

Functional Programming

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Efficiency of Functional Programs

Avoid unnecessary recomputations by ...

- ▶ tupling

Introduce tail recursion by ...

- ▶ parameter accumulation

This Week

Practice I

OCaml introduction, lists, strings, trees

Theory I

lambda-calculus, evaluation strategies, induction, reasoning about functional programs

Practice II

efficiency, tail-recursion, **combinator-parsing**

Theory II

type checking, type inference

Advanced Topics

lazy evaluation, infinite data structures, monads, . . .

What is Parsing?

*Parsing is the decomposition of a **linear sequence** into a **structure**, given by a **grammar**. This linear sequence may be a text in some natural language, a computer program, a web site, a piece of music, a sequence of genes, ...*

- ▶ linear sequence: `'t list` (list of tokens)
- ▶ structure: some user-defined type (abstract syntax tree)
- ▶ grammar: BNF (Backus-Naur form)

Note

- ▶ BNF can express context-free grammars (CFG)
- ▶ combinator parsers can parse context-sensitive grammars
- ▶ however, for the purpose of this lecture CFG suffice

Usual Two Phases - Lexing and Parsing

Lexing

- ▶ divide original input (list of `chars`) into tokens
- ▶ white space and comments may be dropped at this stage
- ▶ syntactic check

Parsing

- ▶ work on list of tokens
- ▶ check if list of tokens corresponds to grammar
- ▶ produce an abstract syntax tree (AST)
- ▶ semantic check

Propositional Formulas

Grammar

$$\phi ::= p \mid (! \phi) \mid (\phi \& \phi)$$

Lexing

token	corresponds to
LPAR	(
RPAR)
NOT	!
AND	&
ID of string	propositional atom

Example

"(a &(!b))" $\xrightarrow{\text{lexing}}$ [LPAR; ID "a"; AND; LPAR; NOT; ID "b"; RPAR; RPAR]
 "(a a)" $\xrightarrow{\text{lexing}}$ [LPAR; ID "a"; ID "a"; RPAR]

Propositional Formulas (cont'd)

Grammar

$$\phi ::= p \mid (! \phi) \mid (\phi \& \phi)$$

Parsing (AST)

```

type t = Atom of string
         | And of t * t
         | Not of t
  
```

Example

```

[LPAR; ID" a"; AND; LPAR; NOT; ID" b"; RPAR; RPAR]  $\xrightarrow{\text{parsing}}$ 
And(Atom" a", Not(Atom" b"))
[LPAR; ID" a"; ID" a"; RPAR]  $\xrightarrow{\text{parsing}}$  ✗
  
```

Parsers

First Attempt

- ▶ functions of type `'t list -> ('a * 't list)`
- ▶ e.g., `digit ['1';'2']` results in `('1', ['2'])`
- ▶ but what about **errors**, i.e., `digit ['x';'y'] = ?`

Type of Parsers

```
type ('a,'t)t = 't list -> ('a * 't list)option
```

- ▶ a parser works on a list of tokens of arbitrary type `'t`
- ▶ a successful parse yields `Some(x,ts)`
 - ➡ result `x` and remaining tokens `ts`
- ▶ a parse error is represented by `None`
 - ➡ no further information about the error

Preparation

Applying a Parser

```
let test p s = match p (Strng.of_string s) with
  | None          -> failwith "parse_error"
  | Some (x,ts)  -> (x,ts)
```

Example

- ▶ `letter ...` parses a single letter (see page 12)
test letter "hello_world"
- char * char list = ('h', "ello_world")
test letter "1234"
Exception: Failure "parse_error".
- ▶ `digit ...` parses a single digit (see page 12)
test digit "1234"
- char * char list = ('1', "234")

Preparation (cont'd)

Type of a Parser

- ▶ type of parser:

```
type ('a,'t)t = 't list -> ('a * 't list)option
```

- ▶ often we want to process single tokens only, i.e., $f : 't \rightarrow 'a$
- ▶ `token` lifts f to parser, i.e., `token f : ('a,'t)t`

