

# Logic Programming

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# Summary of Last Lecture

NB: search trees are a tree representation of SLD-derivations

#### **Definitions**

- a type is a (possible infinite) set of terms
- types are conveniently defined by unary relations
- a type is complete if closed under the instance relation
- with every complete type T one associates an incomplete type IT
  which is a set of terms with instances in T and instances not in T

#### **Definitions**

- a list is complete if every instances satisfies the above type for lists
- otherwise it is incomplete

## Example

the lists [a,b,c] and [a,X,c] are complete; the list [a,b|Xs] is not

### Outline of the Lecture

# Monotone Logic Programs

introduction, basic constructs, logic foundations, unification, semantics, database and recursive programming, termination, complexity

# Incomplete Data Structures and Constraints

incomplete data structures, definite clause grammars, constraint logic programming, answer set programming

### Full Prolog

semantics (revisted), correctness proofs, meta-logical predicates, cuts nondeterministic programming, efficient programs, complexity

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- a proof tree for a program *P* and a goal *G* is a tree, whose nodes are goals and whose edges represent reduction of goals
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- a proof tree for a conjunction of goals  $G_1, \ldots, G_n$  is the set of proof trees for  $G_i$

#### **Definitions**

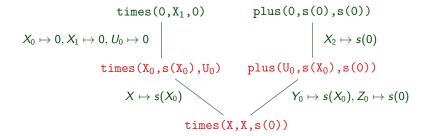
- a proof tree for a program *P* and a goal *G* is a tree, whose nodes are goals and whose edges represent reduction of goals
- the root is the query G
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#### Remark

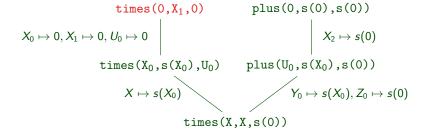
a proof tree is a different representation of one successful solution represented by a search tree combining all possible selection functions

```
\begin{array}{ll} \text{plus}(0,\textbf{X},\textbf{X})\,. & \text{times}(0,\textbf{X},0)\,. \\ \text{plus}(\textbf{s}(\textbf{X}),\textbf{Y},\textbf{s}(\textbf{Z}))\,:-\,\text{plus}(\textbf{X},\textbf{Y},\textbf{Z})\,. & \text{times}(\textbf{s}(\textbf{X}),\textbf{Y},\textbf{Z})\,:-\,\text{times}(\textbf{X},\textbf{Y},\textbf{U})\,, \\ & & \text{plus}(\textbf{U},\textbf{Y},\textbf{Z})\,. \end{array}
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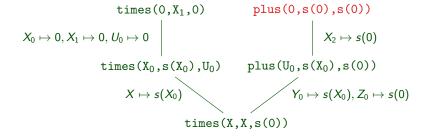
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```



```
\begin{array}{ll} plus(0,X,X)\,. & & times(0,X,0)\,. \\ plus(s(X),Y,s(Z)) :- plus(X,Y,Z)\,. & times(s(X),Y,Z) :- times(X,Y,U)\,, \\ & & plus(U,Y,Z)\,. \end{array}
```



```
\begin{array}{ll} plus(0,X,X)\,. & times(0,X,0)\,. \\ plus(s(X),Y,s(Z))\,:-\,plus(X,Y,Z)\,. & times(s(X),Y,Z)\,:-\,times(X,Y,U)\,, \\ & plus(U,Y,Z)\,. \end{array}
```



### Structured Data and Data Abstraction

Example (Unstructured Data)

course(discrete\_mathematics, tuesday, 8, 11, sandor, szedmak,
 victor\_franz\_hess, d).

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## Example (Structured Data)

```
course(discrete_mathematics,time(tuesday,8,11),
  lecturer(sandor,szedmak),location(victor_franz_hess,d)).
```

### Structured Data and Data Abstraction

# Example (Unstructured Data)

course(discrete\_mathematics, tuesday, 8, 11, sandor, szedmak,
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# Example (Structured Data)

```
course(discrete_mathematics,time(tuesday,8,11),
  lecturer(sandor,szedmak),location(victor_franz_hess,d)).
```

```
lecturer(Lecturer,Course) : -
  course(Course,Time,Lecturer,Location).
duration(Course,Length) : -
  course(Course,time(Day,Start,Finish),Lecturer,Location),
  plus(Start,Length,Finish).
```

# Example (cont'd)

