

Functional Programming

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Lecture 4



Topics

abstract data types, algebraic data types, binary search trees, combinator parsing, efficiency, encoding data types as lambda-terms, evaluation strategies, formal verification, first steps, guarded recursion, Haskell introduction, higher-order functions, historical overview, implementing a type checker, induction, infinite data structures, input and output, lambda-calculus, lazy evaluation, list comprehensions, lists, modules, pattern matching, polymorphism, property-based testing, reasoning about functional programs, recursive functions, sets, strings, tail recursion, trees, tupling, type checking, type inference, types, types and type classes, unification, user-defined types

Overview

• Intermediate Wrap-Up

• User-Defined Types / Trees

• Input and Output

Functions You Should Know

infix operators and special syntax

```
(<=), (<), (==), (>=), (>), (||), (-), (,), (:), (/=), (.), (*), (&&), (+), (++), [], [m..n]
```

other Prelude functions

abs, compare, concat, const, div, drop, error, even, filter, foldr, foldr, fromInteger, fromIntegral, fst, head, init, last, length, lines, map, max, min, mod, negate, not, null, product, putStr, putStrLn, read, replicate, reverse, show, showList, showsPrec, signum, snd, splitAt, sum, tail, take, unlines, unwords, words, zip, zipWith

• other Prelude constants

False, otherwise, True

other functions

Data.Char.chr, Data.Char.isDigit, Data.Char.ord, System.Environment.getArgs

Syntax You Should Recognize

 anonymous functions / functions without names $(\x -> 2 * x)$ -- an anonymous function for doubling

infix operators and sections

$$(+) = (\x y \rightarrow x + y) \qquad \text{infix to prefix}$$

$$x `f` y = f x y \qquad \text{prefix to infix}$$

$$(a >) = (\x \rightarrow a > x) \qquad \text{argument smaller than } a?$$

$$(> b) = (\x \rightarrow x > b) \qquad \text{argument greater than } b?$$

 patterns and guards headIfPositive xs = case xs of x: | x > 0 -> x

 list comprehensions $== [x \mid x \leftarrow xs, p x]$ filter p xs == [f x | x < -xs]map f xs concat (map f xs) == [y | x <- xs, y <- f x] map $(\x -> map ((,) x) ys) xs ==$ $[(x, y) \mid x \leftarrow xs, y \leftarrow ys]$

Types and Type Classes

```
• type signatures - annotate functions by types
range :: Int -> Int -> [Int]
range m n | m > n = []
```

type synonyms - mnemonic names for types
 type Height = Int
 type Width = Int

• type classes and class constraints – for every function f, specific to class C, type inference adds a C-constraint to type

| otherwise = m : range (m + 1) n

Example – Type Constraints

 without type signature, we get ghci> :t range

range :: (Ord a, Num a) => a -> a -> [a]

- m > n, hence m and n of class Ord and m and n of same type
- m + 1, hence m of class Num
- m and n of same type, hence n of class Num

Equational Reasoning

- a function definition in Haskell is a (set of conditional) equation(s)
- if conditions are met, we may "replace equals by equals"
- in this way we may evaluate function calls by applying equations stepwise, until we reach final result

Kinds of Conditions

- "if b then t else e" is t, when b is true; and e, otherwise
- "case e of { $p_1 \rightarrow e_1$; ...; $p_n \rightarrow e_n$ }" is e_i , if e first matches p_i

Primitive Operations

- for primitive operations (like (+), (*), ...), we assume predefined equations
- e.g., 1 + 2 = 3, 0 * 10 = 0, ...

