

# Logic Programming

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## Summary of Last Lecture

### Example (meta-variable facility)

X; Y :- X.

X; Y :- Y.

### Definitions (second-order programming)

- the predicate *bagof( Template, Goal, Bag )* unifies *Bag* with the alternatives of *Template* that meet *Goal*
- if *Goal* has free variables besides the one sharing with *Template* *bagof* will backtrack
- fails if *Goal* has no solutions
- construct  $\text{Var}^{\exists} \text{Goal}$  tells *bagof* to existentially quantify *Var*
- the predicate *setof( Template, Goal, Bag )* is similar to *bagof* but sorts the obtained multi-set (bag) and removed duplicates

# Outline of the Lecture

## Monotone Logic Programs

introduction, basic constructs, logic foundations, unification, semantics, database and recursive programming, termination, complexity

## Incomplete Data Structures and Constraints

incomplete data structures, definite clause grammars, constraint logic programming, answer set programming

## Full Prolog

semantics (revisted), cuts, correctness proofs, meta-logical predicates, pragmatics, efficient programs, meta programming

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- **space usage depends also on the number of data structures created**
- the former may be a major problem: **stack overflow**

## Example

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sublist(Xs,AXBs) :- suffix(XBs,AXBs), prefix(Xs,XBs).  
sublist(Xs,AXBs) :- prefix(AXs,AXBs), suffix(Xs,AXs).
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What is better, if we argue wrt a linked-list implementation of cons lists?

## Answer

the first alternative:

- consider

```
sublist([1,2,3,4],[1,2,3,4,1,2,3,4,1,2,3,4,1,2,3,4])
```

- the 1st clause iterates over the 2nd list to find a suitable suffix
- then iterates over the first list
- no intermediate data structures are created

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```
- the 1st clause iterates over the 2nd list to find a suitable suffix
- then iterates over the first list
- no intermediate data structures are created
- in the 2nd clause an auxilliary list is created

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we say: the first clause doesn't **cons**

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- equivalently the number of unifications (performed and attempted) asymptotically bounds the runtime
- on the other hand, if unification needs to be taken into account time complexity analysis is more involved
- in general size of search space and size of input terms needs to be taken into account

# Howto Improve Performance

Suggestion ①

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reverse([X|Xs],Zs) :-  
    reverse(Xs,Ys),  
    append(Ys,[X],Zs).  
reverse([],[]).
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reverse([X|Xs],Acc,Ys) :-  
    reverse(Xs,[X|Acc],Ys).  
reverse([],Ys,Ys).
```

# Excursion: Transforming Recursion into Iteration

## Definitions

- a Prolog clause is called **iterative** if
  - 1 it has one recursive call, and
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## Example (Factorial Iterative, Version 1)

```
factorial(N,F) :- factorial(0,N,1,F).  
  
factorial(I,N,T,F) :-  
    I < N, I1 is I + 1, T1 is T*I1, factorial(I1,N,T1,F).  
factorial(N,N,F,F).
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## Example (Factorial Iterative, Version 2)

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factorial(N,T,F) :-  
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factorial(0,F,F).
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between(I,J,I) :- I ≤ J.  
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## Example

```
sumlist([ ],Sum,Sum).  
sumlist([I|Is],Temp,Sum) :-  
    Temp1 is Temp + I, sumlist(Is,Temp1,Sum).  
sumlist([],Sum,Sum).
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## Example

```
maximum([X|Xs],M) :- maximum(Xs,X,M).  
  
maximum([X|Xs],Y,M) :-  
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length([X|Xs],N) :-  
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length([],0).
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## Suggestion ②

tuning, via:

- 1 good goal order
- 2 elimination of (unwanted) nondeterminism by using explicit conditions and cuts
- 3 exploit clause indexing (order arguments suitably)  
**indexing** performs static analysis to detect clauses which are applicable for reduction

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append([X|Xs], Ys, [X|Zs]) :-  
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*By default, SWI-Prolog, as most other implementations, indexes predicates on their first argument.*

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## Definition (tail recursion optimisation)

- consider a generic clause for  $A$

$$A' : -B_1, \dots, B_n$$

such that  $A$  and  $A'$  unify with  $\sigma$

- suppose the goal  $B_1\sigma, \dots, B_{n-1}\sigma$  is deterministic
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## Definition

**clause indexing** is used to detect which clauses are applicable for reduction: 2nd clause in append need only be considered for empty lists

# How to Implement Functions

## Functions vs Relations

- often, we want to compute functions:
  - 1 addition:  $\mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$
  - 2 sorting:  $list \rightarrow list$

