

Functional Programming

WS 2012/13

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week 5

Week 5 - λ -Calculus

Summary of Week 4

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Binary Trees

- ▶ at most 2 children per node
- ▶ applications
 - ▶ search trees
 - ▶ Huffman coding

Huffman Coding

- ▶ Idea: use shortest codewords for most frequent symbols
- ▶ Application: (lossless) data compression

This Week

Practice I

OCaml introduction, lists, strings, trees

Theory I

lambda-calculus, evaluation strategies, induction, reasoning about functional programs

Practice II

efficiency, tail-recursion, combinator-parsing

Theory II

type checking, type inference

Advanced Topics

lazy evaluation, infinite data structures, monads, ...

Origin

Goal

- ▶ find a framework in which **every** algorithm can be defined
- ▶ universal language

Result

- ▶ Turing machines (Turing, 1930s)
- ▶ **λ -Calculus** (Church, 1930s)
- ▶ ...

Syntax

λ -Terms

$$t ::= \overbrace{x}^{\text{Variable}} \mid \underbrace{(\lambda x.t)}_{\text{Abstraction}} \mid \overbrace{(t t)}^{\text{Application}}$$

$\mathcal{T}(\mathcal{V})$ set of **all** λ -terms over set of variables \mathcal{V}

Conventions

$$\begin{aligned} &(\lambda x.x) \\ &(\lambda x.(\lambda y.x)) \\ &(\lambda x.(\lambda y.(\lambda z.((x z) (y z)))) \\ &(\lambda x.((\lambda y.(\lambda z.(z y))) x)) \end{aligned}$$

Syntax

λ -Terms

$$t ::= \overbrace{x}^{\text{Variable}} \mid \underbrace{(\lambda x.t)}_{\text{Abstraction}} \mid \overbrace{(t t)}^{\text{Application}}$$

$\mathcal{T}(\mathcal{V})$ set of **all** λ -terms over set of variables \mathcal{V}

Conventions (omit outermost parentheses)

$$\begin{aligned} &\lambda x.x \\ &\lambda x.(\lambda y.x) \\ &\lambda x.(\lambda y.(\lambda z.((x z) (y z)))) \\ &\lambda x.((\lambda y.(\lambda z.(z y))) x) \end{aligned}$$

Syntax

λ -Terms

$$t ::= \overbrace{x}^{\text{Variable}} \mid \underbrace{(\lambda x.t)}_{\text{Abstraction}} \mid \overbrace{(t t)}^{\text{Application}}$$

$\mathcal{T}(\mathcal{V})$ set of **all** λ -terms over set of variables \mathcal{V}

Conventions (combine nested lambdas)

$$\lambda x.x$$

$$\lambda xy.x$$

$$\lambda xyz.((x z) (y z))$$

$$\lambda x.((\lambda yz.(z y)) x)$$

Syntax

λ -Terms

$$t ::= \overbrace{x}^{\text{Variable}} \mid \underbrace{(\lambda x.t)}_{\text{Abstraction}} \mid \overbrace{(t t)}^{\text{Application}}$$

$\mathcal{T}(\mathcal{V})$ set of **all** λ -terms over set of variables \mathcal{V}

Conventions (application is left-associative and binds strongest)

$$\lambda x.x$$

$$\lambda xy.x$$

$$\lambda xyz.x z (y z)$$

$$\lambda x.(\lambda yz.z y) x$$

Intuition

Example

λ -terms

- ▶ $\lambda x. \text{add } x \bar{1}$
- ▶ $(\lambda x. \text{add } x \bar{1}) \bar{2}$
- ▶ $\text{if true } \bar{1} \bar{0}$
- ▶ $\text{pair } \bar{2} \bar{4}$
- ▶ $\text{fst}(\text{pair } \bar{2} \bar{4})$
- ▶ $\lambda xy. \text{add } x y$
- ▶ $\lambda x. (\lambda y. \text{add } x y)$

OCaml

- ▶ `fun x -> x+1`
- ▶ `(fun x -> x+1) 2 \rightarrow^+ 3`
- ▶ `if true then 1 else 0 \rightarrow 1`
- ▶ `(2,4)`
- ▶ `fst(2,4) \rightarrow 2`
- ▶ `fun x y -> x + y`
- ▶ `fun x -> fun y -> x + y`

Remark

' $\bar{0}$ ', ' $\bar{1}$ ', ' $\bar{2}$ ', ' $\bar{3}$ ', ' $\bar{4}$ ', 'add', 'fst', 'if', 'pair', and 'true' are just abbreviations for more complex λ -terms

