

Logic Programming

Georg Moser

Department of Computer Science @ UIBK

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Summary of Last Lecture

Definitions

- a Prolog clause is called iterative if
 - 1 it has one recursive call, and
 - 2 zero or more calls to system predicates, before the recursive call
- a Prolog procedure is iterative if it contains only facts and iterative clauses

Observation

- iterative programs are tail recursive
- sometimes tail recursion in general can be implemented as iteration which doesn't require a stack

```
built_in (+,2).
user_def(fib,1).
:- eval(fib(13),N), N=233.
```

Outline of the Lecture

Monotone Logic Programs

introduction, basic constructs, logic foundations, unification, semantics, database and recursive programming, termination, complexity

Incomplete Data Structures and Constraints

incomplete data structures, definite clause grammars, constraint logic programming, answer set programming

Full Prolog

semantics (revisted), cuts, correctness proofs, meta-logical predicates, nondeterministic programming, pragmatics, efficient programs, meta programming

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Definition

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- a meta-interpreter for a language is an interpreter for the language written in the language itself
- for example, relation solve (Goal) is true, if *Goal* is true with respect to the program interpreted

```
Example (simple meta-interpreter)
solve(true).
solve((A,B)) :- solve(A), solve(B).
solve(A) :- clause(A,B), solve(B).
```

Meta-Program We Have Already Seen

```
Example
accept(S) :-
    initial(Q),
    accept(Q,S).
accept(Q,[X|Xs]) :-
    delta(Q,X,Q_1),
    accept(Q_1, Xs).
accept(Q,[]) :-
    final(0).
initial (q_0).
final(q_2).
delta(q_0, 0, q_0).
delta(q_0, 0, q_1).
delta(q_0,1,q_0).
delta(q_1, 1, q_2).
```

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delta(q_1, 1, q_2).
```

(Sort of) Meta-Program We'll See Soon

```
prove(and(A,B), UnExp, Lits, FreeV, VarLim): -!,
    prove (A, [B| UnExp], Lits, FreeV, VarLim).
prove(or(A,B), UnExp, Lits, FreeV, VarLim): -!,
    prove (A. UnExp. Lits, FreeV, VarLim),
    prove (B, UnExp, Lits, FreeV, VarLim).
prove(all(X,Fml),UnExp,Lits,FreeV,VarLim):-!,
    \+ length (FreeV, VarLim),
    copy_term ((X, Fml, FreeV), (X1, Fml1, FreeV)),
    append (UnExp.[all(X,Fml)], UnExp1).
    prove (Fml1, UnExp1, Lits, [X1 | FreeV], VarLim).
prove(Lit,_UnExp,[L|Lits],_FreeV,_VarLim) :- !,
    (Lit = neg Neg; neg Lit = Neg) \rightarrow
         (unify_with_occurs_check(Neg,L);
         prove(Lit,[], Lits, FreeV0, VarLim0)).
prove(Lit,[Next|UnExp], Lits, FreeV, VarLim) :- !,
    prove (Next, UnExp, [Lit | Lits], FreeV, VarLim).
```

```
Example (meta-interpreter with proofs)
solve(true,true).
solve((A,B),(ProofA,ProofB)) :-
    solve(A,ProofA),
    solve(B,ProofB).
solve(A,(A :- Proof)) :-
    clause(A,B),
    solve(B,Proof).
```

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solve(A,(A :- Proof)) :-
    clause(A,B),
    solve(B,Proof).
```

```
father(andreas,boris). female(doris). male(andreas).
father(andreas,christian). female(eva). male(boris).
father(andreas,doris). male(christian).
father(boris,eva). mother(doris,franz). male(franz).
father(franz,georg). mother(eva,georg). male(georg).
son(X,Y) :- father(Y,X), male(X).
```

```
Example (Tracing Pure Prolog) trace(Goal) :- trace(Goal,0).
```

```
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trace(Goal) :- trace(Goal,0).
trace(true,Depth).
trace((A,B),Depth) :-
    trace(A,Depth), trace(B,Depth).
trace(A,Depth) :-
    clause(A,B),
    display(A,Depth),
    Depth1 is Depth + 1,
    trace(B,Depth1).
```

```
Example (Tracing Pure Prolog)
trace(Goal) :- trace(Goal,0).
trace(true, Depth).
trace((A,B),Depth) :-
    trace(A,Depth), trace(B,Depth).
trace(A,Depth) :-
    clause(A,B),
    display(A, Depth),
    Depth1 is Depth + 1,
    trace(B,Depth1).
display(A,Depth) :- tab(Depth), write(A), nl.
```

```
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```