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- in logic programming we specify relations and every function can be seen as a relation

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- that is, we implement **functions**  $f(i_1, \dots, i_n) = (o_1, \dots, o_m)$  by **relations**  $f_{rel}/(n + m)$
- result is obtained by **query**  $f_{rel}(i_1, \dots, i_n, X_1, \dots, X_m)$ 
  - 1 addition:  $plus(n, m, Z)$   $Z = n + m$
  - 2 sorting:  $sort(list, Xs)$   $Xs = \text{sorted version of list}$

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- solution: **store result of each sub-expression in fresh variable**

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$$\underbrace{\qquad\qquad\qquad}_{u}$$

$$\underbrace{\qquad\qquad\qquad}_{v}$$

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$f(X, Y) :- \text{times}(X, X, Z), \text{minus}(Z, 5, V),$

$$\begin{array}{c}
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 \qquad\qquad\qquad \underbrace{\phantom{x^2 - 5}_z}_v \\
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 $f(X, Y) :- \text{times}(X, X, Z), \text{minus}(Z, 5, V), \text{times}(7, V, U), \text{plus}(Z, U, Y).$

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ack (n+1) m = if m == 0 then ack n 1 else
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### Example (Ackermann function as logic program)

```
ack(0,M,s(M)).  
ack(s(N),M,R) :- =(M,0,B), cond(B,N,M,R).  
cond(true,N,M,R) :- ack(N,s(0),R).  
cond(false,N,M,R) :- -(M,s(0),U), ack(s(N),U,V), ack(N,V,R).
```

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can be programmed as

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f(X,Y) :- eval(s(X*X) - X*X, Y).
```

- evaluator is simple logic program (actually a simple **meta interpreter**)

```
eval(0,0).
```

```
eval(s(E),s(N)) :- eval(E,N).
```

```
eval(E+F,K) :- eval(E,N), eval(F,M), plus(N,M,K).
```

```
eval(E-F,K) :- eval(E,N), eval(F,M), plus(M,K,N).
```

```
eval(E*F,K) :- eval(E,N), eval(F,M), times(N,M,K).
```

Example (f(X,Y) :- eval(s(X\*X) - X\*X, Y).)

f(s(s(0)),Y)

Example (`f(X,Y) :- eval(s(X*X) - X*X, Y).`)

```
eval(s(s(s(0))*s(s(0))) - s(s(0))*s(s(0)),Y)
      |
f(s(s(0)),Y)
```

Example (`f(X,Y) :- eval(s(X*X) - X*X, Y).`)

```
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)
    |
eval(s(s(s(0))*s(s(0))) - s(s(0))*s(s(0)),Y)
    |
f(s(s(0)),Y)
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Example (`f(X,Y) :- eval(s(X*X) - X*X, Y).`)

```
eval(s(s(0))*s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N = s(N1) |
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)
    |
    eval(s(s(s(0))*s(s(0)))-s(s(0))*s(s(0)),Y)
    |
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Example (`f(X,Y) :- eval(s(X*X) - X*X, Y).`)

```
eval(s(s(0)),N2), eval(s(s(0)),N3), times(N2,N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
|
eval(s(s(0))*s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
  N = s(N1) |
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)
  |
eval(s(s(s(0))*s(s(0)))-s(s(0))*s(s(0)),Y)
  |
f(s(s(0)),Y)
```

Example ( $f(X, Y) :- eval(s(X*X) - X*X, Y).$ )

```
eval(s(0),N4), eval(s(s(0)),N3), times(s(N4),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    |
    N2 = s(N4) |
eval(s(s(0)),N2), eval(s(s(0)),N3), times(N2,N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    |
    eval(s(s(0))*s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    |
    N = s(N1) |
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)
    |
    eval(s(s(s(0))*s(s(0))) - s(s(0))*s(s(0)),Y)
    |
    f(s(s(0)),Y)
```

Example (`f(X,Y) :- eval(s(X*X) - X*X, Y).`)

```
eval(0,N5), eval(s(s(0)),N3), times(s(s(N5)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N4 = s(N5) |
eval(s(0),N4), eval(s(s(0)),N3), times(s(N4),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N2 = s(N4) |
eval(s(s(0)),N2), eval(s(s(0)),N3), times(N2,N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    |
    eval(s(s(0))*s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N = s(N1) |
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)
    |
    eval(s(s(s(0))*s(s(0))) - s(s(0))*s(s(0)),Y)
    |
f(s(s(0)),Y)
```

Example (`f(X,Y) :- eval(s(X*X) - X*X, Y).`)

