

Interactive Theorem Proving

Week 11

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Summary

So far

Proof Assistants, HOL Light, λ_{\rightarrow} , λ_P , λ_2 , paradoxes, set theory

- ZF(C) Axiomatization
- Bootstrapping Mathematics in set theory

Today

Mizar

Mizar History

Project started in 1973

- Attempt to reconstruct the mathematical vernacular
- In a computer environment

Today

- Formal language for writing proofs
 - Designed to be close to mathematical vernacular
 - Subset of standard English
 - Declarative style
 - Highly structured
 - rigorous and semantically unambiguous
- System for verifying the proofs
- Mizar Mathematical Library (MML)

Systematic collection

- Started 1989
- Based on Tarski-Grothendieck axiomatization
- Incremental Revisions

Biggest library of formal mathematics

- 1177 articles
- 52775 theorems
- 10670 definitions
- 820 schemes
- 11333 registrations

Logic

- classical first order logic
- second order free variables
 - for recursive definitions or induction
- natural deduction
 - usually forward reasoning

independent of the axioms of set theory

- First article: defining set and element
- Second article: axiomatization of set theory

Mizar article

- A `.miz` file (optionally `.voc` and `.abs`)
- Environment
- Theorems (local lemmas)
- Definitions
- Schemes
- Justifications
 - Simple justification
 - Proof
 - Schematization

Justifications and proofs

Local justification

```
A: statement_1;  
  ...  
  statement_2 by A;
```

External justification

```
x in { x } by TARSKI:def 1;
```

Proof structure

```
statement  
proof  
  ...  
  thus statement;  
end;
```

Natural deduction (1/3)

Conjunction

```
A & B
proof
  ...
  thus A;
  ...
  thus B;
end;
```

Implication

```
A implies B
proof
  assume A;
  ...
  thus B;
end;
```


Natural deduction (2/3)

Disjunction

A or B

proof

 assume not A;

 ...

 thus B;

end;

Comment

:: this is a comment

Natural deduction (3/3)

Equivalence

A iff B

proof

 thus A implies B

 proof

 assume A;

 ...

 thus B;

 end;

 thus B implies A

 proof

 assume B;

 ...

 thus A;

 end;

end;

Predicate logic

Universal Quantification

```
for x being T holds P[x]
proof
  let x be T;
  ...
  thus P[x];
end;
```

Existential Quantification

```
ex x being T st P[x]
proof
  ...
  take x = expression;
  ...
  thus P[x];
end;
```

1.5-order logic: Schemes

```
scheme
  Ind { P[Nat] } : for k being Element of NAT holds P[k]
provided
  A1: P[0] and
  A2: for k being Element of NAT st P[k] holds P[k + 1]
proof
  ...
end;
```

1.5-order logic: Schematization

```
2 divides n * (n+1)
proof
  :: local predicate
  defpred P[Nat] means 2 divides $1 * ($1 + 1);
  a1: P[0];
  a2: for k being Nat st P[k] holds P[k + 1];
  :: referring to the scheme
  for k being Nat holds P[k] from NAT_1:sch 2(a1,a2);
  hence 2 divides n * (n + 1);
end;
```

Implementation

- Separate processes
- Parser
 - Environment
 - Operator syntax
- Analyzer
 - Disambiguation
 - **Types!** (adjectives)
- Checker
 - Disprover processes all disjuncts
 - Forms of a term
 - Congruence closure
- Post-processing
 - Relprem, Relinfer, ...

27 special symbols

&
c=

110 reserved words

contradiction
not
or
implies
iff
for x holds a(x)
ex x st a(x)

Mizar Types

A type hierarchy

- Function of X,Y
- PartFunc of X,Y
- Relation of X,Y
- Subset of [:X,Y:]
- set

Adjectives

- Examples
 - one-to-one Function of X,Y
 - finite non-empty proper Subset of X
- Automatic deriving of type information using registrations
- Overloading of notations
- Types must be non-empty

No set of inference rules

“obviousness w.r.t. an algorithm” by M. Davis

de Bruijn criterion is not preserved

- new computation mechanisms (CAS, DS)
- more automation in the equality calculus
- more general statements in an inference

Congruence closure

Is $x = y$ a consequence of $y = z$, $f(y) = z$, and $f(z) = x$?

- Monotonicity
- Transitivity
- Symmetry

Summary

Today

- Mizar project
- Foundations
- Natural deduction
- Checker and verifier

Next time

- Program extraction
- Logical frameworks