Subterms

Definition

$\text{Sub}(t)$ is set of subterms of t

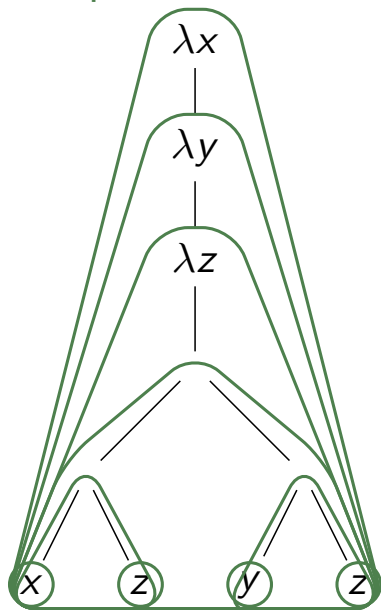
$$\text{Sub}(t) \stackrel{\text{def}}{=} \begin{cases} \{t\} & t = x \\ \{t\} \cup \text{Sub}(u) & t = \lambda x. u \\ \{t\} \cup \text{Sub}(u) \cup \text{Sub}(v) & t = u v \end{cases}$$

Example

$$\begin{aligned} \text{Sub}(\lambda xy. x) &= \{\lambda xy. x\} \cup \text{Sub}(\lambda y. x) \\ &= \{\lambda xy. x, \lambda y. x\} \cup \text{Sub}(x) \\ &= \{\lambda xy. x, \lambda y. x, x\} \end{aligned}$$

Syntax Trees

Example



$$t = \lambda xyz.x z (y z)$$

$$\text{Sub}(t) = \{t, \lambda yz.x z (y z), \lambda z.x z (y z), x z (y z), x z, y z, x, z, y\}$$

Variables

Definition variables

$$\mathcal{V}\text{ar}(t) \stackrel{\text{def}}{=} \begin{cases} \{t\} & t = x \\ \{x\} \cup \mathcal{V}\text{ar}(u) & t = \lambda x.u \\ \mathcal{V}\text{ar}(u) \cup \mathcal{V}\text{ar}(v) & t = u v \end{cases}$$

Free and Bound Variables

Definition

free variables

$$\mathcal{FVar}(t) \stackrel{\text{def}}{=} \begin{cases} \{t\} & t = x \\ \mathcal{FVar}(u) \setminus \{x\} & t = \lambda x.u \\ \mathcal{FVar}(u) \cup \mathcal{FVar}(v) & t = u v \end{cases}$$

bound variables

$$\mathcal{BVar}(t) \stackrel{\text{def}}{=} \begin{cases} \emptyset & t = x \\ \{x\} \cup \mathcal{BVar}(u) & t = \lambda x.u \\ \mathcal{BVar}(u) \cup \mathcal{BVar}(v) & t = u v \end{cases}$$

Examples

t	$\mathcal{Var}(t)$	$\mathcal{FVar}(t)$	$\mathcal{BVar}(t)$
$\lambda x.x$	$\{x\}$	\emptyset	$\{x\}$
$x y$	$\{x, y\}$	$\{x, y\}$	\emptyset
$(\lambda x.x) x$	$\{x\}$	$\{x\}$	$\{x\}$
$\lambda x.x y z$	$\{x, y, z\}$	$\{y, z\}$	$\{x\}$

Computations

Idea

- ▶ rules to manipulate λ -terms
- ▶ a single rule is enough

The β -rule (informal)

$$(\lambda x.s) t \rightarrow_{\beta} \underbrace{s\{x/t\}}_{\text{substitute } x \text{ by } t \text{ in } s}$$

application of a function to some input

Blindly replacing does not suffice

Example

- ▶ consider $\lambda xy.x$ (**fun** **x** **y** **->** **x** in OCaml)
- ▶ behavior: “take 2 arguments, ignore second, return first”
- ▶ $(\lambda xy.x) v w \rightsquigarrow (\lambda y.v) w \rightsquigarrow v$ ✓
- ▶ $(\lambda xy.x) y z \rightsquigarrow (\lambda y.y) z \rightsquigarrow z$ ✗
- ▶ clearly not intended (Problem: **variable capture**)
- ▶ $(\lambda xy.x) y z \rightarrow_{\beta} (\lambda y'.y) z \rightarrow_{\beta} y$

Solution

rename bound variables where necessary

Ocaml

```
let y = 3 and z = 2;;
(fun u -> (fun v -> u)) y z;;
```


Substitutions

Definition

function from variables to terms

$$\sigma: \mathcal{V} \rightarrow \mathcal{T}(\mathcal{V})$$

in our case we only need substitutions replacing a single variable, i.e., only for one $x \in \mathcal{V}$, $\sigma(x) \neq x$

Notation

binding for x such that $\sigma(x) \neq x$

$$\sigma = \{x/t\}$$

Example

$$\sigma = \{x/\lambda x.x\} \text{ hence } \sigma(x) = \lambda x.x \text{ and } \sigma(y) = y$$

Substitutions (cont'd)

Definition (Application)

apply substitution $\sigma = \{x/s\}$ to term t

$$t\sigma \stackrel{\text{def}}{=} \begin{cases} s & t = x \\ y & t = y, x \neq y \\ (u\sigma) (v\sigma) & t = u v \\ \lambda x.u & t = \lambda x.u \\ \lambda y.(u\sigma) & t = \lambda y.u, x \neq y, y \notin \mathcal{FVar}(s) \\ \lambda z.((u\{y/z\})\sigma) & t = \lambda y.u, x \neq y, y \in \mathcal{FVar}(s), z \text{ fresh} \end{cases}$$