```

eval(s(s(0)),N3), times(s(s(0)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
          N5 = 0 |
eval(0,N5), eval(s(s(0)),N3), times(s(s(N5)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
          N4 = s(N5) |
eval(s(0),N4), eval(s(s(0)),N3), times(s(N4),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
          N2 = s(N4) |
eval(s(s(0)),N2), eval(s(s(0)),N3), times(N2,N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
          |
          eval(s(s(0))*s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
          N = s(N1) |
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)
          |
          eval(s(s(s(0))*s(s(0))) - s(s(0))*s(s(0)),Y)
          |
          f(s(s(0)),Y)

```

Example ( $f(X, Y) :- eval(s(X*X) - X*X, Y).$ )

```

times(s(s(0)),s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N3 = s(s(0)) ||
eval(s(s(0)),N3), times(s(s(0)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N5 = 0 |
eval(0,N5), eval(s(s(0)),N3), times(s(s(N5)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N4 = s(N5) |
eval(s(0),N4), eval(s(s(0)),N3), times(s(N4),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N2 = s(N4) |
eval(s(s(0)),N2), eval(s(s(0)),N3), times(N2,N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    |
eval(s(s(0))*s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
    N = s(N1) |
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)
    |
eval(s(s(s(0))*s(s(0))) - s(s(0))*s(s(0)),Y)
    |
f(s(s(0)),Y)

```

Example ( $f(X, Y) :- eval(s(X*X) - X*X, Y).$ )

```

eval(s(s(0))*s(s(0)),M), plus(M,Y,s(s(s(s(0)))))

N1 = s(s(s(0))) ||
times(s(s(0)),s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))

N3 = s(s(0)) ||
eval(s(s(0)),N3), times(s(s(0)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))

N5 = 0 |
eval(0,N5), eval(s(s(0)),N3), times(s(s(N5)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))

N4 = s(N5) |
eval(s(0),N4), eval(s(s(0)),N3), times(s(N4),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))

N2 = s(N4) |
eval(s(s(0)),N2), eval(s(s(0)),N3), times(N2,N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))

| |
eval(s(s(0))*s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))

N = s(N1) |
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)

| |
eval(s(s(s(0))*s(s(0))) - s(s(0))*s(s(0)),Y)

| |
f(s(s(0)),Y)

```

Example ( $f(X, Y) :- eval(s(s(X*X)) - s(s(X*X)), Y).$ )

```

plus(s(s(s(s(0)))),Y,s(s(s(s(s(0))))))
M = s(s(s(s(0)))) ||
eval(s(s(0))*s(s(0)),M), plus(M,Y,s(s(s(s(s(0))))))
N1 = s(s(s(s(0)))) ||
times(s(s(0)),s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
N3 = s(s(0)) ||
eval(s(s(0)),N3), times(s(s(0)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
N5 = 0 |
eval(0,N5), eval(s(s(0)),N3), times(s(s(N5)),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
N4 = s(N5) |
eval(s(0),N4), eval(s(s(0)),N3), times(s(N4),N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
N2 = s(N4) |
eval(s(s(0)),N2), eval(s(s(0)),N3), times(N2,N3,N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
|
eval(s(s(0))*s(s(0)),N1), eval(s(s(0))*s(s(0)),M), plus(M,Y,s(N1))
N = s(N1) |
eval(s(s(s(0))*s(s(0))),N), eval(s(s(0))*s(s(0)),M), plus(M,Y,N)
|
eval(s(s(s(0))*s(s(0))) - s(s(0))*s(s(0)),Y)
|
f(s(s(0)),Y)

```

Example ( $f(X, Y) :- \text{eval}(s(X*X) - X*X, Y).$ )