Examples – Equational Reasoning

```
    definition

 zip (x:xs) (y:ys) = (x, y) : zip xs ys
 zip _
evaluate zip [1,2,3] ['a','b']

    definition

  factorial n \mid n \le 1 = 1
                | otherwise = n * factorial (n - 1)

    evaluate factorial 3

    definition

 head xs = case xs of x:_ -> x

    evaluate head "ab"

    definition

 prod xs = if null xs then 1
             else head xs * prod (tail xs)
• evaluate prod [5,6]
```

live in different name spaces

- data $T \alpha_1 \cdots \alpha_n = C_1 \tau_{11} \cdots \tau_{1m_1}$ \vdots $\mid C_k \tau_{k1} \cdots \tau_{km_k}$
- where T is the name of the new type (starting with a capital letter) taking n type parameters α_1 to α_n
- and C_i is the name of the *i*-th (data) constructor, taking m_i arguments of types τ_{i1} to τ_{im_i} (which may contain only type variables among α_1 to α_n)

Examples

• data Bool = False | True

data Pair a b = Pair a b

• data List a = Nil | Cons a (List a)

Data Declarations / Algebraic Data Types

• data List a = Nii | Cons a (Li

constructors and type names

User-Defined Types / Trees

Automatically Deriving Type Class Instances

- for some type classes it is possible to automatically derive instances for algebraic data types
- e.g.,
 data List a = Nil | Cons a (List a)
 deriving (Eq, Show, Read)
- now, we are able to use (==), show, and read for Lists

Examples

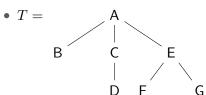
```
ghci> Nil == Cons 1 Nil
False
ghci> show (Cons 1 (Cons 2 Nil))
"Cons 1 (Cons 2 Nil)"
ghci> read it :: List Int
Cons 1 (Cons 2 Nil)
```

Definition – Tree

- (rooted) tree T = (N, E)
- with set of nodes N
- and set of edges/vertices $E \subseteq N \times N$
- unique root of T ($root(T) \in N$) without predecessor
- all other nodes have exactly one predecessor

Example

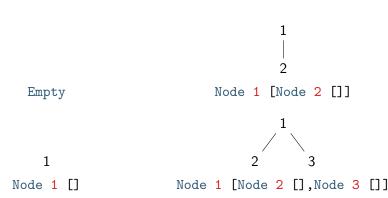
- $N = \{A, B, C, D, E, F, G\}$
- $E = \{(A, B), (A, C), (A, E), (C, D), (E, F), (E, G)\}$
- root(T) = A



Trees in Haskell

- possible type for trees with arbitrary nodes
 data Tree a = Empty | Node a [Tree a]
- a tree is either empty (0 nodes) or there is at least one node with content of type a and an arbitrary number of successor trees

Examples



Binary Trees

- restrict number of successors (maximum 2)
- type
 data BTree a = Empty | Node a (BTree a) (BTree a)
 deriving (Eq, Show, Read)

Functions on Binary Trees

- size number of nodes
 size :: BTree a -> Integer
 size Empty = 0
 size (Node _ l r) = size l + size r + 1
- height length of longest path from root to some leaf
 height :: BTree a -> Integer
 height Empty = 0
 height (Node _ l r) = max (height l) (height r) + 1

Creating Trees from Lists

```
    the easy way

 fromList [] = Empty
 from List (x:xs) = Node x Empty (from List xs)

    the fair way

 make [] = Empty
 make xs = Node z (make ys) (make zs)
   where
                 = length xs 'div' 2
      (ys, z:zs) = splitAt m xs

    the orderly way

 searchTree = foldr insert Empty
   where
      insert x Empty = Node x Empty Empty
      insert x (Node y l r)
        | x < y = Node y (insert x 1) r
        | otherwise = Node y l (insert x r)
```

Transforming Trees into Lists

```
flatten Empty = []
flatten (Node x l r) = flatten l ++ [x] ++ flatten r
```

A Sorting Algorithm for Lists

```
sort = flatten . searchTree
```

CS,HZ,EZ (DCS @ UIBK)

An Initial Example

```
• write the file welcomeIO.hs
main = do

putStrLn "Greetings! What's your name?"
name <- getLine
putStrLn (
    "Welcome to Haskell's IO, " ++ name ++ "!")</pre>
```

- compile it with GHC via
 - \$ ghc --make welcomeIO.hs
- and run it
 - \$./welcomeIO
 - Greetings! What's your name?