```
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trace(true,Depth) :- !.
```

```
system(A is B). system(read(X)). system(integer(X)).
system(clause(A,B)). system(A < B). system(A >= B).
system(write(X)). system(functor(T,F,N)). system(system(X)).
```

```
trace(Goal) :- trace(Goal,0).
trace(true,Depth) :- !.
trace((A,B),Depth) :- !, trace(A,Depth), trace(B,Depth).
```

```
trace(Goal) :- trace(Goal,0).
trace(true,Depth) :- !.
trace((A,B),Depth) :- !, trace(A,Depth), trace(B,Depth).
trace(A,Depth) :- system(A), A, !, display2(A,Depth), nl.
```

```
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trace(true,Depth) :- !.
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trace(A,Depth) :-
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system(A is B). system(read(X)). system(integer(X)).
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trace(Goal) :- trace(Goal,0).
trace(true,Depth) :- !.
trace((A,B),Depth) :- !, trace(A,Depth), trace(B,Depth).
trace(A,Depth) :- system(A), A, !, display2(A,Depth), nl.
trace(A,Depth) :-
    clause(A,B), display(A,Depth), nl,
    Depth1 is Depth + 1, trace(B,Depth1).
trace(A,Depth) :-
    \+ clause(A,B), display(A,Depth),
    tab(8), write(f),nl,fail.
```

```
trace(Goal) :- trace(Goal,0).
trace(true,Depth) :- !.
trace((A,B),Depth) :- !, trace(A,Depth), trace(B,Depth).
trace(A,Depth) :- system(A), A, !, display2(A,Depth), nl.
trace(A,Depth) :-
    clause(A,B), display(A,Depth), nl,
    Depth1 is Depth + 1, trace(B,Depth1).
trace(A,Depth) :-
    \+ clause(A,B), display(A,Depth),
    tab(8), write(f),nl,fail.
display(A,Depth) :- Spacing is 3*Depth, tab(Spacing), write(A).
```

Meta-Interpreters for Debugging

```
Example (Control Execution)
solve(true,_D,no_overflow):-
solve (A, 0, overflow([])).
solve((A,B),D,Overflow):-
        D > 0.
        solve (A, D, Overflow A),
         solve_conjunction (OverflowA, B, D, Overflow).
solve(A, D, no\_overflow) :-
        D > 0.
        system(A), !, A.
solve (A, D, Overflow) :-
        D > 0.
         clause (A,B),
        D1 is D-1,
         solve (B, D1, Overflow B),
         return_overflow (OverflowB, A, Overflow).
```

```
Example (Control Execution (cont'd))
solve_conjunction (overflow (S),_B,_D, overflow (S)).
solve_conjunction (no_overflow, B, D, Overflow) :-
         solve (B, D, Overflow).
return_overflow (no_overflow,_A, no_overflow).
return_overflow(overflow(S), A, overflow([A|S])).
\% isort(Xs,Ys) <— Ys is sorted Xs, using insertion sort
isort([X|Xs], Ys) := isort(Xs, Zs), my_insert(X, Zs, Ys).
isort ([],[]).
my_insert(X,[Y|Ys],[X,Y|Ys]) :-
        X < Y
my_insert(X, [Y|Ys], [Y|Zs]) :-
        X >= Y
        my_insert(X, [Y|Ys], Zs).
my_insert(X,[],[X]).
```

Expert Systems in Prolog

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expert systems typically consists of

- knowledge base
- inference engine

this separation is not suitable for a Prolog implementation

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Employ Meta-Interpreters

we implement the following features of expert systems using metainterpreters:

- interaction with the user
- explanation facility
- uncertainty reasoning

Toy Expert System

```
place_in_oven(Dish,top):-
    pastry (Dish), size (Dish, small).
place_in_oven(Dish, middle) :-
    pastry (Dish), size (Dish, big).
place_in_oven(Dish, middle) :-
    main_meal(Dish).
place_in_oven(Dish,low):-
    slow_cooker(Dish).
pastry(Dish) :- type(Dish, cake).
pastry(Dish) :- type(Dish, bread).
main_meal(Dish) :- type(Dish, meat).
slow_cooker(Dish) :- type(Dish, milk_pudding).
```

```
solve1/1
solve1(true):-
solve1((A,B)):-
        solve1(A), solve1(B).
solve1(A) :-
       A = (A1, A2),
        clause (A,B), solve1(B).
solve1(A) :-
        ask (A, Answer),
        respond (Answer, A).
ask(A, Answer) :- display_query(A), read(Answer).
askable(type(_Dish,_Type)).
askable(size(_Dish,_Size)).
respond(yes,A) :- assert(A).
respond(no,A) := assert(untrue(A)), fail.
```

Interaction (in the Naive)

```
interact(Goal) :-
          reset, solve1 (Goal).
reset :- retractall(type(_Dish,_Type)),
          retractall(size(_Sish,_Size)),
          retractall(untrue(_Fact)).
?— interact(place_in_oven(dish,X)).
type(dish, cake)? yes.
size (dish, small)? no.
type(dish, bread)? no.
size (dish, big)? yes.
X = middle
```

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Question

what about explanations for questions?