Example ($\sigma = \{x/\lambda v.v w\}$)

$$\begin{array}{ll} x\sigma = \lambda v.v w & y\sigma = y \\ (x y)\sigma = (\lambda v.v w) y & (\lambda x.x y)\sigma = \lambda x.x y \\ (\lambda v.x w)\sigma = \lambda v.(\lambda v.v w) w & (\lambda w.x w)\sigma = \lambda w'.(\lambda v.v w) w' \end{array}$$

Examples

$$\begin{aligned}
& (\lambda x.x) (\lambda x.x) \rightarrow_{\beta} \lambda x.x \\
& (\lambda xy.y) (\lambda x.x) \rightarrow_{\beta} \lambda y.y \\
& (\lambda xyz.x z (y z)) (\lambda x.x) \rightarrow_{\beta} \lambda yz.(\lambda x.x) z (y z) \\
& (\lambda x.x x) (\lambda x.x x) \rightarrow_{\beta} (\lambda x.x x) (\lambda x.x x) \\
& \lambda x.x \rightarrow_{\beta} \text{no } \beta\text{-step possible} \\
& \lambda x.(\lambda y.y) z \rightarrow_{\beta} \lambda x.\underline{z}
\end{aligned}$$

 β -Reduction

Definition (Context)

context $C \in \mathcal{C}(\mathcal{V})$

$$C ::= \square \mid \lambda x.C \mid C t \mid t C$$

with $x \in \mathcal{V}$ and $t \in \mathcal{T}(\mathcal{V})$

- ▶ $C[s]$ denotes replacing \square by term s in context C

Example

$$\begin{array}{ll}
C_1 = \square & C_1[\lambda x.x] = \lambda x.x \\
C_2 = x \square & C_2[\lambda x.x] = x (\lambda x.x) \\
C_3 = \lambda x.\square x & C_3[\lambda x.x] = \lambda x.(\lambda x.x) x
\end{array}$$

β -Reduction (cont'd)

Definition (β -step)

if exist context C and terms s , u , and v such that

$$s = C[(\lambda x.u) v]$$

then

$$s \rightarrow_{\beta} C[u\{x/v\}]$$

is a β -step with **redex** $(\lambda x.u) v$ and **contractum** $u\{x/v\}$

- ▶ $s \xrightarrow{+}_{\beta} t$ denotes sequence $s = t_1 \rightarrow_{\beta} t_2 \rightarrow_{\beta} \dots \rightarrow_{\beta} t_n = t$ with $n > 0$
- ▶ $s \xrightarrow{*}_{\beta} t$ is sequence with $n \geq 0$ (s **β -reduces** to t)

β -Reduction

Example

$$\Omega = (\lambda x.x x) (\lambda x.x x)$$

$$K_* = \lambda xy.y$$

$$I_2 = \lambda xy.x y$$

$$K_* \Omega \rightarrow_{\beta} K_* \Omega \rightarrow_{\beta} \dots$$

$$K_* \Omega \rightarrow_{\beta} \lambda y.y$$

$$\begin{aligned}
 I_2 I_2 &= (\lambda xy.x y) (\lambda xy.x y) \rightarrow_{\beta} \lambda y.(\lambda xy.x y) y \equiv \lambda y.(\lambda xy'.x y') y \\
 &\rightarrow_{\beta} \lambda y.(\lambda y'.y y') = \lambda yy'.y y' \equiv I_2
 \end{aligned}$$

What Are the Results of Computations?

Idea

- ▶ only **terms** in λ -calculus
- ▶ express functions **and** values through λ -terms

Definition (Normal form)

$t \in \mathcal{T}(\mathcal{V})$ is in **normal form** (NF) if no β -step possible

Example

$$\begin{array}{ll} \lambda x.x & \text{NF} \\ (\lambda x.x) y & \text{not NF} \end{array}$$

Lambda Interpreter for Pure Students

developed by Michael Brunner (bachelor thesis)

λ -Terms

$$t ::= x \mid (\lambda x.t) \mid (t t)$$

Conventions

- ▶ interpreter command `!pretty` toggles use of conventions for printing
- ▶ nested abstractions use spaces to separate variable names, e.g.,

$$\begin{array}{ll} \lambda xy.x & \lambda x y.x \\ \lambda x_1.y & \lambda x_1.y \end{array}$$

Result

Normal Forms

- ▶ result of input is corresponding NF
 - ▶ `> (\x.x) (\x.x)`
NF: `(\x.x)`

Evaluation Strategy

- ▶ `!by_value` activates call-by-value evaluation (next lecture)
- ▶ `!by_name` activates call-by-name evaluation (next lecture)
- ▶ `!trace` toggles tracing

Abbreviations & Initialisation

Interpreter Command

```
!def <name> = t
```

Example

```
> !def I = \x.x
> !def K = \x y.x
> !def S = \x y z.x z (y z)
> S K I
NF: \z.z
```

```
.lambdainit
```

content of file `.lambdainit` is loaded on start-up of `lips`