```


$$\frac{}{Y = s(0) \parallel}
\frac{}{\text{plus}(s(s(s(s(0)))), Y, s(s(s(s(0)))))}
\frac{}{M = s(s(s(s(0)))) \parallel}
\frac{}{\text{eval}(s(s(0))*s(s(s(0))), M), \text{plus}(M, Y, s(s(s(s(0)))))}
\frac{}{N1 = s(s(s(s(0)))) \parallel}
\frac{}{\text{times}(s(s(0)), s(s(0)), N1), \text{eval}(s(s(0))*s(s(s(0))), M), \text{plus}(M, Y, s(N1))}
\frac{}{N3 = s(s(s(0))) \parallel}
\frac{}{\text{eval}(s(s(0)), N3), \text{times}(s(s(0)), N3, N1), \text{eval}(s(s(0))*s(s(s(0))), M), \text{plus}(M, Y, s(N1))}
\frac{}{N5 = 0 \mid}
\frac{}{\text{eval}(0, N5), \text{eval}(s(s(0)), N3), \text{times}(s(s(N5)), N3, N1), \text{eval}(s(s(0))*s(s(s(0))), M), \text{plus}(M, Y, s(N1))}
\frac{}{N4 = s(N5) \mid}
\frac{}{\text{eval}(s(0), N4), \text{eval}(s(s(0)), N3), \text{times}(s(N4), N3, N1), \text{eval}(s(s(0))*s(s(s(0))), M), \text{plus}(M, Y, s(N1))}
\frac{}{N2 = s(N4) \mid}
\frac{}{\text{eval}(s(s(0)), N2), \text{eval}(s(s(0)), N3), \text{times}(N2, N3, N1), \text{eval}(s(s(0))*s(s(s(0))), M), \text{plus}(M, Y, s(N1))}
\frac{}{|}{\text{eval}(s(s(0))*s(s(s(0))), N1), \text{eval}(s(s(0))*s(s(s(0))), M), \text{plus}(M, Y, s(N1))}
\frac{}{N = s(N1) \mid}
\frac{}{\text{eval}(s(s(s(0))*s(s(s(0))), N), \text{eval}(s(s(0))*s(s(s(0))), M), \text{plus}(M, Y, N)}
\frac{}{|}{\text{eval}(s(s(s(0))*s(s(s(0))), - s(s(s(0))*s(s(s(0))), Y)
\frac{}{|}{f(s(s(0)), Y)$$

```

Example ( $f(X, Y) :- \text{eval}(s(X*X) - X*X, Y).$ )

```

 $\square$ 
Y = s(0) ||
plus(s(s(s(s(0)))), Y, s(s(s(s(0))))))
M = s(s(s(s(0)))) ||
eval(s(s(0))*s(s(s(0))), M), plus(M, Y, s(s(s(s(0))))))
N1 = s(s(s(s(0)))) ||
times(s(s(0)), s(s(0)), N1), eval(s(s(0))*s(s(s(0))), M), plus(M, Y, s(N1))
N3 = s(s(0)) ||
eval(s(s(0)), N3), times(s(s(0)), N3, N1), eval(s(s(0))*s(s(s(0))), M), plus(M, Y, s(N1))
N5 = 0 |
eval(0, N5), eval(s(s(0)), N3), times(s(s(N5)), N3, N1), eval(s(s(0))*s(s(s(0))), M), plus(M, Y, s(N1))
N4 = s(N5) |
eval(s(0), N4), eval(s(s(0)), N3), times(s(N4), N3, N1), eval(s(s(0))*s(s(s(0))), M), plus(M, Y, s(N1))
N2 = s(N4) |
eval(s(s(0)), N2), eval(s(s(0)), N3), times(N2, N3, N1), eval(s(s(0))*s(s(s(0))), M), plus(M, Y, s(N1))
|
eval(s(s(0))*s(s(s(0))), N1), eval(s(s(0))*s(s(s(0))), M), plus(M, Y, s(N1))
N = s(N1) |
eval(s(s(s(0))*s(s(s(0))), N), eval(s(s(0))*s(s(s(0))), M), plus(M, Y, N)
|
eval(s(s(s(0))*s(s(s(0)))) - s(s(s(0))*s(s(s(0))), Y)
|
f(s(s(0)), Y)

```

## Speeding up evaluation using “let”

- consider sub-expression X\*X

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- `let(X,E,F)` encodes *let x = e in f*

`eval(0,0).`

`eval(s(E),s(N)) :- eval(E,N).`

`eval(E+F,K) :- eval(E,N), eval(F,M), plus(N,M,K).`

`eval(E-F,K) :- eval(E,N), eval(F,M), plus(M,K,N).`

`eval(E*F,K) :- eval(E,N), eval(F,M), times(N,M,K).`

`eval(let(X,E,F),K) :- eval(E,N), X = N, eval(F,K).`

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- consider sub-expression  $X*X$
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- `let(X,E,F)` encodes  $\text{let } x = e \text{ in } f$

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### Example

`f(X,Y) :- eval(s(X*X) - X*X, Y).`

`f(X,Y) :- eval(let(X2, X*X, s(X2) - X2), Y).`

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y).)

f(s(s(0)),Y)

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y).)

```
eval(let(X2,s(s(0))*s(s(0)),s(X2)-X2),Y)
      |
      f(s(s(0)),Y)
```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y).)