```
solve2/1
solve2(Goal) :- solve2(Goal,[]).
solve2(true,_Rules):-
solve2((A,B),Rules):-
         solve2(A, Rules), solve2(B, Rules).
solve2(A, Rules):-
        A = (A1, A2),
         clause (A,B).
         solve2(B,[rule(A,B)|Rules]).
solve2(A, Rules):-
         askable(A), \setminus + known(A),
         ask (A, Answer), respond (Answer, A, Rules).
respond (why, A, [Rule | Rules]) :-
         display_rule (Rule),
         ask (A, Answer),
         respond (Answer, A, Rules).
```

Interaction with Explanations

 $interact_why(Goal) :- reset, solve2(Goal).$

Interaction with Explanations

```
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```

```
?- interact_why(place_in_oven(dish,X)).
type(dish,cake)? yes.
size(dish,small)? no.
type(dish,bread)? no.
size(dish,big)? why.
if pastry(dish) and size(dish,big)
then place_in_oven(dish,middle)
size(dish,big)? yes.
X = middle
```

Interaction with Explanations

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```
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if pastry(dish) and size(dish,big)
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```

Question

how to obtain general explanations

```
interpret / 1
interpret ((Proof1, Proof2)) :-
        interpret(Proof1), interpret(Proof2).
interpret (Proof) :-
        fact (Proof, Fact),
        nl, write(Fact),
        writeln(' is a fact in the database').
interpret (Proof) :-
         rule (Proof, Head, Body, Proof1),
         nl, write (Head).
         writeln(' is proved using the rule'),
         display_rule (rule (Head, Body)),
         interpret (Proof1).
extract_body((Proof1, Proof2),(Body1, Body2)):-
         !, extract_body(Proof1, Body1),
         extract_body (Proof2, Body2).
extract_body((Goal <-- _Proof), Goal).
```

```
how/1
how(Goal) :- solve(Goal, Proof), interpret(Proof).
?— interact(place_in_oven(dish,X)).
% required for type and size of dish
?- how(place_in_oven(dish,top)).
place_in_oven(dish,top) is proved using the rule
if pastry(dish) and size(dish, small)
then place_in_oven(dish,top)
pastry(dish) is proved using the rule
if type (dish, bread)
then pastry (dish)
type(dish, bread) is a fact in the database
size (dish, small) is a fact in the database
```

• the explanation is exhaustive

Prolog computation is mirrored

 the explanation is exhaustive not intelligible for a knowledge base with 100 rules

Prolog computation is mirrored

- the explanation is exhaustive not intelligible for a knowledge base with 100 rules
- restrict explanation to one level:

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pastry (dish) can be further explained
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- Prolog computation is mirrored
- take expert knowledge into account:

```
interpret((Goal <-- Proof)) :-
  classification(Goal),
  write(Goal),
  writeln(' is a classification example').</pre>
```

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Exercise

Modify the implementation of $\ensuremath{\mathsf{how}}/1$ such that the partial answer proposed is generated

Definition

the certainty of a goal is computed as follows

$$\mathsf{cert}(G) = \begin{cases} \min\{\mathsf{cert}(A), \mathsf{cert}(B)\} & G = (A, B) \\ \max\{\mathsf{cert}(B) \cdot \mathit{Factor} \mid \mathsf{exists} \ \langle A : -B, \mathit{Factor} \rangle\} & G = A \end{cases}$$

Definition

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Definition (clauses with certification factor)

```
solve3/1
solve3(true,1):-
solve3((A,B),C):-
        solve3(A,C1),
        solve3(B,C2),
        minimum (C1, C2, C).
solve3(A,C):-
        clause_cf(A,B,C1),
        solve3 (B, C2),
        C is C1 * C2.
?— interact(place_in_oven(dish,X)).
% required for type and size of dish
?— solve3(place_in_oven(dish,top),C).
C = 0.7
```

Thank You for Your Attention!