Operating on a Single Token

```
let token f = function  
  | []      -> None  
  | t::ts  -> match f t with  
    | Some x -> Some (x,ts)  
    | None   -> None
```

Primitive Parsers

Any Token Satisfying a Condition

```
let sat p ts = token(fun t -> if p t then Some t else None) ts
```

Checking for End of Input

```
let eoi = function [] -> Some ((),[])
             | _ -> None
```

Example

```
# test (sat ((=) 'h')) "hello_world"
- char * char list = ('h', "ello_world")
# test (sat ((=) 'e')) "hello_world"
Exception Failure "parse_error".
# test eoi ""
- : unit * char list = ((), "")
```

Character Parsers

Reading a Given Character

```
let any ts = sat (fun _ -> true) ts
```

```
let char c ts = sat ((=) c) ts
```

Reading Letters and Digits

```
let letter =  
  sat(fun c -> ('a' <= c && c <= 'z') || ('A' <= c && c <= 'Z'))
```

```
let digit = sat (fun c -> '0' <= c && c <= '9')
```

Character Parsers (cont'd)

Choosing From a List of Tokens

```
let oneof s = sat (fun c -> Lst.mem c (Strng.of_string s))
```

Every Token Except a Given List

```
let noneof s = sat (fun c -> not(Lst.mem c (Strng.of_string s)))
```

White Space

```
let space = oneof "_\n\r\t"
```

Example

```
# test space "_ _ _ hello _ world"  
- : char * char list = (' ', "_ _ _ hello _ world")
```

Parser Combinators

Choice

`('a,'t)t -> ('a,'t)t -> ('a,'t)t`

```
let (<|>) p q ts = match p ts with None          -> q ts
                        | Some _ as r -> r
```

- ▶ choice takes two parsers `p` and `q` of same result type
- ▶ if `p` is successful on `ts` then its result is returned
- ▶ otherwise `q` is applied on `ts`

Example

```
# test (letter <|> digit) "hello_world"
- : char * char list = ('h', "hello_world")
# test (letter <|> digit) "1234"
- : char * char list = ('1', "234")
```

Parser Combinators (cont'd)

Bind - Binding a Parser to the Result of Another

$$('a, 't)t \rightarrow ('a \rightarrow ('b, 't)t) \rightarrow ('b, 't)t$$

```
let (>>=) p f ts = match p ts with None      -> None
                        | Some(x,ts) -> f x ts
```

- ▶ bind takes two arguments
- ▶ first a parser `p` with results of type `'a`
- ▶ then a function `f` taking an `'a` and producing a parser with results of type `'b`
- ▶ `p >>= f` executes `p` and then feeds the function `f` with its result
- ▶ since `f x` is a parser, the result of `p >>= f` is a parser

Parser Combinators (cont'd)

Then - Sequential Composition of Parsers

$(\text{'a}, \text{'t})t \rightarrow (\text{'b}, \text{'t})t \rightarrow (\text{'b}, \text{'t})t$

```
let (>>) p q ts = (>>=) p (fun _ -> q) ts
```

- ▶ then takes two parsers `p` and `q`
- ▶ first a parser `p` with results of type `'a`
- ▶ then a parser `q` with result type `'b`
- ▶ `p >> q` first executes `p` and then executes `q`
- ▶ the result of `p` is ignored, the result of `q` is returned

Example

```
# test (space >> space) "   hello_world"
- : char * char list = (' ', "hello_world")
```


Turning Values Into Parsers

Return

```
'a -> ('a,'t)t
```

```
let return x = fun ts -> Some (x,ts)
```

- ▶ `return x` takes the value `x` and yields a parser that returns `x` without consuming any input

Example

```
# test (return 3) "hello_world"
```

```
- : int * char list = (3, "hello_world")
```

Counting Spaces

First try

```
let rec count_spaces =  
  (space >>  
    count_spaces >>= fun i -> return (i+1))  
<|>  
  (any >>  
    count_spaces >>= fun i -> return i)  
<|>  
  (eoi >> return 0)
```

Problem

- ▶ OCaml interpreter detects that `count_spaces` would not terminate (eager evaluation)
- ▶ Error: This kind of expression is not allowed as right-hand ...