```
eval(s(s(0))*s(s(0)),N), X2 = N, eval(s(X2)-X2,Y)
|
eval(let(X2,s(s(0))*s(s(0)),s(X2)-X2),Y)
|
f(s(s(0)),Y)
```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y). )

```
X2 = s(s(s(s(0)))), eval(s(X2)-X2,Y)
      N = s(s(s(s(0)))) ||
eval(s(s(0))*s(s(0)),N), X2 = N, eval(s(X2)-X2,Y)
|
eval(let(X2,s(s(0))*s(s(0)),s(X2)-X2),Y)
|
f(s(s(0)),Y)
```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y).)

```
eval(s(s(s(s(s(0)))))-s(s(s(s(0)))),Y)
      X2 = s(s(s(s(0)))) |
X2 = s(s(s(s(0)))), eval(s(X2)-X2,Y)
      N = s(s(s(s(0)))) ||
eval(s(s(0))*s(s(0)),N), X2 = N, eval(s(X2)-X2,Y)
      |
eval(let(X2,s(s(0))*s(s(0)),s(X2)-X2),Y)
      |
f(s(s(0)),Y)
```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y).)

```
eval(s(s(s(s(s(0))))),N), eval(s(s(s(s(0)))),M), plus(M,Y,N)
    |
    eval(s(s(s(s(s(0)))))-s(s(s(s(0)))),Y)
        |
        X2 = s(s(s(s(0))))
        |
        X2 = s(s(s(s(0)))), eval(s(X2)-X2,Y)
            |
            N = s(s(s(s(0))) ||
eval(s(s(0))*s(s(s(0))),N), X2 = N, eval(s(X2)-X2,Y)
    |
    eval(let(X2,s(s(s(0)))*s(s(s(0))),s(X2)-X2),Y)
        |
        f(s(s(s(0))),Y)
```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y). )

```
eval(s(s(s(s(0)))),M), plus(M,Y,s(s(s(s(s(0))))))
      N = s(s(s(s(0)))) ||

eval(s(s(s(s(s(0)))),N), eval(s(s(s(s(0)))),M), plus(M,Y,N)
      |
      eval(s(s(s(s(s(0)))))-s(s(s(s(0)))),Y)
      X2 = s(s(s(s(0)))), eval(s(X2)-X2,Y)
      N = s(s(s(s(0)))) ||

eval(s(s(0))*s(s(0)),N), X2 = N, eval(s(X2)-X2,Y)
      |
      eval(let(X2,s(s(0))*s(s(0)),s(X2)-X2),Y)
      |
      f(s(s(0)),Y)
```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y). )

```
plus(s(s(s(s(0)))),Y,s(s(s(s(s(0))))))
M = s(s(s(s(0)))) ||
eval(s(s(s(s(0)))),M), plus(M,Y,s(s(s(s(s(0)))))))
N = s(s(s(s(s(0)))) ||
eval(s(s(s(s(s(0))))),N), eval(s(s(s(s(0)))),M), plus(M,Y,N)
|
eval(s(s(s(s(s(0)))))-s(s(s(s(0)))),Y)
X2 = s(s(s(s(0)))), eval(s(X2)-X2,Y)
N = s(s(s(s(0)))) ||
eval(s(s(0))*s(s(0)),N), X2 = N, eval(s(X2)-X2,Y)
|
eval(let(X2,s(s(0))*s(s(0)),s(X2)-X2),Y)
|
f(s(s(0)),Y)
```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y).)

```

    □
    Y = s(0) ||
plus(s(s(s(s(0)))),Y,s(s(s(s(s(0))))))
    M = s(s(s(0))) ||
eval(s(s(s(s(0)))),M), plus(M,Y,s(s(s(s(s(0)))))))
    N = s(s(s(s(0)))) ||
eval(s(s(s(s(s(0))))),N), eval(s(s(s(s(0)))),M), plus(M,Y,N)
    |
eval(s(s(s(s(s(0)))))-s(s(s(s(0)))),Y)
    X2 = s(s(s(s(0)))) |
X2 = s(s(s(s(0)))), eval(s(X2)-X2,Y)
    N = s(s(s(s(0)))) ||
eval(s(s(0))*s(s(0)),N), X2 = N, eval(s(X2)-X2,Y)
    |
eval(let(X2,s(s(0))*s(s(0)),s(X2)-X2),Y)
    |
f(s(s(0)),Y)

```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y).)

