

All Solutions

2. Let $P' = P \cup \{x :: \text{Int}\}$.

1	$\text{True} :: \text{Bool}$	$\text{ins } P'$
2	$x :: \text{Int}$	$\text{ins } P'$
3	$+ :: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$	$\text{ins } P'$
4	$- :: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$	$\text{ins } P'$
5	$1 :: \text{Int}$	$\text{ins } P'$
6	$(+) x :: \text{Int} \rightarrow \text{Int}$	$\text{app } 3, 2$
7	$x + 1 :: \text{Int}$	$\text{app } 6, 5$
8	$(-) x :: \text{Int} \rightarrow \text{Int}$	$\text{app } 4, 2$
9	$x - 1 :: \text{Int}$	$\text{app } 8, 5$
10	$\text{if True then } x + 1 \text{ else } x - 1 :: \text{Int}$	$\text{ite } 1, 7, 9$

3.	1	$x :: \alpha_0$	assumption
	2	$y :: \alpha_1$	assumption
	3	$x :: \alpha_0$	copy 1
	4	$\lambda y. x :: \alpha_1 \rightarrow \alpha_0$	abs 2-3
	5	$\lambda xy. x :: \alpha_0 \rightarrow \alpha_1 \rightarrow \alpha_0$	abs 1-4

4.

$$\begin{aligned}
 \text{Pair}(\text{Bool}, \alpha_0) \approx \text{Pair}(\alpha_1, \text{Int}) &\stackrel{(d_1)}{\Rightarrow} \{\} && \text{Bool} \approx \alpha_1; \alpha_0 \approx \text{Int} \\
 &\stackrel{(v_2)}{\Rightarrow} \{\alpha_1/\text{Bool}\} && \alpha_0 \approx \text{Int} \\
 &\stackrel{(v_1)}{\Rightarrow} \{\alpha_0/\text{Int}\} && \square
 \end{aligned}$$

Hence the solution is the composition of $\{\alpha_1/\text{Bool}\}$ and $\{\alpha_0/\text{Int}\}$, that is,

$$\sigma = \{\alpha_0/\text{Int}, \alpha_1/\text{Bool}\}$$

5. The only possible first step is (d_1) , resulting in

$$\text{Bool} \approx \alpha_0; \alpha_0 \approx \text{Int}$$

Then we can either apply (v_2) on the first equality or (v_1) on the second one. In both cases the result is

$$\text{Bool} \approx \text{Int}$$

We have not reached \square yet, but still no further unification rule is applicable. Hence, the initial problem is not unifiable.

