universität innsbruck

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Functional Programming

Lecture 11

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Topics

abstract data types, algebraic data types, binary search trees, combinator parsing, efficiency, encoding data types as lambda-terms, evaluation strategies, formal verification, first steps, guarded recursion, Haskell introduction, higher-order functions, historical overview, induction, infinite data structures, input and output, lambda-calculus, lazy evaluation, list comprehensions, lists, modules, pattern matching, polymorphism, property-based testing, reasoning about functional programs, recursive functions, sets, strings, tail recursion, trees, tupling, type checking, type inference, types, types and type classes, unification, user-defined types

Overview

- Laziness and Infinite Data Structures
- Examples of (Infinite) Laziness
- Recommendations

Laziness – Motivation

- only compute values needed for final result
- avoid computing same value twice (memoization)

Example

- in the program

```
main = do
  let f1 x = x + 1
      f2 x = f2 x   -- a nonterminating function
  input <- getLine
  let i = read input
  print $ head [f1 i, f2 i]
```

- value of `f2 i` not needed
- however, without laziness program would not terminate

Laziness and Infinite Data Structures Facilitate Modularity

- separation of concerns
- use potentially infinite data structures
- write small functions with specific tasks

Find Index of First List Element Satisfying Property

- function `firstIndex :: (a -> Bool) -> [a] -> Int`
- in Haskell (only using Prelude functions!) `firstIndex p = fst . head . filter (p . snd) . zip [0..]`
- (lazy) evaluation

```
fst . head . filter ((==1) . snd) $ zip [0..] [1..9]
  = ... $ (0,1) : zip [1..] [2..9]
  = fst . head $ (0,1) : filter ((==1) . snd) (...)
  = fst (0,1)
  = 0
```

- without laziness (ignoring nontermination of `[0..]`) three list traversals (generation of `[0 .. length xs]`, `zip`, and `filter`)

Watch out for Memory Leaks

with laziness even tail recursive programs may run out of memory

Function (Tail Recursive)

```
length' acc []      =  acc
length' acc (_:xs)  =  length' (acc + 1) xs
```

Evaluation

```
length' 0 [1,2,3,4]
  =  length' (0+1) [2,3,4]
  =  length' (0+1+1) [3,4]
  =  length' (0+1+1+1) [4]
  =  length' (0+1+1+1+1) []
  =  (0+1+1+1+1)
```


Being Strict – The seq Function

- type `seq :: a -> b -> b`
- ``x `seq` y`` reduces `x` to WHNF and returns `y`

Function (still Tail Recursive)

```
length'' acc []      =  acc
length'' acc (_:xs)  =  acc `seq` length'' (acc + 1) xs
```

Strict Folding of Lists – Data.List.foldl'

```
foldl' _ b []     = b
foldl' f b (x:xs) = c `seq` foldl' f c xs
  where
    c = f b x
```

Example

```
length'' = foldl' (const . (+1))
```


Example 1 – Fibonacci Numbers (this time starting from 0)

$$\mathsf{fib}(i) = \begin{cases} 0 & \text{if } i \leq 0 \\ 1 & \text{if } i = 1 \\ \mathsf{fib}(i-1) + \mathsf{fib}(i-2) & \text{otherwise} \end{cases}$$

Sequence

$0\ 1\ 1\ 2\ 3\ 5\ 8\ 13\ 21\ 34\ 55\ 89\ 144\ 233\ 377\ 610\ 987\ 1597\ 2584\ \ldots$

One Way of Computing Fibonacci Numbers

starting at 0	0	1	1	2	3	5	8	13	21 ...
starting at 1	1	1	2	3	5	8	13	21	...
(+)	1	2	3	5	8	13	21	34	...

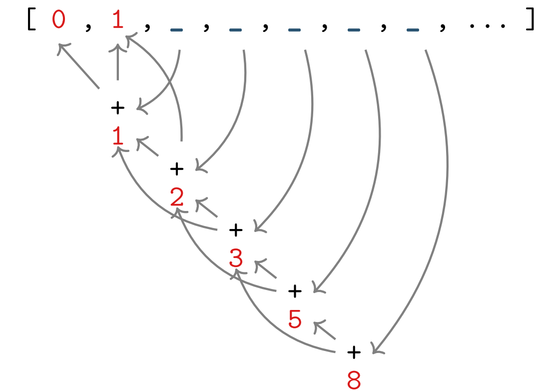
Ingredients

- function to shift sequence to left
- function to add two sequences

Fibonacci Numbers in Haskell

```
fibs :: [Integer]
fibs = 0 : 1 : zipWith (+) fibs (tail fibs)
```

Resulting Data Dependencies



Example 2 – A Variation of the "Sieve of Eratosthenes"

Goal
generate list of all prime numbers

Procedure
start with list of all natural numbers (from 2 on)

1. mark first element x as prime
2. remove all multiples of x
3. go to Step 1

The Sieve in Haskell

```
primes :: [Integer]
primes = sieve [2..]
  where
    sieve (x:xs) =
      x : sieve [y | y <- xs, y `mod` x /= 0]
```

Example 3 – Dropping n Elements from Tail of List

- naive solution: `dropEnd n xs = take (length xs - n) xs`
- problem: does not work on infinite lists
  (nontermination even if only finite prefix of result is used)
- another attempt: `dropEnd n = reverse . drop n . reverse`
- problem: does not work on infinite lists and requires three traversals
- solution: `dropEnd n xs = zipWith const xs (drop n xs)`

Why does this work?

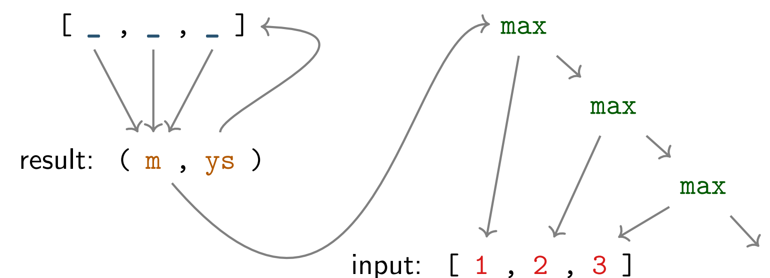
- result of `zipWith f xs ys` has length of shorter of `xs` and `ys`
- `zipWith const xs ys` ignores elements of `ys`
- `zipWith const` $[x_1, \ldots, x_m]$ $[y_1, \ldots, y_n] =$
  $[x_1, \ldots, x_{\min\{m,n\}}]$

Example 4 – Lazy Shenanigans

- task – replace all elements of list by its maximum in single traversal
- in Haskell

```
replaceAllByMax xs = ys
  where
    (m, ys) = maxAndReplaceAll m xs
    maxAndReplaceAll c =
      foldr (\x (m, ys) -> (max x m, c : ys)) (0, [])
```

Resulting Data Dependencies for [1,2,3]



Courses Related to FP in SS2019

- Term Rewriting (VO2 + PS1)

Bachelor Projects Related to FP

- Fast Multiset-Comparison
- Domain Specific Language for Machine Learning Theorem Proving
- Intelligent search-time proof strategy selector for Isabelle/HOL
- Meta-Tool Based Proof Search in Coq
- Alternative Approaches to Termination Analysis
- Maximal Interpretations
- A Certified Decision Procedure for Termination of Right-Ground Term Rewrite Systems
- Commutation of Term Rewrite Systems
- Unification in Booltool
- ...