Functional Programming

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week 8

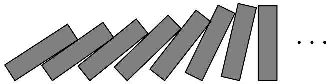
Mathematical Induction

Induction Principle

$$(\underbrace{P(m)}_{\text{base case}} \wedge \underbrace{\forall k \geq m.(P(k) \to P(k+1))}_{\text{step case}}) \to \forall n \geq m.P(n)$$

Example

- first domino will fall
- if a domino falls also its right neighbor falls



Induction on Lists

Induction Principle (without Types)

$$(\underbrace{P([])}_{\text{base case}} \wedge \underbrace{\forall x.\forall xs.(P(xs) \to P(x :: xs))}_{\text{step case}}) \to \forall ls.P(ls)$$

Lemma (append is associative)

$$xs \mathbin{@} (ys \mathbin{@} zs) = (xs \mathbin{@} ys) \mathbin{@} zs$$

where

```
let rec (@) xs ys = match xs with
  | []    -> ys
  | x::xs -> x :: (xs @ ys)
```

Proof.

Blackboard □

Structural Induction

Usage

- can be used on every variant type
- base cases correspond to non-recursive constructors
- step cases correspond to recursive constructors

Example

- lists
- trees
- λ-terms
- ...

This Week

Practice I
OCaml introduction, lists, strings, trees

Theory I
lambda-calculus, evaluation strategies, induction,
reasoning about functional programs

Practice II
efficiency, tail-recursion, combinator-parsing, dynamic programming

Theory II
type checking, type inference

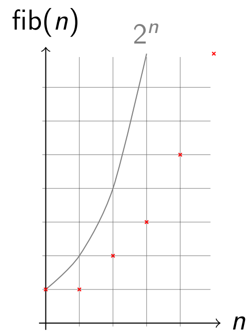
Advanced Topics
lazy evaluation, infinite data structures, monads, ...

Mathematical

Definition (n-th Fibonacci number)

$$\mathrm{fib}(n) \stackrel{\text{def}}{=} \begin{cases} 1 & \text{if } n \leq 1 \\ \mathrm{fib}(n-1) + \mathrm{fib}(n-2) & \text{otherwise} \end{cases}$$

Graph



Mathematical (cont'd)

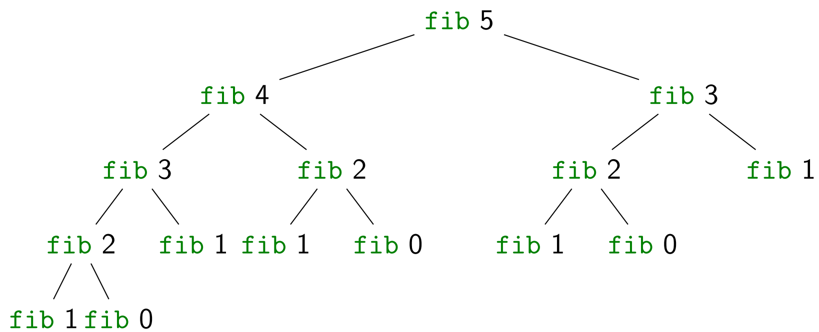
Example
1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, 2584, 4181 ,6765, 10946, 17711, 28657, 46368, 75025, 121393, 196418, 317811, 514229, 832040, 1346269, 2178309, 3524578, 5702887, 9227465, 14930352, 24157817, 39088169, 63245986, 102334155, 165580141, 267914296, 433494437, 701408733, 1134903170, 1836311903, 2971215073, ...

OCaml

Definition

```
let rec fib n = if n < 2 then 1 else fib(n-1) + fib(n-2)
```

Example



Tupling

Idea

- use tuples to return more than one result
- make results available as return values instead of recomputing them

Fibonacci Numbers

Example

```
let rec fibpair n = if n < 1 then (0,1) else (
 if n = 1 then (1,1)
          else let (f1,f2) = fibpair (n-1) in (f2,f1+f2)
)
```

- this function is linear

Lemma

$$\texttt{fibpair}(n+1) = (\texttt{fib}\ n, \texttt{fib}(n+1))$$

Proof.

Blackboard □

A Second Example

Goal

compute average value of an integer list (module IntLst)

Naive Approach

- `let average xs = sum xs / Lst.length xs`
- 2 traversals of `xs` are done

Combined Function

- ```
  let rec sumlen = function
   | []    -> (0,0)
   | x::xs -> let (sum,len) = sumlen xs in (sum+x,len+1)
  ```
- `let average1 xs = let (sum,len) = sumlen xs in sum/len`
- one traversal of `xs` suffices

Recursion vs. Tail Recursion

Idea

- a function calling itself is recursive
- functions that mutually call each other are mutually recursive
- special kind of recursion is tail recursion

Definition (Tail recursion)

a function is called tail recursive if the recursive call is last action in the function body

Reward

`http://xkcd.com/1270/`
  ```

Examples

Length

```
let rec length = function []    -> 0
                        | _::xs -> 1 + length xs
```

- not tail recursive

Even/Odd

```
let rec is_even = function 0 -> true
                         | 1 -> false
                         | n -> is_odd(n-1)
and is_odd      = function 0 -> false
                         | 1 -> true
                         | n -> is_even(n-1)
```

- mutually recursive (btw: also tail recursive)

Parameter Accumulation

Idea

- make function tail recursive
- provide data as input instead of computing it before recursive call
- Why? (tail recursive functions can automatically be transformed into space-efficient loops)

Example (Sumlen)

```
let rec sumlen = function
 | []    -> (0,0)
 | x::xs -> let (sum,len) = sumlen xs in (sum+x,len+1)
```

- not tail recursive

```
let sumlen_tl xs =
 let rec sumlen sum len = function
  | []    -> (sum,len)
  | x::xs -> sumlen (sum+x) (len+1) xs
 in
 sumlen 0 0 xs
```

- tail recursive

```
let sumlen_fold xs =
 Lst.foldl (fun (sum,len) x -> (sum+x,len+1)) (0,0) xs
```

- tail recursive

Example (Range)

```
let rec range m n = if m >= n then []
                            else m::range (m+1) n
```

- not tail recursive

```
let range_tl m n =
 let rec range acc m n =
  if m >= n then acc else range ((n-1)::acc) m (n-1)
 in
 range [] m n
```

- tail recursive

Examples (Reverse)

- ```
  let rec reverse = function    [] -> []
                            | x::xs -> (reverse xs) @ [x]
  ```
- not tail recursive
- ```
  let rev xs =
   let rec rev acc = function []    -> acc
                            | x::xs -> rev (x::acc) xs
   in
   rev [] xs
  ```
- tail recursive
- ```
  let rev xs = Lst.foldl (fun acc x -> x::acc) [] xs
  ```
- tail recursive
  ```