

All Solutions

2 We prove the property by structural induction over xs .

Base Case ($xs = []$). The base case is shown by the derivation

$$\begin{aligned}
 \text{map } f (\text{map } g []) &= \text{map } f [] && \text{by def. of map} \\
 &= [] && \text{by def. of map} \\
 &= \text{map } (f \circ g) [] && \text{by def. of map}
 \end{aligned}$$

Step Case ($xs = z : zs$). The IH is

$$\text{map } f (\text{map } g zs) = \text{map } (f \circ g) zs$$

We conclude the step case by the derivation

$$\begin{aligned}
 \text{map } f (\text{map } g (z : zs)) &= \text{map } f (g z : \text{map } g zs) && \text{by def. of map} \\
 &= f (g z) : (\text{map } f (\text{map } g zs)) && \text{by def. of map} \\
 &\stackrel{\text{IH}}{=} f (g z) : \text{map } (f \circ g) zs && \\
 &= \text{map } (f \circ g) (z : zs) && \text{by def. of map}
 \end{aligned}$$

3 We prove the property by structural induction over xs .

Base Case ($xs = []$). The base case is shown by the derivation

$$\begin{aligned}
 \text{filter } p (\text{map } f []) &= \text{filter } p [] && \text{by def. of map} \\
 &= [] && \text{by def. of filter} \\
 &= \text{map } f [] && \text{by def. of map} \\
 &= \text{map } f (\text{filter } (p \circ f) []) && \text{by def. of filter}
 \end{aligned}$$

Step Case ($xs = z : zs$). The IH is

$$\text{filter } p (\text{map } f zs) = \text{map } f (\text{filter } (p \circ f) zs)$$

We start the step case by

$$\text{filter } p (\text{map } f (z : zs)) = \text{filter } p (f z : \text{map } f zs) \quad \text{by def. of map}$$

and proceed by a case analysis on whether $p (f z)$:

- Assume $p (f z)$ holds. Then we conclude by

$$\begin{aligned}
 \text{filter } p (f z : \text{map } f zs) &= f z : \text{filter } p (\text{map } f zs) && \text{by def. of filter} \\
 &\stackrel{\text{IH}}{=} f z : \text{map } f (\text{filter } (p \circ f) zs) && \\
 &= \text{map } f (z : \text{filter } (p \circ f) zs) && \text{by def. of map} \\
 &= \text{map } f (\text{filter } (p \circ f) (z : zs)) && \text{by def. of filter}
 \end{aligned}$$

- Assume $p (f z)$ does not hold. Then we conclude by

$$\begin{aligned}
\text{filter } p (f z : \text{map } f zs) &= \text{filter } p (\text{map } f zs) && \text{by def. of filter} \\
&\stackrel{\text{IH}}{=} \text{map } f (\text{filter } (p \circ f) zs) \\
&= \text{map } f (\text{filter } (p \circ f) (z : zs)) && \text{by def. of filter}
\end{aligned}$$

4 We prove the property by structural induction on xs .

Base Case ($xs = []$). The base case is shown by the derivation

$$\begin{aligned}
\text{map } f ([] ++ ys) &= \text{map } f ys && \text{by def. of map} \\
&= [] ++ \text{map } f ys && \text{by def. of ++} \\
&= \text{map } f [] ++ \text{map } f ys && \text{by def. of map}
\end{aligned}$$

Step Case ($xs = z : zs$). The IH is

$$\text{map } f (zs ++ ys) = \text{map } f zs ++ \text{map } f ys$$

We conclude the step case by the derivation

$$\begin{aligned}
\text{map } f ((z : zs) ++ ys) &= \text{map } f (z : (zs ++ ys)) && \text{by def. of ++} \\
&= f z : \text{map } f (zs ++ ys) && \text{by def. of map} \\
&\stackrel{\text{IH}}{=} f z : (\text{map } f zs ++ \text{map } f ys) \\
&= (f z : \text{map } f zs) ++ \text{map } f ys && \text{by def. of ++} \\
&= \text{map } f (z : zs) ++ \text{map } f ys && \text{by def. of map}
\end{aligned}$$

5 If $n < 0$ the property trivially holds. Thus, we restrict n to be a natural number and prove the property for arbitrary xs by induction on n .

Base Case ($n = 0$). The base case is shown by the derivation

$$\begin{aligned}
\text{take } 0 (\text{map } f xs) &= [] && \text{by def. of take} \\
&= \text{map } f [] && \text{by def. of map} \\
&= \text{map } f (\text{take } 0 xs) && \text{by def. of take}
\end{aligned}$$

Step Case ($n = k + 1$). The IH is

$$\forall xs. \text{take } k (\text{map } f xs) = \text{map } f (\text{take } k xs)$$

Now if $xs = []$ the property trivially holds, thus let us assume $xs = z : zs$.

$$\begin{aligned}
\text{take } (k + 1) (\text{map } f (z : zs)) &= \text{take } (k + 1) (f z : \text{map } f zs) && \text{by def. of map} \\
&= f z : \text{take } k (\text{map } f zs) && \text{by def. of take} \\
&\stackrel{\text{IH}}{=} f z : \text{map } f (\text{take } k zs) \\
&= \text{map } f (z : \text{take } k zs) && \text{by def. of map} \\
&= \text{map } f (\text{take } (k + 1) (z : zs)) && \text{by def. of take}
\end{aligned}$$

6 If $n < 0$ the property trivially holds. Thus, we restrict n to be a natural number and and prove the property for arbitrary xs by induction on n .

Base Case ($n = 0$). The base case is shown by the derivation

$$\begin{aligned} \text{take } 0 \text{ } xs \text{ ++ drop } 0 \text{ } xs &= [] \text{ ++ } xs && \text{by def. of take and drop} \\ &= xs && \text{by def. of ++} \end{aligned}$$

Step Case ($n = k + 1$). The IH is

$$\forall xs. \text{take } k \text{ } xs \text{ ++ drop } k \text{ } xs = xs$$

Now if $xs = []$ the property trivially holds, thus let us assume $xs = z : zs$.

$$\begin{aligned} \text{take } (k + 1) \text{ } (z : zs) \text{ ++ drop } (k + 1) \text{ } (z : zs) &= (z : \text{take } k \text{ } zs) \text{ ++ drop } (k + 1) \text{ } (z : zs) && \text{by def. of take} \\ &= z : (\text{take } k \text{ } zs \text{ ++ drop } (k + 1) \text{ } (z : zs)) && \text{by def. of ++} \\ &= z : (\text{take } k \text{ } zs \text{ ++ drop } k \text{ } zs) && \text{by def. of drop} \\ &\stackrel{\text{IH}}{=} z : zs \end{aligned}$$