```

    □
    Y = s(0) ||
plus(s(s(s(s(0)))),Y,s(s(s(s(s(0))))))
    M = s(s(s(0))) ||
eval(s(s(s(s(0)))),M), plus(M,Y,s(s(s(s(s(0)))))))
    N = s(s(s(s(0)))) ||
eval(s(s(s(s(s(0))))),N), eval(s(s(s(s(0)))),M), plus(M,Y,N)
    |
eval(s(s(s(s(s(0)))))-s(s(s(s(0)))),Y)
    X2 = s(s(s(s(0)))) |
X2 = s(s(s(s(0)))), eval(s(X2)-X2,Y)
    N = s(s(s(0))) ||
eval(s(s(0))*s(s(0)),N), X2 = N, eval(s(X2)-X2,Y)
    |
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```

Example (f(X,Y) :- eval(let(X2,X\*X,s(X2)-X2), Y).)

```

    □
    Y = s(0) ||
plus(s(s(s(s(0)))),Y,s(s(s(s(s(0))))))
    M = s(s(s(0))) ||
eval(s(s(s(s(0)))),M), plus(M,Y,s(s(s(s(s(0)))))))
    N = s(s(s(s(0)))) ||
eval(s(s(s(s(s(0))))),N), eval(s(s(s(s(0)))),M), plus(M,Y,N)
    |
eval(s(s(s(s(s(0)))))-s(s(s(s(0)))),Y)
    X2 = s(s(s(s(0)))) |
X2 = s(s(s(s(0)))), eval(s(X2)-X2,Y)
    N = s(s(s(s(0)))) ||
eval(s(s(0))*s(s(0)),N), X2 = N, eval(s(X2)-X2,Y)
    |
eval(let(X2,s(s(0))*s(s(0)),s(X2)-X2),Y)
    |
f(s(s(0)),Y)

```

## Speeding up “let” even further

- detected problems:
  - 1 after computing  $x^2$ , result is evaluated again  
`eval(s(s(s(s(0)))) ,M)`
  - 2 eval also steps into **initial input**

## Speeding up “let” even further

- detected problems:
  - 1 after computing  $x^2$ , result is evaluated again  
`eval(s(s(s(s(0)))),M)`
  - 2 eval also steps into **initial input**
- solution: add new constructor *num* which states that the argument is a number, and hence, does not have to be evaluated

```
eval(0,0).  
eval(s(E),s(N)) :- eval(E,N).  
eval(E+F,K) :- eval(E,N), eval(F,M), plus(N,M,K).  
eval(E-F,K) :- eval(E,N), eval(F,M), plus(M,K,N).  
eval(E*F,K) :- eval(E,N), eval(F,M), times(N,M,K).  
eval(num(N),N).  
eval(let(X,E,F),K) :- eval(E,N), X = num(N), eval(F,K).
```

## Speeding up “let” even further

- detected problems:
  - 1 after computing  $x^2$ , result is evaluated again  
`eval(s(s(s(s(0)))),M)`
  - 2 eval also steps into **initial input**
- solution: add new constructor *num* which states that the argument is a number, and hence, does not have to be evaluated

```
eval(0,0).  
eval(s(E),s(N)) :- eval(E,N).  
eval(E+F,K) :- eval(E,N), eval(F,M), plus(N,M,K).  
eval(E-F,K) :- eval(E,N), eval(F,M), plus(M,K,N).  
eval(E*F,K) :- eval(E,N), eval(F,M), times(N,M,K).  
eval(num(N),N).  
eval(let(X,E,F),K) :- eval(E,N), X = num(N), eval(F,K).
```

Example `(f(X,Y) :- GX = num(X), eval(let(X2,GX*GX,s(X2)-X2),Y))`

`f(s(s(0)),Y)`

Example `(f(X,Y) :- GX = num(X), eval(let(X2,GX*GX,s(X2)-X2),Y))`

```
GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
|
f(s(s(0)),Y)
```

Example `(f(X,Y) :- GX = num(X), eval(let(X2,GX*GX,s(X2)-X2),Y))`

```
eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
     GX = num(s(s(0))) |
GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
                           |
                           f(s(s(0)),Y)
```

Example `(f(X,Y) :- GX = num(X), eval(let(X2,GX*GX,s(X2)-X2),Y))`

```
eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
|
eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
    |
    GX = num(s(s(0))) |
GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
    |
    f(s(s(0)),Y)
```