```
teaches(Lecturer,Day) : -
  course(Course,time(Day,Start,Finish),Lecturer,Location).
occupied(Location,Day,Time) : -
  course(Course,time(Day,Start,Finish),Lecturer,Location),
  Start 
  Time, Time 
  Finish.
```

NB: rules for comparision are as expected

# Example (cont'd)

```
teaches(Lecturer,Day) : -
  course(Course,time(Day,Start,Finish),Lecturer,Location).
occupied(Location,Day,Time) : -
  course(Course,time(Day,Start,Finish),Lecturer,Location),
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  Time, Time 
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```

NB: rules for comparision are as expected

### Why structure Data?

- helps to organise data; databases are usually structured . . .
- rules can be written abstractly, hiding irrelevant detail
- modularity becomes possible or is improved

# Logic Programs and the Relational Database Model

#### Observation

the basic operations of relational algebras, namely:

- 1 union
- 2 difference
- 3 cartesian product
- 4 projection
- 5 selection
- 6 intersection

can easily be expressed within logic programming

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#### Observation

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```
r_{union_s}(X_1,...,X_n) := r(X_1,...,X_n).

r_{union_s}(X_1,...,X_n) := s(X_1,...,X_n).
```

```
Example (Type Condition)
is_number(0).
is_number(s(X)) : - is_number(X).
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```
\begin{array}{lll} & \texttt{plus}(0,\textbf{X},\textbf{X}) : - & \texttt{is\_number}(\textbf{X}) \,. \\ & \texttt{plus}(\textbf{s}(\textbf{X}),\textbf{Y},\textbf{s}(\textbf{Z})) : - & \texttt{plus}(\textbf{X},\textbf{Y},\textbf{Z}) \,. \\ & \texttt{times}(\textbf{0},\textbf{X},\textbf{0}) \,. \\ & \texttt{times}(\textbf{s}(\textbf{X}),\textbf{Y},\textbf{Z}) : - & \texttt{times}(\textbf{X},\textbf{Y},\textbf{U}) \,, & \texttt{plus}(\textbf{U},\textbf{Y},\textbf{Z}) \,. \end{array}
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Example (Type Condition)
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\begin{split} & \text{plus}(0,X,X) \,. \\ & \text{plus}(s(X),Y,s(Z)) \,:-\, \text{plus}(X,Y,Z) \,. \\ & \text{times}(0,X,0) \,. \\ & \text{times}(s(X),Y,Z) \,:-\, \text{times}(X,Y,U) \,,\, \text{plus}(U,Y,Z) \,. \end{split}
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### Example

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```

```
factorial(0,s(0)). factorial(s(N),F): - factorial(N,F<sub>1</sub>), times(s(N),F<sub>1</sub>,F).
```

```
\begin{array}{l} 0 \, \leqslant \, X \, :- \, \, \text{is.number(X).} \\ s(X) \, \leqslant \, s(Y) \, :- \, X \, \leqslant \, Y. \\ \\ \text{minimum(N1,N2,N1)} \, :- \, N_1 \, \leqslant \, N_2. \\ \\ \text{minimum(N1,N2,N2)} \, :- \, N_2 \, \leqslant \, N_1. \end{array}
```

```
\begin{array}{l} \texttt{0} \; \leqslant \; \texttt{X} . \\ \\ \texttt{s}(\texttt{X}) \; \leqslant \; \texttt{s}(\texttt{Y}) \; :- \; \texttt{X} \; \leqslant \; \texttt{Y} . \\ \\ \texttt{minimum}(\texttt{N}_1,\texttt{N}_2,\texttt{N}_1) \; :- \; \texttt{N}_1 \; \leqslant \; \texttt{N}_2 . \\ \\ \texttt{minimum}(\texttt{N}_1,\texttt{N}_2,\texttt{N}_2) \; :- \; \texttt{N}_2 \; \leqslant \; \texttt{N}_1 \, . \end{array}
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\begin{array}{l} 0 \;\leqslant\; X \,. \\ \text{s}(X) \;\leqslant\; \text{s}(Y) \;:-\; X \;\leqslant\; Y \,. \\ \\ \text{minimum}(N_1,N_2,N_1) \;:-\; N_1 \;\leqslant\; N_2 \,. \\ \\ \\ \text{minimum}(N_1,N_2,N_2) \;:-\; N_2 \;\leqslant\; N_1 \,. \end{array}
```

