

This exam consists of **four** exercises. The available points for each item are written in the margin. You need at least 50 points to pass.

- [25] 1 Consider the Haskell functions:
- ```
map :: (a -> b) -> [a] -> [b]
map f [] = []
map f (x:xs) = f x : map f xs

takeChain :: (a -> a -> Bool) -> [a] -> [a]
takeChain p (x:y:ys) | p x y = x : takeChain p (y:ys)
 | otherwise = [x]
takeChain p xs = xs
```
- [9] (a) Use equational reasoning to evaluate `takeChain (<) [1,2,3,1]` and `map (\x -> 0) [1,2,3,1]`, giving each intermediate step.
- [6] (b) Determine whether `map` and `takeChain` are tail recursive and guardedly recursive, respectively.
- [10] (c) Give a tail recursive implementation of `map`.
- [25] 2 Consider `map` from Exercise 1. Use structural induction to prove `map (\x -> x) xs = xs` for all finite lists `xs`.
- [6] (a) Prove the base case, applying a single equation at a time.
- [9] (b) State the induction hypothesis and the statement you have to show in the step case.
- [10] (c) Prove the step case, applying a single equation at a time.
- [25] 3 Consider the  $\lambda$ -term  $t = ((\lambda x. x) (\lambda x. (\lambda y. y))) (((\lambda x. x) (\lambda x. (\lambda y. y)))) a$
- [4] (a) Write  $t$  as compactly as possible, using the notational conventions (for omitting parentheses etc.).
- [10] (b) Stepwise reduce  $t$  to normal form using applicative order reduction.
- [6] (c) Only using the variables  $x$  and  $y$ , give three *different*  $\lambda$ -terms that are  $\alpha$ -equivalent to  $(\lambda xy. y) y$ .
- [5] (d) Give a  $\lambda$ -term that has a  $\beta$ -normal form but also admits an infinite  $\beta$ -reduction.
- [25] 4 Consider the typing environment  $E = P \cup \{f :: \text{Int} \rightarrow \text{Int}, x :: \text{Int}\}$ .
- [9] (a) Use type checking to show that  $E \vdash f (x + 1) :: \text{Int}$ .
- [8] (b) If possible, solve the unification problem  $\alpha_0 \rightarrow \text{List}(\alpha_1) \approx \alpha_1 \rightarrow \alpha_0$  and compute the resulting mgu.
- [8] (c) Apply the typing constraint rules to transform  $E \triangleright f f :: \alpha_0$  into a unification problem.