Example `(f(X,Y) :- GX = num(X), eval(let(X2,GX*GX,s(X2)-X2),Y))`

```
eval(num(s(s(0)),N1), eval(num(s(s(0)),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)
    |
    eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
    |
    eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
        |
        GX = num(s(s(0))) |
        GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
            |
            f(s(s(0)),Y)
```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```
eval(num(s(s(0)),N2), times(s(s(0)),N2,N), X2 = num(N), eval(s(X2)-X2,Y)
      N1 = s(s(0)) |
eval(num(s(s(0)),N1), eval(num(s(s(0)),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)
      |
      eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
      |
      eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
          GX = num(s(s(0))) |
          GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
          |
          f(s(s(0)),Y)
```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```
times(s(s(0)),s(s(0)),N), X2 = num(N), eval(s(X2)-X2,Y)
    N2 = s(s(0)) |
eval(num(s(s(0)),N2), times(s(s(0)),N2,N), X2 = num(N), eval(s(X2)-X2,Y)
    N1 = s(s(0)) |
eval(num(s(s(0)),N1), eval(num(s(s(0)),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)
    |
eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
    |
eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
    GX = num(s(s(0))) |
GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
    |
f(s(s(0)),Y)
```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```
X2 = num(s(s(s(s(0))))), eval(s(X2)-X2,Y)
      N = s(s(s(s(0)))) ||
times(s(s(0)),s(s(0)),N), X2 = num(N), eval(s(X2)-X2,Y)
      N2 = s(s(0)) |
eval(num(s(s(0)),N2), times(s(s(0)),N2,N), X2 = num(N), eval(s(X2)-X2,Y)
      N1 = s(s(0)) |
eval(num(s(s(0)),N1), eval(num(s(s(0)),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)
      |
eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
      |
eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
      GX = num(s(s(0))) |
GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
      |
f(s(s(0)),Y)
```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```

eval(s(num(s(s(s(s(0))))))-num(s(s(s(s(0))))),Y)
    X2 = num(s(s(s(s(0))))) |
    X2 = num(s(s(s(s(0))))), eval(s(X2)-X2,Y)
        N = s(s(s(s(0)))) ||
    times(s(s(0)),s(s(0)),N), X2 = num(N), eval(s(X2)-X2,Y)
        N2 = s(s(0)) |
    eval(num(s(s(0)),N2), times(s(s(0)),N2,N), X2 = num(N), eval(s(X2)-X2,Y)
        N1 = s(s(0)) |
    eval(num(s(s(0)),N1), eval(num(s(s(0)),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)
        |
    eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
        |
    eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
        GX = num(s(s(0))) |
    GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
        |
    f(s(s(0)),Y)

```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```

eval(s(num(s(s(s(s(0)))))),N), eval(num(s(s(s(s(0))))),M), plus(M,Y,N)
|
eval(s(num(s(s(s(s(0))))))-num(s(s(s(s(0))))),Y)
    X2 = num(s(s(s(s(0)))) |
    X2 = num(s(s(s(s(0))))), eval(s(X2)-X2,Y)
    N = s(s(s(s(0)))) ||
times(s(s(0)),s(s(0)),N), X2 = num(N), eval(s(X2)-X2,Y)
    N2 = s(s(0)) |
eval(num(s(s(0)),N2), times(s(s(0)),N2,N), X2 = num(N), eval(s(X2)-X2,Y)
    N1 = s(s(0)) |
eval(num(s(s(0)),N1), eval(num(s(s(0)),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)
|
eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
|
eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
    GX = num(s(s(0))) |
GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
|
f(s(s(0)),Y)

```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```

eval(num(s(s(s(s(0))))),N1), eval(num(s(s(s(s(0))))),M), plus(M,Y,s(N1))
    N = s(N1) |
eval(s(num(s(s(s(s(0))))),N), eval(num(s(s(s(s(0))))),M), plus(M,Y,N)
    |
eval(s(num(s(s(s(s(0))))))-num(s(s(s(s(0))))),Y)
    X2 = num(s(s(s(s(0)))))| 
    X2 = num(s(s(s(s(0))))), eval(s(X2)-X2,Y)
    N = s(s(s(s(0)))) ||
times(s(s(0)),s(s(0)),N), X2 = num(N), eval(s(X2)-X2,Y)
    N2 = s(s(0)) |
eval(num(s(s(0)),N2), times(s(s(0)),N2,N), X2 = num(N), eval(s(X2)-X2,Y)
    N1 = s(s(0)) |
eval(num(s(s(0)),N1), eval(num(s(s(0)),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)
    |
eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
    |
eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
    GX = num(s(s(0))) |
    GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
    |
f(s(s(0)),Y)