### Example

mod(X,Y,Z) : -Z < Y, times(Y,Q,W), plus(W,Z,X).

```
\begin{array}{l} 0 \;\leqslant\; X \,. \\  \text{s(X)} \;\leqslant\; \text{s(Y)} \;:-\; X \;\leqslant\; Y \,. \\  \text{minimum}(\text{N}_1,\text{N}_2,\text{N}_1) \;:-\; \text{N}_1 \;\leqslant\; \text{N}_2 \,. \\  \text{minimum}(\text{N}_1,\text{N}_2,\text{N}_2) \;:-\; \text{N}_2 \;\leqslant\; \text{N}_1 \,. \end{array}
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```
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```

```
\begin{array}{l} 0 \;\leqslant\; X \,. \\ s(X) \;\leqslant\; s(Y) \;:-\; X \;\leqslant\; Y \,. \\ \\ \text{minimum} \left(N_1\,,N_2\,,N_1\right) \;:-\; N_1 \;\leqslant\; N_2\,. \\ \\ \text{minimum} \left(N_1\,,N_2\,,N_2\right) \;:-\; N_2 \;\leqslant\; N_1\,. \end{array}
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```
\begin{split} & \text{mod}(X,Y,Z) \; := \; Z \; < \; Y, \; \text{times}(Y,\mathbb{Q},\mathbb{W}) \,, \; \text{plus}(\mathbb{W},Z,X) \,. \\ & \text{mod}(X,Y,X) \; := \; X \; < \; Y. \\ & \text{mod}(X,Y,Z) \; := \; \text{plus}(X1,Y,X) \,, \; \text{mod}(X1,Y,Z) \,. \end{split}
```

```
ackermann(0,N,s(N)).
```

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### Example

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\begin{split} & \operatorname{mod}(X,Y,Z) \,:\, - \, Z \,<\, Y, \, \operatorname{times}(Y,\mathbb{Q},\mathbb{W}) \,, \, \operatorname{plus}(\mathbb{W},Z,\mathbb{X}) \,. \\ & \operatorname{mod}(X,Y,\mathbb{X}) \,:\, - \, X \,<\, Y. \\ & \operatorname{mod}(X,Y,Z) \,:\, - \, \operatorname{plus}(X1,Y,\mathbb{X}) \,, \, \operatorname{mod}(X1,Y,Z) \,. \end{split}
```

```
ackermann(0,N,s(N)).

ackermann(s(M),0,Val) :- ackermann(M,s(0),Val).
```

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```

```
\begin{split} & \texttt{member}(\texttt{X}, \texttt{[X|Xs]}) \, . \\ & \texttt{member}(\texttt{X}, \texttt{[Y|Xs]}) \, :- \, \texttt{member}(\texttt{X}, \texttt{Xs}) \, . \end{split}
```

```
\label{eq:member} \begin{split} & \texttt{member}(\texttt{X}, [\texttt{X}|\texttt{Xs}]) \, . \\ & \texttt{member}(\texttt{X}, [\texttt{Y}|\texttt{Xs}]) \, :- \, \texttt{member}(\texttt{X}, \texttt{Xs}) \, . \\ & :- \, \texttt{member}(\texttt{X}, [\texttt{a}, \texttt{b}, \texttt{a}]) \, . \end{split}
```

```
\label{eq:member} \begin{split} & \texttt{member}(\texttt{X}, [\texttt{X}|\texttt{X}\texttt{s}]) \; . \\ & \texttt{member}(\texttt{X}, [\texttt{Y}|\texttt{X}\texttt{s}]) \; :- \; \texttt{member}(\texttt{X}, \texttt{X}\texttt{s}) \, . \\ & :- \; \texttt{member}(\texttt{X}, [\texttt{a}, \texttt{b}, \texttt{a}]) \, . \end{split}
```

```
append(Xs,Ys,Zs) : -
    Xs = [],
append(Xs,Ys,Zs) : -
    Xs = [H|Ts],
```