6. First we turn the resulting typing inference problem into a unification problem:

$$\begin{aligned}
& \frac{P \triangleright \mathbf{let} \text{ suc} = \lambda x. x + 1 \mathbf{ in} \ \mathbf{let} \ d = \lambda x. \text{ suc} (\text{ suc } x) \mathbf{ in} \ d \ 2 :: \alpha_0}{\Rightarrow^{(\text{let})} P \triangleright \lambda x. x + 1 :: \alpha_1; P, \text{ suc} :: \alpha_1 \triangleright \mathbf{let} \ d = \lambda x. \text{ suc} (\text{ suc } x) \mathbf{ in} \ d \ 2 :: \alpha_0} \\
& \Rightarrow^{(\text{let})} \frac{P \triangleright \lambda x. x + 1 :: \alpha_1; P, \text{ suc} :: \alpha_1 \triangleright \lambda x. \text{ suc} (\text{ suc } x) :: \alpha_2; P, \text{ suc} :: \alpha_1, d :: \alpha_2 \triangleright d \ 2 :: \alpha_0}{\Rightarrow^{(\text{asb})^2} \alpha_1 \approx \alpha_3 \rightarrow \alpha_4; \alpha_2 \approx \alpha_5 \rightarrow \alpha_6;} \\
& \frac{P, x :: \alpha_3 \triangleright x + 1 :: \alpha_4; P, \text{ suc} :: \alpha_1, x :: \alpha_5 \triangleright \text{ suc} (\text{ suc } x) :: \alpha_6; P, \text{ suc} :: \alpha_1, d :: \alpha_2 \triangleright d \ 2 :: \alpha_0}{\Rightarrow^{(\text{app})^3} \alpha_1 \approx \alpha_3 \rightarrow \alpha_4; \alpha_2 \approx \alpha_5 \rightarrow \alpha_6;} \\
& \quad P, x :: \alpha_3 \triangleright (+) \ x :: \alpha_7 \rightarrow \alpha_4; P, x :: \alpha_3 \triangleright 1 :: \alpha_7; \\
& \quad \frac{P, \text{ suc} :: \alpha_1, x :: \alpha_5 \triangleright \text{ suc} :: \alpha_8 \rightarrow \alpha_6; P, \text{ suc} :: \alpha_1, x :: \alpha_5 \triangleright \text{ suc } x :: \alpha_8;}{P, \text{ suc} :: \alpha_1, d :: \alpha_2 \triangleright d :: \alpha_9 \rightarrow \alpha_0; P, \text{ suc} :: \alpha_1, d :: \alpha_2 \triangleright 2 :: \alpha_9} \\
& \Rightarrow^{(\text{con})^4} \alpha_1 \approx \alpha_3 \rightarrow \alpha_4; \alpha_2 \approx \alpha_5 \rightarrow \alpha_6; \mathbf{Int} \approx \alpha_7; \alpha_1 \approx \alpha_8 \rightarrow \alpha_6; \alpha_2 \approx \alpha_9 \rightarrow \alpha_0; \mathbf{Int} \approx \alpha_9 \\
& \quad \frac{P, x :: \alpha_3 \triangleright (+) \ x :: \alpha_7 \rightarrow \alpha_4; P, \text{ suc} :: \alpha_1, x :: \alpha_5 \triangleright \text{ suc } x :: \alpha_8;}{\Rightarrow^{(\text{con})^2} \alpha_1 \approx \alpha_3 \rightarrow \alpha_4; \alpha_2 \approx \alpha_5 \rightarrow \alpha_6; \mathbf{Int} \approx \alpha_7; \alpha_1 \approx \alpha_8 \rightarrow \alpha_6; \alpha_2 \approx \alpha_9 \rightarrow \alpha_0; \mathbf{Int} \approx \alpha_9} \\
& \quad \frac{P, x :: \alpha_3 \triangleright (+) :: \alpha_{10} \rightarrow \alpha_7 \rightarrow \alpha_4; P, x :: \alpha_3 \triangleright x :: \alpha_{10};}{P, \text{ suc} :: \alpha_1, x :: \alpha_5 \triangleright \text{ suc} :: \alpha_{11} \rightarrow \alpha_8; P, \text{ suc} :: \alpha_1, x :: \alpha_5 \triangleright x :: \alpha_{11};} \\
& \Rightarrow^{(\text{con})^4} \alpha_1 \approx \alpha_3 \rightarrow \alpha_4; \alpha_2 \approx \alpha_5 \rightarrow \alpha_6; \mathbf{Int} \approx \alpha_7; \alpha_1 \approx \alpha_8 \rightarrow \alpha_6; \alpha_2 \approx \alpha_9 \rightarrow \alpha_0; \mathbf{Int} \approx \alpha_9; \\
& \quad \mathbf{Int} \rightarrow \mathbf{Int} \rightarrow \mathbf{Int} \approx \alpha_{10} \rightarrow \alpha_7 \rightarrow \alpha_4; \alpha_3 \approx \alpha_{10}; \alpha_1 \approx \alpha_{11} \rightarrow \alpha_8; \alpha_5 \approx \alpha_{11}
\end{aligned}$$

The resulting most general unifier is

$$\{\alpha_0 \mapsto \text{Int}, \alpha_1 \mapsto \text{Int} \rightarrow \text{Int}, \alpha_2 \mapsto \text{Int} \rightarrow \text{Int}, \alpha_3 \mapsto \text{Int}, \alpha_4 \mapsto \text{Int}, \alpha_5 \mapsto \text{Int}, \\ \alpha_6 \mapsto \text{Int}, \alpha_7 \mapsto \text{Int}, \alpha_8 \mapsto \text{Int}, \alpha_9 \mapsto \text{Int}, \alpha_{10} \mapsto \text{Int}, \alpha_{11} \mapsto \text{Int}\}$$