Notes

- putStrLn prints a string + newline
- getLine reads a line from standard input
- new syntax: do and <-

IO and the Type System

- consider ghci> :load welcomeIO.hs
 - ghci> :t putStrLn

putStrLn :: String -> IO ()
ghci> :t getLine
getLine :: IO String

ghci> :t main
main :: IO ()

• IO a is type of IO actions delivering results of type a (in addition to their IO operations)

• String -> IO () - after supplying a string, we obtain an IO action

Examples

- (in the case of putStrLn, "printing")
- IO () just IO (in the case of main, run our program)
- IO String do some IO and deliver a string (in the case of getLine, the user-input)

Further Notes

- IO actions (everything of type IO a) are just descriptions of what should be done; nothing is actually done at time of specification
- only main may start execution of IO actions
- inside IO actions, order is important; IO actions are executed in order of appearance (once execution starts); the result of a sequence of IO actions is the result of the last action
- inside IO actions, x <- action (where action :: IO a) may be used to bind the result value of action (which has type a) to the name x (but seriously, this is actually only done, once execution starts)
- $x \leftarrow a$ is not available outside IO actions

Implications

- once we are inside an IO action, we cannot escape
 - strict separation between purely functional code and IO
 - when IO a does not appear inside type signature, we can be absolutely sure that no IO ("side-effect") is performed

Using Pure Code Inside IO Actions

• consider the program reply.hs reply :: String -> String reply name = "Pleased to meet you, " ++ name ++ ".\n" ++ "Your name contains " ++ n ++ " characters." where n = show \$ length namemain :: IO () main = doputStrLn "Greetings again. What's your name?" name <- getLine</pre> let niceReply = reply name putStrLn niceReply

 that is, we may use let x = f (there is no in here!) to bind the result of the pure function f to the name x

Some Simple IO Functions

- return :: a -> IO a turn anything into an IO action
- System.Environment.getArgs :: IO [String] get command line arguments
- putChar :: Char -> IO () print character
- putStr :: String -> IO () print string
- putStrLn :: String -> IO () print string + newline
- getChar :: IO Char read single character from stdin
- getLine :: IO String read line (excluding newline)
- interact :: (String -> String) -> IO () use function that gets input as string and produces output as string
- type FilePath = String
- readFile :: FilePath -> IO String read file content
- writeFile :: FilePath -> String -> IO ()
- appendFile :: FilePath -> String -> IO ()

Examples – Imitating Some GNU Commands

```
    cat.hs - print file contents

 main = do
    [file] <- getArgs
    s <- readFile file
    putStr s

    wc.hs - count newlines/words/characters in input

  count s = ns ++ " " ++ ws ++ " " ++ bs ++ "\n"
    where ns = show $ length $ lines s
          ws = show $ length $ words s
          bs = show $ length s
```

```
main = interact count
```

- uniq.hs omit repeated lines of input
 main = interact (unlines . nub . lines)
- sort.hs sort input
 main = interact (unlines . sort . lines)

Notes

• getArgs :: IO [String] is in System.Environment

foreach :: [a] -> (a -> IO ()) -> IO ()

- nub :: Eq a => [a] -> [a] is in Data.List; eliminates duplicates
- sort :: Ord a => [a] -> [a] is in Data.List; sorts a list

Do Some IO Action for Each Argument

```
• foreach [] io = return ()
  foreach (a:as) io = do { io a; foreach as io }
• better cat.hs
  main = do
    files <- getArgs
  foreach files readAndPrint
  where readAndPrint file = do
        s <- readFile file
        putStr s</pre>
```

Exercises (for November 10th)

- 1. Read Chapter 7 of Real World Haskell.
- 2. Evaluate the two function calls foldr (-) 0 [1,2,3] and foldl (-) 0 [1,2,3] by equational reasoning using the definitions: foldr f b []

```
foldl f b = b
foldl f b (x:xs) = foldl f (b `f` x) xs
```

foldr f b (x:xs) = x `f` foldr f b xs

- 3. Implement the predicate isSorted :: Ord a => BTree a -> Bool, checking whether the given tree is a search tree.
- 4. Write a program Grep. hs that, given a string, echos every line of its standard input, containing this string.
- 5. Modify Grep.hs to also print line numbers of matching lines.
- 6. Implement a function showBTree :: Show a => BTree a -> String that gives an ASCII representation of a binary tree.

lecture 4