

Automated Theorem Proving

Georg Moser

Institute of Computer Science @ UIBK

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Summary of Last Lecture

Definition

$$\frac{\gamma}{\gamma(t)} \quad t \text{ term in } \mathcal{L}^+ \qquad \frac{\delta}{\delta(k)} \quad k \text{ fresh constant in } \mathcal{L}^+$$

- 1 \mathcal{L}^+ denotes extension of base language \mathcal{L}
- 2 new individual constants are introduced in δ rules
- 3 **fresh** means new to the branch of the tableau

Theorem

a sentence F is valid iff F has a tableau proof

Outline of the Lecture

Early Approaches in Automated Reasoning

Herbrand's theorem for dummies, Gilmore's prover, method of Davis and Putnam

Starting Points

resolution, tableau provers, Skolemisation, ordered resolution, redundancy and deletion

Automated Reasoning with Equality

paramodulation, ordered completion and proof orders, superposition

Applications of Automated Reasoning

Neuman-Stubblebinde Key Exchange Protocol, Robbins problem

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First-Order Tableau

Example

consider the tableau proof of

$$\exists x \forall y R(x, y) \rightarrow \forall y \exists x R(x, y)$$

on the whiteboard



Free-Variable Semantic Tableaux

Definition (expansion rules)

$$\frac{\gamma}{\gamma(x)} \quad x \text{ a free variable} \qquad \frac{\delta}{\delta(f(x_1, \dots, x_n))} \quad f \text{ a Skolem function}$$

- x_1, \dots, x_n denote all free variables of the formula δ
- Skolem function f must be new on the branch



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Remark

- δ -rule still leaves a lot of room for improvement
- requirement that f must be new on the branch forces the introduction of inefficiently many new Skolem functions
- prevented with cleverer notions of the δ -rule

Definition (atomic closure rule)

- 1 \exists branch in tableau T that contains two literals A and $\neg B$
- 2 \exists mgu σ of A and B
- 3 then $T\sigma$ is also a tableau



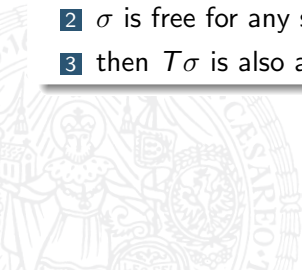
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- 1 T is a tableau for \mathcal{G}
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Remark

completeness of **free-variable tableaux** can (eventually) be proven via model existence

Example

consider the tableau proof of

$$\exists x \forall y R(x, y) \rightarrow \forall y \exists x R(x, y)$$

and

$$\forall x \forall y (P(x) \wedge P(y)) \rightarrow \forall x \forall y (P(x) \vee P(y))$$

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Soundness of Free-Variable Tableaux

Definition

- a branch in a free-variable tableau is called **satisfiable**, if \exists structure \mathcal{A} and \forall environment $\ell: (\mathcal{A}, \ell) \models \mathcal{G}$
- a free-variable tableau is **satisfiable**, if there exists a satisfiable branch



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Proof

the lemma follows by case-distinction on the applied expansion rule, it suffices to consider the δ -rule all other cases are similar

Proof (cont'd).

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$$f^{\mathcal{A}'}(\ell(x_1), \dots, \ell(x_n)) := a$$

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$$(\mathcal{A}, \ell) \models \delta \quad \Longrightarrow \quad (\mathcal{A}', \ell) \models \delta(f(x_1, \dots, x_n))$$

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if the **substitution rule** is applied to a satisfiable tableau T , then its result is satisfiable



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Strong Completeness of Free-Variable Tableaux

NB: may consider a sequence of atomic closure rules that leads to an (atomically closed) tableau as one block



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Definition

- T be a tableau with branches B_1, \dots, B_n
- $\forall i A_i$ and $\neg B_i$ are literals on B_i
- if σ is a mgu of $A_1 = B_1, \dots, A_n = B_n$
- then σ is called **most general atomic closure substitution**



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Lemma (Lifting Lemma)

- 1 τ a substitution free for tableau T such that each branch in $T\tau$ is atomically closed
- 2 then \exists a most general atomic closure substitution σ and
- 3 $T\sigma$ is closed by n applications of the atomic closure rule

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a **strategy** S details:

- 1 which expansion rule is supposed to be applied
- 2 or that no expansion rule can be applied

a strategy may use extra information which is updated



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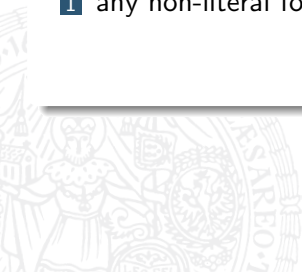
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Example

strategy employed in the implementation of free-variable tableaux is fair

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- for each tableau the extra information includes
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this strategy is **not** fair

Theorem (Strong Completeness)

- 1 S be a fair strategy
- 2 F be a valid sentence
- 3 F has a tableau proof with the following properties:
 - all tableau expansion rules are considered first and follow strategy S
 - a block of atomic closure rules closes the tableau



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^athe formulas on the branch form a Hintikka set

Implementation of γ -Rule

γ -rule (simplified)

```
singlestep([OldBranch | Rest], NewTree) :-  
    member(NotatedGamma, OldBranch),  
    notation(NotatedGamma, Free),  
    fmla(NotatedGamma, Gamma),  
    is_universal(Gamma),  
    remove(NotatedGamma, OldBranch, TempBranch),  
    NewFree = [V | Free],  
    instance(Gamma, V, GammaInstance),  
    notation(NotatedGammaInstance, NewFree),  
    fmla(NotatedGammaInstance, GammaInstance),  
    append([NotatedGammaInstance | TempBranch],  
          [NotatedGamma], NewBranch),  
    append(Rest, [NewBranch], NewTree).
```