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\label{eq:member} \begin{split} & \texttt{member}(\texttt{X}, [\texttt{X}|\texttt{X}\texttt{s}]) \: . \\ & \texttt{member}(\texttt{X}, [\texttt{Y}|\texttt{X}\texttt{s}]) \: : \: - \: \texttt{member}(\texttt{X}, \texttt{X}\texttt{s}) \: . \end{split} \\ & : \: - \: \texttt{member}(\texttt{X}, [\texttt{a}, \texttt{b}, \texttt{a}]) \: . \end{split}
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```
append(Xs,Ys,Zs) : -
    Xs = [],
    Zs = Ys.
append(Xs,Ys,Zs) : -
    Xs = [H|Ts],
    append(Ts,Ys,Us),
    Zs = [H|Us].
```

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\label{eq:member} \begin{split} & \texttt{member}(\texttt{X}, [\texttt{X}|\texttt{X}\texttt{s}]) \, . \\ & \texttt{member}(\texttt{X}, [\texttt{Y}|\texttt{X}\texttt{s}]) \, :- \, \texttt{member}(\texttt{X}, \texttt{X}\texttt{s}) \, . \\ & :- \, \texttt{member}(\texttt{X}, [\texttt{a}, \texttt{b}, \texttt{a}]) \, . \end{split}
```

#### Example

```
append(Xs,Ys,Zs) :- append([],Ys,Ys).
    Xs = [], append([H|Ts],Ys,[H|Zs]) :-
    Zs = Ys. append(Ts,Ys,Zs).

append(Xs,Ys,Zs) :-
    Xs = [H|Ts],
    append(Ts,Ys,Us),
    Zs = [H|Us].
```

five steps to implement relation R

 $\blacksquare$  look up existing definitions of relation R

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#### five steps to implement relation R

- $\blacksquare$  look up existing definitions of relation R
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- $\blacksquare$  look up existing definitions of relation R
  - family relations
  - train tables
- 2 define types of individual data
  - is\_number
  - · mainly for documentation
- think up a suitable name
  - convert a verbose description into a name
  - child\_of
- 4 write queries (use cases)
- 5 write the actual program

```
prefix(Xs,Ys) : - append(Xs,As,Ys).
suffix(Xs,Ys) : - append(As,Xs,Ys).
member(X,Ys) : - append(As,[X|Xs],Ys).
```

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prefix(Xs,Ys) : - append(Xs,As,Ys).
suffix(Xs,Ys) : - append(As,Xs,Ys).
member(X,Ys) : - append(As,[X|Xs],Ys).
```

```
reverse([],[]).
reverse([X|Xs],Zs) : - reverse(Xs,Ys), append(Ys,[X],Zs).
```

```
prefix(Xs,Ys) : - append(Xs,As,Ys).
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reverse([],[]).
reverse([X|Xs],Zs) : - reverse(Xs,Ys), append(Ys,[X],Zs).
reverse(Xs,Ys) : - reverse(Xs,[],Ys).
reverse([X|Xs],Acc,Ys) : - reverse(Xs,[X|Acc],Ys).
reverse([],Ys,Ys).
```

