

This exam consists of **four** exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

[25] 1 Consider the two Haskell functions:

```

[] ++ ys      = ys
(x:xs) ++ ys = x : (xs ++ ys)

```

```

concat []      = []
concat (x:xs) = x ++ concat xs

```

Use structural induction to prove that `concat (xs ++ ys) = concat xs ++ concat ys` for all lists `xs` and `ys`. Apart from the defining equations above, you may assume associativity of `++`.

- [6] (a) Prove the base case, applying a single equation at a time.
 [9] (b) State the induction hypothesis and the statement you have to show in the step case.
 [10] (c) Prove the step case, applying a single equation at a time.

[25] 2 Consider the four λ -terms $A = \lambda f x. f x$, $K = \lambda x y. x$, $\Omega = (\lambda x. x x) (\lambda x. x x)$, and y .

- [9] (a) Use normal order reduction to reduce $A K y \Omega$ to normal form, giving each intermediate β -step.
 [8] (b) State which of A , K , Ω , and y from above are in weak head normal form and which are not.
 [8] (c) Determine for each of the following λ -terms if they are α -equivalent to A , K , Ω , and y , respectively.

$$t_1 = y \quad t_2 = \lambda x y. y \quad t_3 = \lambda g y. g y \quad t_4 = \lambda y x. y$$

[25] 3 Consider the Haskell functions:

```

diffs []      = 0
diffs (x:xs) = diffs xs - x

```

```

rev acc [] = acc
rev acc (x:xs) = rev (x:acc) xs

```

```

downto m n = if m < n then [] else m : downto (m - 1) n

```

- [9] (a) Use equational reasoning to evaluate `diffs [1,2]`, `rev [1,2,3] [1,2,3]`, and `downto 1 2`.
 [6] (b) Determine for each of the above functions whether it is tail recursive and/or guardedly recursive.
 [10] (c) Give a tail recursive implementation of `downto`.

[25] 4 Consider the typing environment $E = \{1 :: \text{Int}\}$.

- [10] (a) Use type checking to prove the typing judgment $E \vdash (\lambda x y. 1) :: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$.
 [6] (b) Apply the typing constraint rules to transform $E \triangleright \lambda x y. 1 :: \alpha_0$ into a unification problem.
 [9] (c) Solve the unification problem $\alpha_0 \approx \alpha_1 \rightarrow \alpha_2; \alpha_2 \approx \alpha_3 \rightarrow \alpha_4; \text{Int} \approx \alpha_4$ and compute the resulting mgu.