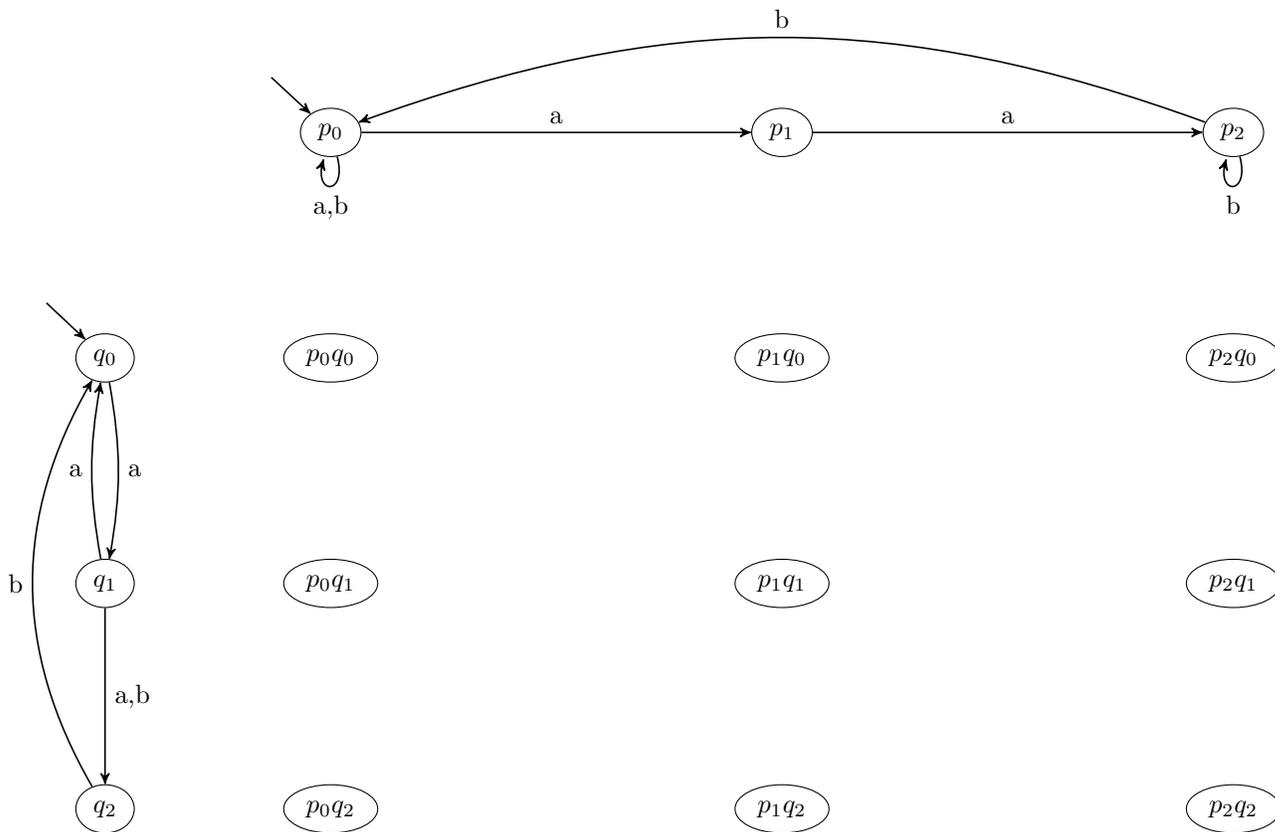
- Please answer all exercises in a readable and precise way.
- Please cross out solution attempts which are replaced by another solution.
- Please do not remove the staples of the exam.
- Cheating is not allowed. Everyone who is caught will fail the exam.

Exercise	Maximal points	Points
1	21	
2	20	
3	20	
4	9	
$\Sigma$	70	
Grade		

**Exercise 1 (14+5+2 points)**

Consider the GNBA  $\mathcal{A}_1 = (\{p_0, p_1, p_2\}, \Sigma, p_0, \delta_1, \{p_0, p_2\}, \{p_1\})$  and  $\mathcal{A}_2 = (\{q_0, q_1, q_2\}, \Sigma, q_0, \delta_2, \{q_0, q_1\})$ .

(i) Construct the GNBA  $\mathcal{A}$  for the intersection of  $\mathcal{A}_1$  and  $\mathcal{A}_2$ .



(ii) Write down the final states set(s) of  $\mathcal{A}$  explicitly.

(iii) Is  $\mathcal{L}(\mathcal{A}) = \emptyset$ ? If not, provide a word which is contained in  $\mathcal{L}(\mathcal{A})$ .

**Exercise 2 (20 points)**

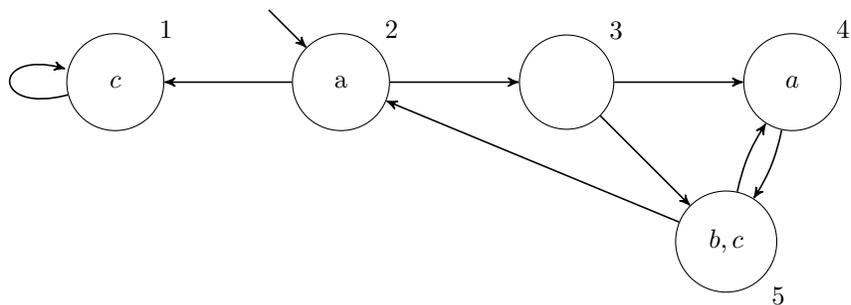
Consider the following nanoPromela statement.

```
do
  :: a => if :: d => e ! f fi ; g ! h
  :: b => i ! j ; if :: k => l ! m fi
  :: c => skip
od
```

Formally derive all transitions that are possible from this initial statement. You may use abbreviations like “do ... od” and “if ... fi”.

### Exercise 3 (20 points)

Consider the following transition system  $TS$ .



Perform CTL\*-model checking for the formula

$$\Phi = (\mathbf{A}((\mathbf{E}\mathbf{X}a) \Rightarrow \mathbf{X}a)) \wedge (\mathbf{E}(\neg a \wedge \neg b) \mathbf{U} c)$$

Here, the sets  $Sat(\Psi)$  should be indicated for every non-atomic state-subformula  $\Psi$  of  $\Phi$ . Note that the subformula  $\neg a \wedge \neg b$  of  $\Phi$  should be seen as a state-formula. It is not necessary to perform the LTL-model checking explicitly, but write down each LTL-formula that is checked.



**Exercise 4 (9 points)**

Each correct answer is worth 3 points. A wrong answer results in zero points (for that question, not for the whole exercise). Giving no answer is worth 1 point.

	Yes	No
$F(a \cup b) \equiv \neg E G \neg b$ .		
The LTL formula $F(a \cup X(b \cup X a))$ describes the following property: Every path contains two $a$ 's with a $b$ in between.		
The number of states of the GNBA $\mathcal{A}_\varphi$ for some LTL-formula $\varphi$ using the improved translation is $1 + 2^n$ where $n$ is the number of temporal operators in $\varphi$ .		