Counting Spaces (cont'd)

Second try (dummy argument)

```
let rec count_spaces () =  
  (space >>  
    count_spaces () >>= fun i -> return (i+1))  
<|>  
  (any >>  
    count_spaces () >>= fun i -> return i)  
<|>  
  (eoi >> return 0)
```

Problem

- ▶ OCaml interpreter doesn't detect nontermination (eager evaluation)
- ▶ `# test (count_spaces ()) "hello_world"`
Stack overflow during evaluation (looping recursion?).

Counting Spaces (cont'd)

Third try (enforce evaluation)

```
let rec count_spaces () =  
  (space >>= fun _ ->  
    count_spaces () >>= fun i -> return (i+1))  
<|>  
  (any >>= fun _ ->  
    count_spaces () >>= fun i -> return i)  
<|>  
  (eoi >> return 0)
```

Success

```
# test (count_spaces ()) "hello_world"  
- : int * char list = (1, "")  
# test (count_spaces ()) "   hello_world"  
- : int * char list = (4, "")
```

Parser Combinators (cont'd)

Many

```
let rec many p = (  
  p      >>= fun x ->  
  many p >>= fun xs ->  
  return (x::xs)  
) <|> return []
```

- ▶ `many p` applies `p` **zero** or more times
- ▶ result is list of results of `p`
- ▶ greedy (as many applications of `p` as possible)

Example

```
# test (many space) "   hello_world"  
- : char list * char list = ("   ", "hello_world")
```

Arithmetic

A Grammar for Arithmetic

$$\begin{array}{lll}
 e ::= e + t \mid t & t ::= t * f \mid f & f ::= (e) \mid n \\
 n ::= d n \mid d & d ::= 0 \mid \dots \mid 9 &
 \end{array}$$

Type of the AST

```

type arith = Num of int
           | Add of arith * arith
           | Mul of arith * arith
  
```

Example

```

# test ??? "3+2*5"
- : arith * char list = (Add (Num 3, Mul (Num 2, Num 5)), "")
  
```

A Parser for Arithmetic (does not terminate)

```
let rec e() =
  (e() >>= fun e1 -> char '+' >> t() >>= fun e2 ->
    return(Add(e1,e2)))
  <|> (t())
and t() =
  (t() >>= fun t1 -> char '*' >> f() >>= fun t2 ->
    return(Mul(t1,t2)))
  <|> (f())
and f() = (
  char '(' >>= fun _ ->
  e() >>= fun e1 ->
  char ')') >>
  return e1
) <|> n
and n = many1 digit >>= fun r ->
  return(Num (int_of_string (Strng.to_string r)))
```

Problem & Solution

- ▶ problem: left recursive grammar
- ▶ solution: eliminate left recursion

Non-left recursive Grammar for Arithmetic

$$e ::= t e'$$

$$e' ::= + t e' \mid \epsilon$$

$$t ::= f t'$$

$$t' ::= * f t' \mid \epsilon$$

$$f ::= (e) \mid n$$

$$n ::= d n \mid d$$

$$d ::= 0 \mid \dots \mid 9$$

A Parser for Arithmetic

```

let rec e() = t() >>= e'
and e' term = (char '+' >>= fun _ -> t() >>= e' >>= fun t2 ->
  return(Add(term,t2))
) <|> return term
and t() = f() >>= t'
and t' fact = (char '*' >>= fun _ -> f() >>= t' >>= fun f2 ->
  return(Mul(fact,f2))
) <|> return fact
and f() = (char '(' >>= fun _ -> e() >>= fun e1 -> char ')') >>
  return e1
) <|> n
and n = many1 digit >>= fun r ->
  return(Num (int_of_string (Strng.to_string r)))

```

Example

```

# test (e ()) "3+2*5";;
- : arith * char list = (Add (Num 3, Mul (Num 2, Num 5)), "")

```