```
prefix(Xs,Ys) : - append(Xs,As,Ys).
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#### Example

```
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reverse([X|Xs],Zs) : - reverse(Xs,Ys), append(Ys,[X],Zs).
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reverse([X|Xs],Acc,Ys) : - reverse(Xs,[X|Acc],Ys).
reverse([],Ys,Ys).
```

```
length([],0).
length([X|Xs],s(N)) : - length(Xs,N).
```

```
\begin{split} & \texttt{select}(\texttt{X}, [\texttt{X}|\texttt{Xs}], \texttt{Xs}) \,. \\ & \texttt{select}(\texttt{X}, [\texttt{Y}|\texttt{Ys}], [\texttt{Y}|\texttt{Zs}]) \,:-\, \texttt{select}(\texttt{X}, \texttt{Ys}, \texttt{Zs}) \,. \end{split}
```

```
permutationsort(Xs,Ys) : - permutation(Xs,Ys), ordered(Ys).
```

```
\begin{split} & \texttt{select}(\texttt{X}, [\texttt{X}|\texttt{Xs}], \texttt{Xs}) \,. \\ & \texttt{select}(\texttt{X}, [\texttt{Y}|\texttt{Ys}], [\texttt{Y}|\texttt{Zs}]) \,:-\, \texttt{select}(\texttt{X}, \texttt{Ys}, \texttt{Zs}) \,. \end{split}
```

```
\label{eq:permutation} \begin{split} & \text{permutationsort(Xs,Ys)} : - \text{ permutation(Xs,Ys), ordered(Ys).} \\ & \text{permutation(Xs,[Z|Zs])} : - \text{ select(Z,Xs,Ys), permutation(Ys,Zs).} \\ & \text{permutation([],[]).} \end{split}
```

```
select(X,[X|Xs],Xs).

select(X,[Y|Ys],[Y|Zs]) : - select(X,Ys,Zs).
```

```
\begin{split} & \text{permutationsort}(Xs,Ys) : - \text{ permutation}(Xs,Ys), \text{ ordered}(Ys). \\ & \text{permutation}(Xs,[Z|Zs]) : - \text{ select}(Z,Xs,Ys), \text{ permutation}(Ys,Zs). \\ & \text{permutation}([],[]). \\ & \text{ordered}([X]). \\ & \text{ordered}([X,Y|Ys]) : - X \leqslant Y, \text{ ordered}([Y|Ys]). \\ & \text{select}(X,[X|Xs],Xs). \\ & \text{select}(X,[Y|Ys],[Y|Zs]) : - \text{ select}(X,Ys,Zs). \end{split}
```

```
Example (Permutation Sort)  \begin{array}{lll} & \text{permutationsort}(Xs,Ys) : - \text{ permutation}(Xs,Ys), \text{ ordered}(Ys). \\ & \text{permutation}(Xs,[Z|Zs]) : - \text{ select}(Z,Xs,Ys), \text{ permutation}(Ys,Zs). \\ & \text{permutation}([],[]). \\ & \text{ordered}([X]). \\ & \text{ordered}([X,Y|Ys]) : - X \leqslant Y, \text{ ordered}([Y|Ys]). \\ & \text{select}(X,[X|Xs],Xs). \\ & \text{select}(X,[Y|Ys],[Y|Zs]) : - \text{ select}(X,Ys,Zs). \\ \end{array}
```

```
Example (Permutation Sort)  \begin{array}{lll} & \text{permutationsort}(Xs,Ys) : - \text{ permutation}(Xs,Ys), \text{ ordered}(Ys). \\ & \text{permutation}(Xs,[Z|Zs]) : - \text{ select}(Z,Xs,Ys), \text{ permutation}(Ys,Zs). \\ & \text{permutation}([],[]). \\ & \text{ordered}([X]). \\ & \text{ordered}([X,Y|Ys]) : - X \leqslant Y, \text{ ordered}([Y|Ys]). \\ & \text{select}(X,[X|Xs],Xs). \end{array}
```

select(X,[Y|Ys],[Y|Zs]) : - select(X,Ys,Zs).

```
Example (Quick Sort)
quicksort([X|Xs],Ys) : -
   partition(Xs,X,Littles,Bigs),
   quicksort(Littles,Ls), quicksort(Bigs,Rs),
   append(Ls,[X|Rs],Ys).
```

```
Example (Quick Sort)
quicksort([X|Xs],Ys) : -
    partition(Xs,X,Littles,Bigs),
    quicksort(Littles,Ls), quicksort(Bigs,Rs),
    append(Ls,[X|Rs],Ys).

partition([X|Xs],Y,[X|Ls],Bs) : -
    X =< Y, partition(Xs,Y,Ls,Bs).
partition([X|Xs],Y,Ls,[X|Bs]) : -
    X > Y, partition(Xs,Y,Ls,Bs).
partition([],Y,[],[]).
```

```
Example (Quick Sort)
 quicksort([X|Xs],Ys):-
     partition(Xs,X,Littles,Bigs),
     quicksort(Littles, Ls), quicksort(Bigs, Rs),
     append(Ls, [X|Rs], Ys).
 partition([X|Xs],Y,[X|Ls],Bs):-
     X =< Y, partition(Xs,Y,Ls,Bs).</pre>
 partition([X|Xs],Y,Ls,[X|Bs]):-
     X > Y, partition(Xs, Y, Ls, Bs).
 partition([],Y,[],[]).
```

```
isotree(nil,nil).
isotree(tree(X,Left1,Right1),tree(X,Left2,Right2)) : -
    isotree(Left1,Left2), isotree(Right1,Right2).
isotree(tree(X,Left1,Right1),tree(X,Left2,Right2)) : -
    isotree(Left1,Right2), isotree(Right1,Left2).
```