```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```

eval(num(s(s(s(s(0))))),M), plus(M,Y,s(s(s(s(s(0))))))
    N1 = s(s(s(s(0)))) |
eval(num(s(s(s(s(0))))),N1), eval(num(s(s(s(s(0))))),M), plus(M,Y,s(N1))
    N = s(N1) |
eval(s(num(s(s(s(s(0)))))),N), eval(num(s(s(s(s(0))))),M), plus(M,Y,N)
    |
eval(s(num(s(s(s(s(0))))))-num(s(s(s(s(0))))),Y)
    X2 = num(s(s(s(s(0)))))| 
    X2 = num(s(s(s(s(0))))), eval(s(X2)-X2,Y)
    N = s(s(s(s(0)))) ||
times(s(s(0)),s(s(0)),N), X2 = num(N), eval(s(X2)-X2,Y)
    N2 = s(s(0)) |
eval(num(s(s(0)),N2), times(s(s(0)),N2,N), X2 = num(N), eval(s(X2)-X2,Y)
    N1 = s(s(0)) |
eval(num(s(s(0))),N1), eval(num(s(s(0))),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)
    |
eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)
    |
eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
    GX = num(s(s(0))) |
GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)
    |
f(s(s(0)),Y)

```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```

plus(s(s(s(s(0)))),Y,s(s(s(s(s(0))))))

M = s(s(s(s(0)))) |
eval(num(s(s(s(s(0))))),M), plus(M,Y,s(s(s(s(s(0))))))

N1 = s(s(s(s(0)))) |
eval(num(s(s(s(s(0))))),N1), eval(num(s(s(s(s(0))))),M), plus(M,Y,s(N1))

N = s(N1) |
eval(s(num(s(s(s(s(0)))))),N), eval(num(s(s(s(s(0))))),M), plus(M,Y,N)

eval(s(num(s(s(s(s(0))))))-num(s(s(s(s(0))))),Y)
X2 = num(s(s(s(s(0)))))|

X2 = num(s(s(s(s(0))))), eval(s(X2)-X2,Y)
N = s(s(s(s(0))))||

times(s(s(0)),s(s(0)),N), X2 = num(N), eval(s(X2)-X2,Y)
N2 = s(s(0)) |

eval(num(s(s(0))),N2), times(s(s(0)),N2,N), X2 = num(N), eval(s(X2)-X2,Y)
N1 = s(s(0)) |

eval(num(s(s(0))),N1), eval(num(s(s(0))),N2), times(N1,N2,N), X2 = num(N), eval(s(X2)-X2,Y)

eval(num(s(s(0)))*num(s(s(0))),N), X2 = num(N), eval(s(X2)-X2,Y)

eval(let(X2,num(s(s(0)))*num(s(s(0))),s(X2)-X2),Y)
GX = num(s(s(0))) |

GX = num(s(s(0))), eval(let(X2,GX*GX,s(X2)-X2),Y)

f(s(s(0)),Y)

```

Example  $(f(X, Y) :- GX = \text{num}(X), \text{eval}(\text{let}(X2, GX*GX, s(X2)-X2), Y))$

```

 $\square$ 
Y = s(0) ||
plus(s(s(s(s(0)))), Y, s(s(s(s(s(0))))))
M = s(s(s(s(0)))) |
eval(num(s(s(s(s(0))))), M), plus(M, Y, s(s(s(s(s(0))))))

N1 = s(s(s(s(0)))) |
eval(num(s(s(s(s(0))))), N1), eval(num(s(s(s(s(0))))), M), plus(M, Y, N1)

N = s(N1) |
eval(s(num(s(s(s(s(0)))))), N), eval(num(s(s(s(s(0))))), M), plus(M, Y, N)
|
eval(s(num(s(s(s(s(0)))))-num(s(s(s(s(0))))), Y)
X2 = num(s(s(s(s(0)))) | 
X2 = num(s(s(s(s(0))))), eval(s(X2)-X2, Y)
N = s(s(s(s(0)))) ||
times(s(s(0)), s(s(0)), N), X2 = num(N), eval(s(X2)-X2, Y)
N2 = s(s(0)) |
eval(num(s(s(0))), N2), times(s(s(0)), N2, N), X2 = num(N), eval(s(X2)-X2, Y)
N1 = s(s(0)) |

eval(num(s(s(0))), N1), eval(num(s(s(0))), N2), times(N1, N2, N), X2 = num(N), eval(s(X2)-X2, Y)
|
eval(num(s(s(0)))*num(s(s(0))), N), X2 = num(N), eval(s(X2)-X2, Y)
|
eval(let(X2, num(s(s(0)))*num(s(s(0))), s(X2)-X2), Y)
GX = num(s(s(0))) |
GX = num(s(s(0))), eval(let(X2, GX*GX, s(X2)-X2), Y)
|
f(s(s(0)), Y)

```